

HP Assembler

Programmer's Reference Manual



HEWLETT-PACKARD COMPANY 11000 WOLFE ROAD, CUPERTINO, CALIFORNIA, 95014

Part No. 02116-9014

Printed in U.S.A. 8/75

HP Computer Museum www.hpmuseum.net

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PREFACE

This publication is the reference manual for the Hewlett-Packard Assembly Language for the 2100 family of computers. Since Hewlett-Packard provides assemblers with all of its operating systems, this manual covers only the specifications of assembly language, not operating procedures for the assemblers. The user should refer to the appropriate system manual or operator's guide listed below:

SOFTWARE OPERATING PROCEDURES	
SIO SUBSYSTEMS module	(5951-1390)
DISC OPERATING SYSTEM	(02116-91748)
MOVING-HEAD DISC OPERATING SYSTEM	(02116-91779)
MAGNETIC TAPE SYSTEM	(02116-91752)

In addition, the Formatter and other relocatable subroutines that can be called by relocatable assembly language programs are described in full in the RELOCATABLE SUBROUTINES manual (02116-91780). Interaction between relocatable programs and operating systems is described in:

BASIC CONTROL SYSTEM	(02116-9017)
MOVING-HEAD DISC OPERATING SYSTEM	(02116-91779)
DISC OPERATING SYSTEM	(02116-91748)
MAGNETIC TAPE SYSTEM	(02116-91752)

Interaction between absolute programs and SIO drivers is described in an appendix to this book.

NEW AND CHANGED INFORMATION

All known errors in this manual have been corrected. In addition, the Assembler operating procedures (formally in Section V) are now contained in the SOFTWARE OPERATING PROCEDURES, SIO SUBSYSTEMS module (5951-1390).

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INTRODUCTION

The Assembler and the Extended Assembler translate symbolic source language instructions into an object program for execution on the computer. The source language provides mnemonic machine operation codes, assembler directing pseudo codes, and symbolic addressing. The assembled program may be absolute or relocatable.

The source program may be assembled as a complete entity or it may be subdivided into several relocatable subprograms (or a main program and several subroutines), each of which may be assembled separately. The relocating loader loads the program and Links the subprograms as required. The Basic Binary Loader or Basic Binary Disc Loader loads absolute programs.

Input for the Assembler is prepared on paper tape or cards; the Assembler punches the binary program on paper tape in a format acceptable to the loader.

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SECTION I GENERAL DESCRIPTION

ASSEMBLY PROCESSING

The Assembler is a two pass system, or, if both punch and list output are requested, a three pass system on a minimum configuration. A pass is defined as a processing cycle of the source program input.

In the first pass, the Assembler creates a symbol table from the names used in the source statements. It also checks for certain possible error conditions and generates diagnostic messages if necessary.

During pass two, the Assembler again examines each statement in the source program along with the symbol table and produces the binary program and a program listing. Additional diagnostic messages may also be produced.

If only the output device is available and if both the binary output and the list output are requested, the listing function is deferred and performed as pass three.

When using the Assembler with a mass storage device the source program is written on the device during the first pass; the second pass of the source is read from the mass storage.

SYMBOLIC ADDRESSING

Symbols may be used for referring to machine instructions, data, constants, and certain other pseudo operations. A symbol represents the address for a computer word in memory. A symbol is defined when it is used as a label for a location in the program, a name of a common storage segment, the label of a data storage area or constant, the label of an absolute or relocatable value, or a location external to the program.

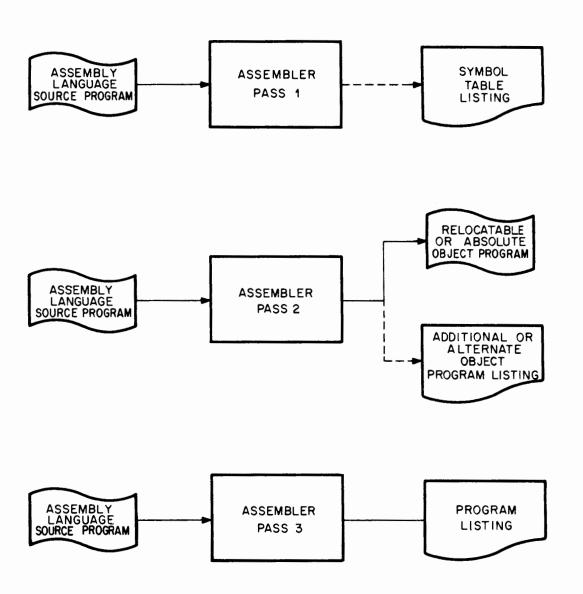


Figure 1-1. HP Assembler Processing

GENERAL DESCRIPTION

Through use of simple arithmetic operators, symbols may be combined with other symbols or numbers to form an expression which may identify a location other than that specifically named by a symbol. Symbols appearing in operand expressions, but not specifically defined, and symbols that are defined more than once are considered to be in error by the Assembler.

PROGRAM RELOCATION

Relocatable programs may be relocated in core by the relocating loader; the location of the program origin and all subsequent instructions is determined at the time the program is loaded.

A relocatable program is assembled assuming a starting location of zero. All other instructions and data areas are assembled relative to this zero base. When the program is loaded, the relocatable operands are adjusted to correspond with the actual locations assigned by the loader.

The starting locations of the common storage area and the base page portion of the program are always established by the loader. References to the common area are common relocatable. References to the base page portion of the program are base page relocatable. If a program refers to the common area or makes use of the base page via the ORB pseudo instruction, the program must also be relocatable.

If a program is to be relocatable, all subprograms comprising the program must be relocatable; all memory reference operands must be relocatable expressions or literals, or have an absolute value of less than 100_{o} .

PROGRAM LOCATION COUNTERS

The Assembler maintains a counter, called the program location counter, that assigns consecutive memory addresses to source statements.

GENERAL DESCRIPTION

The initial value of the program location counter is established according to the use of either the NAM or ORG pseudo operation at the start of the program. The NAM operation causes the program location counter to be set to zero for a relocatable program; the ORG operation specifies the absolute starting location for an absolute program.

Through use of the ORB pseudo operation a relocatable program may specify that certain operations or data areas be allocated to the base page. If so, a separate counter, called the base page location counter, is used in assigning these locations.

ASSEMBLY OPTIONS

Parameters specified with the first statement, the control statement, define the output to be produced by the Assembler:[†]

- Absolute The addresses generated by the Assembler are to be interpreted as absolute locations in memory. The program is a complete entity; external symbols, common storage references, and entry points are not permitted.
- Relocatable The program may be located anywhere in memory. All operands which refer to memory locations are adjusted as the program is loaded. Operands, other than those referring to the first 64 locations, must be relocatable expressions. Subprograms may contain external symbols and entry points, and may refer to common storage.
- Binary output An absolute or relocatable program is to be punched on paper tape.
- List output A program listing is produced either during pass two or pass three.

†See Section V for complete details.

Table print - List the symbol table at the end of the first pass.

Selective assembly - Sections of the program may be included or excluded at assembly time depending on the option used.

SECTION II INSTRUCTION FORMAT

A source language statement consists of a label, an operation code, an operand, and comments. The label is used when needed as a reference by other statements. The operation code may be a mnemonic machine operation or an assembly directing pseudo code. An operand may be an expression consisting of an alphanumeric symbol, a number, a special character, or any of these combined by arithmetic operations. (For the Extended Assembler, an operand may also be a literal.) Indicators may be appended to the operand to specify certain functions such as indirect addressing. The comments portion of the statement is optional.



STATEMENT CHARACTERISTICS

The fields of the source statement appear in the following order:

Label Opcode Operand Comments

Field Delimiters

One or more spaces separate the fields of a statement. An end-of-statement mark terminates the entire statement. On paper tape this mark is a return, (CR), and line feed, (LF). + A single space following the end-of-statement mark from the previous source statement is the null field indicator of the label field.

+A circled symbol (e.g., (CR)) represents an ASCII code or Teleprinter key.

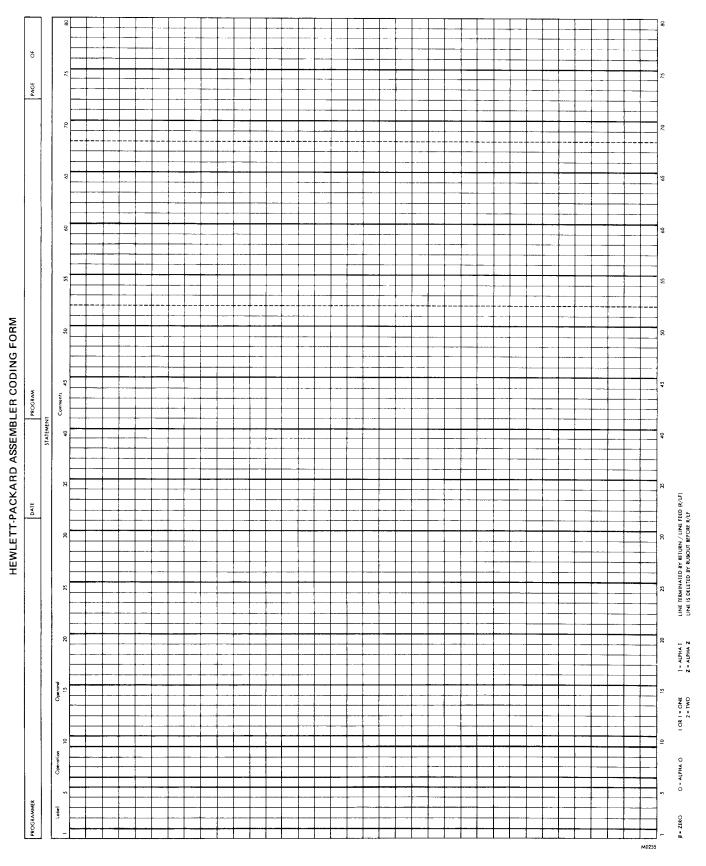


Figure 2-1. Sample Coding Form (Actual Size 11 x 13-1/2)

Character Set

The characters that may appear in a statement are as follows:

- A through Z
- 0 through 9
- . (period)
 - * (asterisk)
 - + (plus)
 - (minus)
 - , (comma)
 - = (equals)
 - () (parentheses)
 - (space)

Any other ASCII characters may appear in the Remarks field. (See Appendix A.)

The letters A through Z, the numbers O through 9, and the period may be used in an alphanumeric symbol. In the first position in the Label field, an asterisk indicates a comment; in the Operand field, it represents the value of the program location counter for the current instruction. The plus and minus are used as operators in arithmetic address expressions. The comma separates several operation codes, or an expression and an indicator in the Operand field. An equals sign indicates a literal value. The parentheses are used only in the COM pseudo instruction.

Spaces separate fields of a statement. They may also be used to establish the format of the output list. Within a field they may be used freely when following +, -, ,, or (.

STATEMENT LENGTH

A statement may contain up to 80 characters including blanks, but excluding the end-of-statement mark. Fields beginning in characters 73 - 80 are not processed by the Assembler.

LABEL FIELD

The Label field identifies the statement and may be used as a reference point by other statements in the program.

The field starts in position one of the statement; the first position following an end-of-statement mark for the preceding statement. It is terminated by a space. A space in position one is the null field indicator for the label field; the statement is unlabeled.

Label Symbol

A label must be symbolic. It may have one to five characters consisting of A through Z, O through 9, and the period. The first character must be alphabetic or a period. A label of more than five characters could be entered on the source language tape, but the Assembler flags this condition as an error and truncates the label from the right to five characters.

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Examples:

[†]The caret symbol, $_{\scriptscriptstyle \wedge}$, indicates the presence of a space.

Each label must be unique within the program; two or more statements may not have the same symbolic name. Names which appear in the Operand field of an EXT or COM pseudo instruction may not also be used as statement labels in the same subprogram.

Examples:

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Asterisk

An asterisk in position one indicates that the entire statement is a comment. Positions 2 through 80 are available; however, positions 1 through 68 only are printed as part of the assembly listing on the 2752A Teleprinter. An asterisk within the Label field is illegal in any position other than one.

OPCODE FIELD

The operation code defines an operation to be performed by the computer or the Assembler. The Opcode field follows the Label field and is separated from it by at least one space. If there is no label, the operation code may

begin anywhere after position one. The Opcode field is terminated by a space immediately following an operation code. Operation codes are organized in the following categories:

Machine operation codes Memory Reference Register Reference Input/Output, Overflow, and Halt Extended Arithmetic Unit Pseudo operation codes Assembler control Object program linkage Address and symbol definition Constant definition Storage allocation Arithmetic subroutine calls Assembly Listing Control (Extended Assembler) Operation codes are discussed in detail in Sections III and IV.

OPERAND FIELD

The meaning and format of the Operand field depend on the type of operation code used in the source statement. The field follows the Opcode field and is separated from it by at least one space. It is terminated by a space except when the space follows, + - (or, if there are no comments, by an end-of-statement mark.

The Operand field may contain an expression consisting of one of the following:

Single symbolic term Single numeric term Asterisk Combination of symbolic terms, numeric terms, and the asterisk jointed by the arithmetic operators + and -. An expression may be followed by a comma and an indicator.

Programs being assembled by the Extended Assembler may also contain a literal value in the Operand field.

Symbolic Terms

A symbolic term may be one to five characters consisting of A through Z, O through 9, and the period. The first character must be alphabetic or a period.

Examples:

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A symbol used in the Operand field must be a symbol that is defined elsewhere in the program in one of the following ways:

As a label in the Label field of a machine operation

As a label in the Label field of a BSS, ASC, DEC, DEX, OCT, DEF, ABS, EQU or REP pseudo operation

As a name in the Operand field of a COM or EXT pseudo operation As a label in the Label field of an arithmetic subroutine pseudo operation

The value of a symbol is absolute or relocatable depending on the assembly option selected by the user. The Assembler assigns a value to a symbol as it appears in one of the above fields of a statement. If a program is to be loaded in absolute form, the values assigned by the assembler remain fixed. If the program is to be relocated, the actual value of a symbol is established on loading. A symbol may also be made absolute through use of the EQU pseudo instruction.

A symbolic term may be preceded by a plus or minus sign. If preceded by a plus or no sign, the symbol refers to its associated value. If preceded by a minus sign, the symbol refers to the two's complement of its associated value. A single negative symbolic operand may be used only with the ABS pseudo operation.

Numeric Terms

A numeric term may be decimal or octal. A decimal number is represented by one to five digits within the range 0 to 32767. An octal number is represented by one to six octal digits followed by the letter B; (0 to 177777B).

If a numeric term is preceded by a plus or no sign, the binary equivalent of the number is used in the object code. If preceded by a minus sign, the two's complement of the binary equivalent is used. A negative numeric operand may be used only with the DEX, DEC, OCT, and ABS pseudo operations.

In an absolute program, the maximum value of a numeric operand depends on the type of machine or pseudo instruction. In a relocatable program, the value of a numeric operand may not exceed 77B. Numeric operands are absolute. Their value is not altered by the assembler or the loader.

Asterisk

An asterisk in the Operand field refers to the value in the program location counter (or base page location counter) at the time the source program statement is encountered. The asterisk is considered a relocatable term in a relocatable program.

Expression Operators

The asterisk, symbols, and numbers may be joined by the arithmetic operators + and - to form arithmetic address expressions. The Assembler evaluates an expression and produces an absolute or relocatable value in the object code.

Examples:

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Evaluation of Expressions

An expression consisting of a single operand has the value of that operand. An expression consisting of more than one operand is reduced to a single value. In expressions containing more than one operator, evaluation of the expression proceeds from left to right. The algebraic expression A-(B-C+5) must be represented in the Operand field as A-B+C-5. Parentheses are not permitted in operand expressions for the grouping of operands.

The range of values that may result from an operand expression depends on the type of operation. The Assembler evaluates expressions as follows:

Pseudo Operations	modulo 2 ¹⁵ -1
Memory Reference	modulo 2 ¹⁰ -1
Input/Output	2^6 - 1 (maximum value)

Expression Terms

The terms of an expression are the numbers and the symbols appearing in it. Decimal and octal integers, and symbols defined as being absolute in an EQU pseudo operation are absolute terms. The asterisk and all symbols that are defined in the program are relocatable or absolute depending on the type of assembly. Symbols that are defined as external may appear only as single term expressions.

Within a relocatable program, terms may be program relocatable, base page relocatable, or common relocatable. A symbol that names an area of common storage is a common relocatable term. A symbol that is allocated to the base page is a base page relocatable term. A symbol that is defined in any

⁺The evaluation of expressions by the Assembler is compatible with the addressing capability of the hardware instructions (e.g., up to 32K words through Indirect Addressing). The user must take care not to create addresses which exceed the memory size of the particular configuration.

other statement is a program relocatable term. Within one expression all relocatable terms must be base page relocatable, program relocatable, or common relocatable; the three types may not be mixed.

Absolute and Relocatable Expressions

An expression is absolute if its value is unaffected by program relocation. An expression is relocatable if its value changes according to the location into which the program is loaded. In an absolute program, all expressions are absolute. In a relocatable program, an expression may be base page relocatable, program relocatable, common relocatable, or absolute (if less than 100_{o}) depending on the definition of the terms composing it.

ABSOLUTE EXPRESSIONS

An absolute expression may be any arithmetic combination of absolute terms. It may also contain relocatable terms alone, or in combination with absolute terms. If relocatable terms do appear, there must be an even number of them; they must be of the same type; and they must be paired by sign (a negative term for each positive term). The paired terms do not have to be contiguous in the expression. The pairing of terms by type cancels the effect of relocation; the value represented by the pair remains constant.

An absolute expression reduces to a single absolute value. The value of an absolute multiterm expression may be negative only for ABS pseudo operations. A single numeric term also may be negative in an OCT, DEX, or DEC pseudo instruction. In a relocatable program the value of an absolute expression must be less than 100₈ for instructions that reference memory locations (Memory Reference, DEF, Arithmetic subroutine calls).

Examples:

If P_1 and P_2 are program relocatable terms; B_1 and B_2 , base page relocatable; C_1 and C_2 , common relocatable; and A, an absolute term; then the following are absolute terms:

The asterisk is base page relocatable or program relocatable depending on the location of the instruction.

RELOCATABLE EXPRESSIONS

A relocatable expression is one whose value is changed by the loader. All relocatable expressions must have a positive value.

A relocatable expression may contain any odd number of relocatable terms, alone, or in combination with absolute terms. All relocatable terms must be of the same type. Terms must be paired by sign with the odd term being positive.

A relocatable expression reduces to a single positive relocatable term, adjusted by the values represented by the absolute terms and paired relocatable terms associated with it.

Examples:

If P_1 , P_2 , and P_3 are program relocatable terms; B_1 , B_2 , and B_3 base page relocatable; C_1 , C_2 and C_3 , common relocatable; and A, an absolute term; then the following are relocatable terms:

Pl-A	C ₁ -A	^B 1 ^{+A}
^P 1 ^{-P} 2 ^{+P} 3	^C 1 ^{-C} 2 ^{+C} 3	Cl+y
*+A	*-P1+P2	*-A
A+B1	A+C1	-A-P1+P2+P3
^B 1 ^{-B} 2 ^{+B} 3 ^{-A}	^C 1 ^{-C} 2 ^{+C} 3 ^{-A}	A+*
*+P1 - *	^P 1 ^{-P} 2 ^{+*}	-C1+C2+C3

Literals

Actual literal values may be specified as operands in relocatable programs to be assembled by the Extended Assembler. The Extended Assembler converts the literal to its binary value, assigns an address to it, and substitutes this address as the operand. Locations assigned to literals are those immediately following the last location used by the program.

A literal is specified by using an equal sign and a one-character identifier defining the type of literal. The actual literal value is specified immediately following this identifier; no spaces may intervene.

The identifiers are:

- =D a decimal integer, in the range -32767 to 32767, including zero.⁺
- =F a floating point number; any positive or negative real number in the range 10^{-38} to 10^{38} , including zero.[†]
- =B an octal integer, one to six digits, $b_1b_2b_3b_4b_5b_6'$, where b_1 may be 0 or 1, and b_2-b_7 may be 0 to 7.⁺
- =A two ASCII characters.†
- =L an expression which, when evaluated, will result in an absolute value. All symbols appearing in the expression must be previously defined.

[†] See CONSTANT DEFINITION, Section 4.

If the same literal is used in more than one instruction, only one value is generated, and all instructions using this literal refer to the same location.

Literals may be specified only in the following memory reference instructions and pseudo instructions:

ADA	ADB	AND	мру)	
LDA	LDB	XOR	DIV 👌	may use =D, =B, =A, =L
CPA	CPB	IOR	y	
DLD	FAD			
FMP	FSB >	may us	e ≠F	
FDV)			

Examples:

LDA	=D798Ø	A-Register is loaded with the binary equivalent of
		^{798ø} 1ø.
IOR	=B777	Inclusive OR is performed with contents of A-Register
		and 777 ₈ .
LDA	=ANO	A-Register is loaded with binary representation of
		ASCII characters NO.
LDB	=LZETZ-ZOO	M+68 B-Register is loaded with the value resulting
		from the absolute expression.
FMP	=F39.75	Contents of A- and B-Registers multiplied by float-
		ing point constant 39.75.

Indirect Addressing

The HP computers provide an indirect addressing capability for memory reference instructions. The operand portion of an indirect instruction contains an address of another location rather than an actual operand. The

secondary location may be the operand or it may be indirect also and give yet another location, and so forth. The chaining ceases when a location is encountered that does not contain an indirect address. Indirect addressing provides a simplified method of address modifications as well as allowing access to any location in core.

The Assembler allows specification of indirect addressing by appending a comma and the letter I to any memory reference operand other than one referring to an external symbol. The actual operand of the instruction may be given in a DEF pseudo operation; this pseudo operation may also be used to indicate further levels of indirect addressing.



Examples:

[Lab	el.		5	Operation 10				Operand 15						2	0			25								30	,		35								Comments 40 45										50									
4	NΒ				Ι		L	C	A	T		S	A	N	N	,	Ι							E	A	С	ŀ	ł		T	I	Μ	E	Γ	T	T	Η	Ε		I	S	Z	!	I	S	5	E			E (cl	U	T	E	D	,	Γ	
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Γ	Τ			Τ	T			•	Γ	T			Γ	T	T					Ī	Γ	T	1										Γ		Τ	T		_		Γ	Γ			1		Τ		T	T		T				Π		Γ	
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Γ	Г				Τ				Τ	T				Τ	T						Γ	T	T					T							T								Τ			Γ					1	1					Γ	1
Γ	T			T	1				Τ	T				T	1							T	1				T	T						1	T	T					Γ	1			T	T		T	1	T	1	1			Π		T	
F	1	1	+	+	+		-	t	+	+	-		+	+	+	+	_			1	t	+	+		-	1	1-	t	1			+	1	+	t	+	+		-	1	t	t	+	+	1	1	+	+	t	+	+	-+	-			-	t	+

A relocatable assembly language program, however, may be designed without concern for the pages in which it will be stored; indirect addressing is not required in the source language. When the program is being loaded, the loader provides indirect addressing whenever it detects an operand which does not fall in the current page or the base page. The loader substitutes a reference to the base page and then stores an indirect address in this referenced location. References to the same operand from other pages will be linked through the same location in the base page.

Base Page Addressing

The computer provides a capability which allows the memory reference instructions to address either the current page or the base page. The Assembler or the loader adjusts all instructions in which the operands refer to the base page; specific notation defining an operand as a base page reference is not required in the source program.

Clear Flag Indicator

The majority of the input/output instructions can alter the status of the input/output interrupt flag after execution or after the particular test is performed. In source language, this function is selected by appending a comma and a letter C to the Operand field.

Examples:

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	Γ	Γ		T	Τ	Τ	S	T	C	T	Ţ	I	0	7	Ι,	0					C	L	E	A	F	R		F	L	A	G	Ι	0	7	ľ	A	F	T	E	R		C	0	N	1	r	R	0	L	Π		Г	Т	T
				T	T	1				T	T				Γ	T	T	T			В	I	T	·	1	ſ	s		S	Ε	Т					t			T	T	t			T	T	1	1			Π		Γ	Τ	t
F	F	1	Ţ	T	T	1	0	T	B		Ţ	I	0	5	Ι,	C	Ţ	1			С	L	E	A	ſ	R		F	L	A	G	Ι	0	5		A	F	T	E	R		M	0		/ E		1		-	Π	Γ	F	T	t
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Γ					T	1			1	T	t	1	-				T	T					Ì		T	1	1							Γ	T	T			1		Γ		T	T	T	1	1						T	t

COMMENTS FIELD

The Comments field allows the user to transcribe notes on the program that will be listed with source language coding on the output produced by the Assembler. The field follows the Operand field and is separated from it by at least one space. The end-of-statement mark, (CR) (LF), or the 80th character in the entire statement terminates the field. If the listing is to be produced on the 2752A Teleprinter, the total statement length, excluding

the end-of-statement mark, should not exceed 52 characters, the width of the source language portion of the listing. Statements consisting solely of comments may contain up to 68 characters including the asterisk in the first position. On the list output, statements consisting entirely of comments begin in position 5 rather than 21 as with other source statements.

If there is no operand present, the Comments field should be omitted in the NAM and END pseudo operations and in the input/output statements, SOC, SOS, and HLT. If a comment is used, the Assembler attempts to interpret it as an operand.

SECTION III MACHINE INSTRUCTIONS

The HP Assembler language machine instruction codes take the form of threeletter mnemonics. Each source statement corresponds to a machine operation in the object program produced by the Assembler.

Notation used in representing source language instruction is as follows:

label	Optional statement label
m	Memory location an expression
I .	Indirect addressing indicator
sc	Select code an expression
С	Clear interrupt flag indicator
comments	Optional comments
[]	Brackets defining a field or portion of a field that is optional
{ }	Brackets indicating that one of the set may be selected.
lit	literal

MEMORY REFERENCE

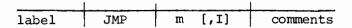
Memory reference instructions perform arithmetic, logical and jump operations on the contents of the locations in core and the registers. An instruction may directly address the 2048 words of the current and base pages. If required, indirect addressing may be utilized to refer to all 32,768 words of memory. Expressions in the Operand field are evaluated modulo 2¹⁰.

If the program is to be assembled in relocatable form, the Operand field may contain relocatable expressions or absolute expressions which are less than 100₈ in value. If the program is to be absolute, the operands may be any expressions consistent with the location of the program. Literals may not be used in an absolute program. Absolute programs must be complete entities; they may not refer to external subroutines or common storage.

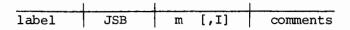
MACHINE INSTRUCTIONS

Jump and Increment-Skip

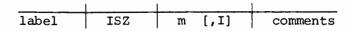
Jump and Increment-Skip instructions may alter the normal sequence of program execution.



Jump to m. Jump indirect inhibits interrupt until the transfer of control is complete.



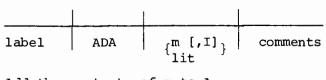
Jump to subroutine. The address for label+1 is placed into the location represented by m and control transfers to m+1. On completion of the sub-routine, control may be returned to the normal sequence by performing a JMP m, I.



Increment, then skip if zero. ISZ adds 1 to the contents of m. If m then equals zero, the next instruction in memory is bypassed.

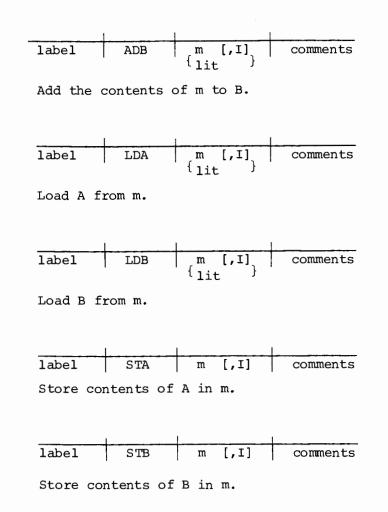
Add, Load and Store

Add, Load, and Store instructions transmit and alter the contents of memory and of the A- and B-Registers. A literal, indicated by "lit", may be either =D, =B, =A, or =I type.



Add the contents of m to A.

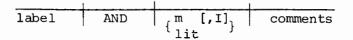
MACHINE INSTRUCTIONS



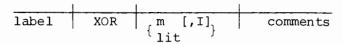
In each instruction, the contents of the sending location is unchanged after execution.

Logical Operations

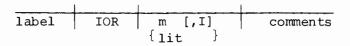
The logical instructions allow bit manipulation and the comparison of two computer words.



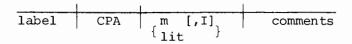
The logical product of the contents of m and the contents of A are placed in A.



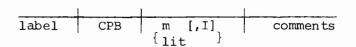
The modulo-two sum (exclusive "or") of the bits in m and the bits in A is placed in A.



The logical sum (inclusive "or") of the bits in m and the bits in A is placed in A.



Compare the contents of m with the contents of A. If they differ, skip the next instruction; otherwise, continue.



Compare the contents of m with the contents of B. If they differ, skip the next instruction; otherwise, continue.

REGISTER REFERENCE

The register reference instructions include a shift-rotate group, an alterskip group, and NOP (no-operation). With the exception of NOP, they have the capability of causing several actions to take place during one memory cycle. Multiple operations within a statement are separated by a comma.

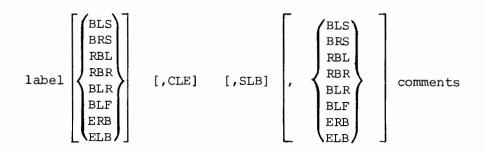
Shift-Rotate Group

This group contains 19 basic instructions that can be combined to produce more than 500 different single cycle operations.

CLE	Clear E to zero
ALS	Shift A left one bit, zero to least significant bit. Sign unaltered
BLS	Shift B left one bit, zero to least significant bit. Sign unaltered
ARS	Shift A right one bit, extend sign; sign unaltered.
BRS	Shift B right one bit, extend sign; sign unaltered.
RAL	Rotate A left one bit
RBL	Rotate B left one bit
RAR	Rotate A right one bit
RBR	Rotate B right one bit
ALR	Shift A left one bit, clear sign, zero to least significant bit
BLR	Shift B left one bit, clear sign, zero to least significant bit
ERA	Rotate E and A right one bit
ERB	Rotate E and B right one bit
ELA	Rotate E and A left one bit
ELB	Rotate E and B left one bit
ALF	Rotate A left four bits
BLF	Rotate B left four bits
SLA	Skip next instruction if least significant bit in A is zero
SLB	Skip next instruction if least significant bit in B is zero

These instructions may be combined as follows:

$1abe1 \begin{bmatrix} ALS \\ ARS \\ RAL \\ RAR \\ ALR \\ ALF \\ ERA \\ ELA \end{bmatrix} [,CLE]$	[,SLA]	ALS ARS RAL RAR ALR ALF ERA ELA	omments
---	--------	---	---------

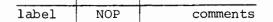


CLE, SLA, or SLB appearing alone or in any valid combination with each other are assumed to be a shift-rotate machine instruction.

The shift-rotate instructions must be given in the order shown. At least one and up to four are included in one statement. Instructions referring to the A-register may not be combined in the same statement with those referring to the B-register.

No-Operation Instruction

When a no-operation is encountered in a program, no action takes place; the computer goes on to the next instruction. A full memory cycle is used in executing a no-operation instruction.



A subroutine to be entered by a JSB instruction should have a NOP as the first statement. The return address can be stored in the location occupied by the NOP during execution of the program. A NOP statement causes the Assembler to generate a word of zeros.

Alter-Skip Group

The alter-skip group contains 19 basic instructions that can be combined to produce more than 700 different single cycle operations.

CLA	Clear	the	A-Register
CLB	Clear	the	B-Register

CMA	Complement the A-Register
CMB	Complement the B-Register
CCA	Clear, then complement the A-Register (set to ones)
CCB	Clear, then complement the B-Register (set to ones)
CLE	Clear the E-Register
CME	Complement the E-Register
CCE	Clear, then complement the E-Register
SEZ	Skip next instruction if E is zero
SSA	Skip if sign of A is positive (0)
SSB	Skip if sign of B is positive (0)
INA	Increment A by one
INB	Increment B by one
SZA	Skip if contents of A equals zero
SZB	Skip if contents of B equals zero
SLA	Skip if least significant bit of A is zero
SLB	Skip if least significant bit of B is zero
RSS	Reverse the sense of the skip instructions. If no skip instructions precede in the statement, skip the next in- struction

These instructions may be combined as follows:

			$\left[, \begin{cases} CLE \\ CME \\ CCE \end{cases}\right] [,SSA] [,SLA] [,INA] [,SZA] [,RSS]$	comments
label	$\left[\begin{cases} CLB \\ CMB \\ CCB \end{cases} \right]$	[,SEZ]	$\left[, \left\{ \begin{array}{c} CLE\\ CME\\ CCE \end{array} \right\} \right] [,SSB] [,SLB] [,INB] [,SZB] [,RSS]$	comments

The alter-skip instructions must be given in the order shown. At least one and up to eight are included in one statement. Instructions referring to the A-register may not be combined in the same statement with those referring to the B-register. When two or more skip functions are combined in a single operation, a skip occurs if any one of the conditions exists. If a word with RSS also includes both SSA and SLA (or SSB and SLB), a skip occurs only when sign and least significant bit are both set (1).

INPUT/OUTPUT, OVERFLOW, AND HALT

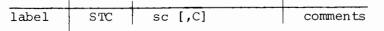
The input/output instructions allow the user to transfer data to and from an external device via a buffer, to enable or disable external interrupt, or to check the status of I/O devices and operations. A subset of these instructions permits checking for an arithmetic overflow condition.

Input/output instructions require the designation of a select code, sc, which indicates one of 64 input/output channels or functions. Each channel consists of a connect/disconnect control bit, a flag bit, and a buffer of up to 16 bits. The setting of the control bit indicates that a device associated with the channel is operable. The flag bit is set automatically when transmission between the device and the buffer is completed. Instructions are also available to test or clear the flag bit for the particular channel. If the interrupt system is enabled, setting of the flag causes program interrupt to occur; control transfers to the interrupt location related to the channel.

Expressions used to represent select codes (channel numbers) must have a value of less than 2^6 . The value specifies the device or operation referenced. Instructions which transfer data between the A or B register and a buffer, access the Switch register when sc = 1. The character C appended to such an instruction clears the overflow bit after the transfer from the switch register is complete.

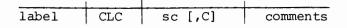
Input/Output

Prior to any input/output data transmission, the control bit is set. The instruction which enables the device may also transfer data between the device and the buffer.

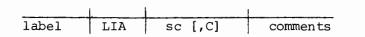


Set I/O control bit for channel specified by sc. STC transfers or enables transfer of an element of data from an input device to the buffer or to an

output device from the buffer. The exact function of the STC depends on the device; for the 2752A Teleprinter, an STC enables transfer or a series of bits. If sc = 1, this statement is treated as NOP. The C option clears the flag bit for the channel.

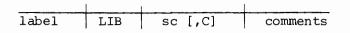


Clear I/O control bit for channel specified by sc. When the control bit is cleared, interrupt on the channel is disabled, although the flag may still be set by the device. If sc = 0, control bits for all channels are cleared to zero; all devices are disconnected. If sc = 1, this statement is treated as NOP.

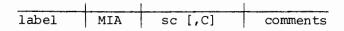




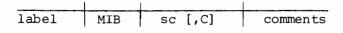
Load into A the contents of the I/O buffer indicated by sc.



Load into B the contents of the I/O buffer indicated by sc.



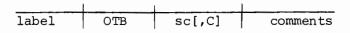
Merge (inclusive "or") the contents of the I/O buffer indicated by sc into A.



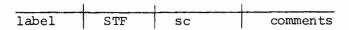
Merge (inclusive "or") the contents of the I/O buffer indicated by sc into B.



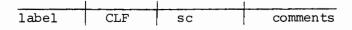
Output the contents of A to the I/O buffer indicated by sc.



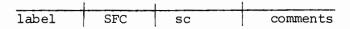
Output the contents of B to the I/O buffer indicated by sc.



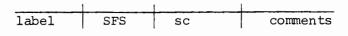
Sets the flag bit of the channel indicated by sc. If sc = 0, STF enables the interrupt system. A sc code of 1 causes the overflow bit to be set.



Clear the flag bit to zero for the channel indicated by sc. If sc = 0, CLF disables the interrupt system. If sc = 1, the overflow bit is cleared to zero.



Skip the next instruction if the flag bit for channel sc is clear. If sc = 1, the overflow bit is tested.



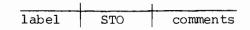
Skip the next instruction if the flag bit for channel sc is set. If sc = 1, the overflow is tested.

Overflow

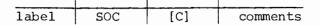
In addition to the use of a select code of 1, the overflow bit may be accessed by the following instructions:



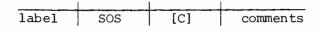
Clear the overflow bit.



Set overflow bit.



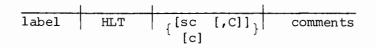
Skip the next instruction if the overflow bit is clear. The C option clears the bit after the test is performed.



Skip the next instruction if the overflow bit is set. The C option clears the bit after the test is performed.

The C option is identified by the sequence "space C space" following either "SOC" or "SOS". Anything else is treated as a comment.

Halt



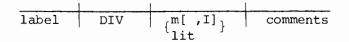
Halt the computer. The machine instruction word is displayed in the T-Register. If the C option is used, the flag bit associated with channel sc is cleared.

If neither the select code nor the C option is used, the comments portion must be omitted.

EXTENDED ARITHMETIC UNIT

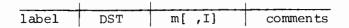
Ten instructions may be used with the EAU version of the Assembler or Extended Assembler to increase the computer's overall efficiency. The computer must include the Extended Arithmetic Unit option to obtain the resulting increase in available core storage and decrease in program run time.

The MPY instruction multiplies the contents of the A-Register by the contents of m. The product is stored in registers B and A. B contains the sign of the product and the 15 most significant bits; A contains the least significant bits.



The DIV instruction divides the contents of registers B and A by the contents of m. The quotient is stored in A and the remainder in B. Initially B contains the sign and the 15 most significant bits of the dividend; A contains the least significant bits.

The DLD instruction loads the contents of locations m and m + 1 into registers A and B, respectively.



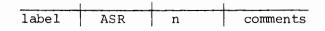
The DST instruction stores the contents of registers A and B in locations m and m + 1, respectively.

MPY, DIV, DLD, DST results in two machine words: a word for the instruction code and one for the operand.

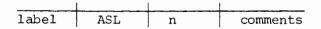
The above four instructions are available without the Extended Arithmetic Unit option as software subroutines.[†] As a part of the Extended Arithmetic option, they require less core storage and can be executed in less time.

† See ARITHMETIC SUBROUTINE CALLS, Section 4.

The following seven instructions can be used only on machines with the Extended Arithmetic Unit. These shift-rotate instructions provide the capability to shift or rotate the B- and A-Registers n number of bit positions, where 1 < n < 16.



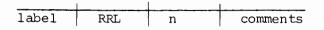
The ASR instruction arithmetically shifts the B- and A-Registers right n bits. The sign bit (bit 15 of B) is extended.



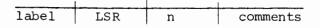
The ASL instruction arithmetically shifts the B- and A-Register left n bits. Zeroes are placed in the least significant bits. The sign bit (bit 15 of B) is unaltered. The overflow bit is set if bit 14 differs from bit 15 before each shift; otherwise, exit with overflow bit cleared.

 label
 RRR
 n
 comments

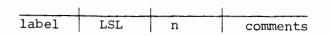
The RRR instruction rotates the B- and A-Registers right n bits.



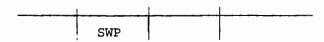
The RRL instruction rotates the B- and A-Registers left n bits.



The LSR instruction logically shifts the B- and A-Registers right n bits. Zeroes are placed in the most significant bits.



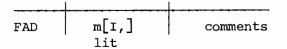
The LSL instruction logically shifts the B- and A-Registers left n bits. Place zeroes into the least significant bits.



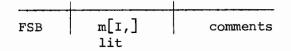
Exchange the contents of the A- and B-Registers. The contents of the A-Register are shifted into the B-Register and the contents of the B-Register are shifted into the A-Register.

FLOATING-POINT INSTRUCTIONS

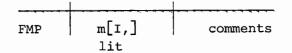
Floating-point instructions provide a means of performing calculations on floating-point values. Computers with the hardware floating-point option should use assemblers and libraries with floating-point capabilities. The floating-point assembler generates calls to the appropriate hardware function instead of the library subroutines. If the computer does not have the hardware floating-point option, then non-floating-point assemblers and libraries should be used.



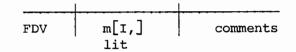
FAD performs an addition between a floating-point number stored in the A- and B-registers and a floating-point number stored in memory locations m and m + 1 The result is returned in the A- and B-registers.



The FSB instruction subtracts a floating-point value in memory locations m + 1 from a floating-point value in the A- and B-registers. The result is returned in the A- and B-registers.



The FMP instruction multiplies a floating-point value in memory locations m and m + 1 with a floating-point value in the A- and B-registers. The result is returned in the A- and B-registers.



The FDV instruction divides the floating-point value in memory locations m and m + 1 into the value stored in the A- and B-registers. The result is returned in the A- and B-registers.



The FIX instruction converts a floating-point number contained in the A- and B-registers to a fixed point number. The result is returned in the A-register. The contents of the B-register are meaningless.



The FLT instruction converts a fixed-point value contained in the A-register to a floating-point value. The result is returned in the A- and B-registers.

SECTION IV PSEUDO INSTRUCTIONS

The pseudo instructions control the Assembler, establish program relocatability, and define program linkage as well as specify various types of constants, blocks of memory, and labels used in the program. With the Extended Assembler, pseudo instructions also control listing output.

ASSEMBLER CONTROL

The Assembler control pseudo instructions establish and alter the contents of the base page and program location counters, and terminate assembly processing. Labels may be used but they are ignored by the Assembler. NAM records produced by the Assemblers are accepted by the DOS, DOS-M and BCS loaders.

NAM defines the name of a relocatable program. A relocatable program must begin with a NAM statement.⁺ A relocatable program is assembled assuming a starting location of zero (i.e., zero relative). The name may be a symbol of one to five alphanumeric characters the first of which must be alphabetic or a period. The program name is printed on the list output. The name is optional and if omitted, the comments must be omitted also.

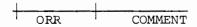


The ORG statement defines the origin of an absolute program, or the origin of subsequent sections of absolute or relocatable programs.

The Control Statement, the HED instruction, and comments may appear prior to the NAM or ORG statements. If the Control Statement (ASMB,...) does not appear on tape preceding the program, it must be entered from the teleprinter.

An absolute program must begin with an ORG statement.⁺ The operand m, must be a decimal or octal integer specifying the initial setting of the program location counter.

ORG statements may be used elsewhere in the program to define starting addresses for portions of the object code. For absolute programs the Operand field, m, may be any expression. For relocatable programs, m, must be a program relocatable expression; it may not be base page or common relocatable or absolute. An expression is evaluated modulo 2¹⁵. Symbols must be previously defined. All instructions following an ORG are assembled at consecutive addresses starting with the value of the operand.



ORR resets the program location counter to the value existing when an ORG or ORB instruction was encountered.

Example:

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+	\neg	-	_	-	-	⊢	ł	╀	+	+-	+	+	+	-		\square			-	-	┝─	\vdash	+	+	+	+	+	+	+-	+	╀	+	+	+-	+	⊢	\vdash	+	+	+-	┞	+	+	+-	+	╀	+	+	Н	\vdash	\vdash	

⁺ The Control Statement, the HED instruction, and comments may appear prior to the NAM or ORG statements. If the Control Statement (ASMB,...) does not appear on tape preceding the program, it must be entered from the teleprinter.

More than one ORG or ORB statement may occur before an ORR is used. If so, when the ORR is encountered, the program location counter is reset to the value it contained when the first ORG or ORB of the string occurred.

Example:

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If a second ORR appears before an intervening ORG or ORB, the second ORR is ignored.

ORR cannot be used to reset the location counter for locations in the base page that are governed by the ORB statement.



ORB defines the portion of a relocatable program that must be assigned to the base page by the Assembler. The Label field (if given) is ignored, and the statement requires no operand. All statements that follow the ORB statement are assigned contiguous locations in the base page. Assignment to the base page terminates when the Assembler detects an ORG, ORR, or END statement.

When more than one ORB is used in a program, each ORB causes the Assembler to resume assigning base page locations at the address following the last assigned base page location.

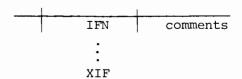
An ORB statement in an absolute program has no significance and is flagged as an error.

Example:

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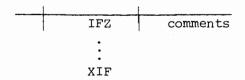
The IFN and IFZ pseudo instructions cause the inclusion of instructions in a program provided that either an "N" or "Z", respectively, is specified as a parameter for the ASMB control statement.[†] The IFN or IFZ instruction precedes the set of statements that are to be included. The pseudo instruction XIF serves as a terminator. If XIF is omitted, END acts as a terminator to both the set of statements and the assembly. IFN and IFZ may be used only when the source program is translated by the Extended Assembler which is provided for 8K or larger machines.

†See CONTROL STATEMENT, Section 5.



All source language statements appearing between the IFN and the XIF pseudo instructions are included in the program if the character "N" is specified on the ASMB control statement.

All source language statements appearing between the IFZ and the XIF pseudo instructions are included in the program if the character "Z" is specified on the ASMB control statement.



When the particular letter is not included on the control statement, the related set of statements appears on the Assembler output listing but is not assembled.

Any number of IFN-XIF and IFZ-XIF sets may appear in a program, however, they may not overlap. An IFZ or IFN intervening between an IFZ or IFN and the XIF terminator results in a diagnostic being issued during compilation; the second pseudo instruction is ignored.

Both IFN-XIF and IFZ-XIF pseudo instructions may be used in the program; however, only one type will be selected in a single assembly. Therefore, if both characters "N" and "Z" appear in the control statement, the character which is listed last will determine the set of coding that is to be included in the program.

Example:

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	IFZ LDAC	AR							
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		TIFS			┥┨┼┾┽				
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┝╋┽┾┼┫	XIF					+ + + + +-			
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┝╋╪┼┼┨					++++++				
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	I F N L D A P				╌┨┤┤╌┾	┿╋┼┾┿╸	┝╋╋┥		┝┼╊┽╉
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┝╋╌┝╌┝╌┝	CMA,S JMP N	L A N E Z A O . G O	┥┟┼┼╴		┼┠┼┼┽	┼┠┼┼╌			┝┼╋┼╋
┝╋┿╋┼┙					┼╋┽┼┼	┼╂┽┼┿━		┥╉┼┥	┝┼╋┼╋
┝╍┨╶┥╌┥	LDA T CPA C	I M E O S T			┿╉┼┼┾	┿╉┈┊┊┾╸	╞╉┊┿┿┈	+ +	┝┿╌┠┼╶╂
┝╋┽┼┊┥		051						┝╺╂╴╎╴╎╍╸	┝┼╂┼╈
	XIF HLT 7	7					+ +	$ \cdot \cdot \cdot $	+++++
NO.GO	┝╼╋╼╪╌┽┈╋╼╋╼╡	1						┝┨╌┿┥╋	┝┼╋┿╋
┝╋┽┿┿┿	•	┝┼┅┿╋╋╋┉┿╼	┝┿╋┥		┿╋┿┿	+	┥╋┽┥┿	+	┝┼╂┼╇
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		┝┼╂┼┠┼┼				++++			
++++	END	┝┼┼┼┼┼			++++	┼╆╬┼┾╸			

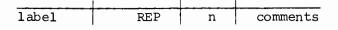
Program TRAVL will perform computations involving either or neither CAR or PLANE considerations depending on the presence or absence of Z or N parameters in the Control Statement.

Example:

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Program WAGES computes a weekly wage value. Overtime consideration will be included in the program if "Z" is included in the parameters of the Control Statement.

The REP pseudo instruction, available in the Extended Assembler only, causes the repetition of the statement immediately following it a specified number of times.



The statement following the REP in the source program is repeated n times. The n may be any absolute expression. Comment lines (indicated by an asterisk in character position 1) are not repeated by REP. If a comment follows a REP instruction, the comment is ignored and the instruction following the comment is repeated.

A label specified in the REP pseudo instruction is assigned to the first repetition of the statement. A label cannot be part of the instruction to be repeated; it would result in a doubly defined symbol error.

Example:

CLA TRIPL REP 3 ADA DATA

The above source code would generate the following:

	CLA		Clear the A-Register; the content of DATA is
TRIPL	ADA ADA ADA	DATA DATA DATA	tripled and stored in the A-Register.

Example:

FILL	REP	100B
	NOP	

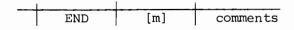
The example above loads 100_8 memory locations with the NOP instruction. The first location is labeled FILL.

Example:

REP 2 MPY DATA

The above source code would generate the following:

MPY DATA MPY DATA



This statement terminates the program; it marks the physical end of the source language statements. The Operand field, m, may contain a name appearing as a statement label in the current program or it may be blank. If a name is entered, it identifies the location to which the loader transfers control after a relocatable program is loaded. A NOP should be stored at that location; the loader transfers control via a JSB.

If the Operand field is blank, the Comments field must be blank also, otherwise, the Assembler attempts to interpret the first five characters of the comments as the transfer address symbol.

The Label field of the END statement is ignored.

OBJECT PROGRAM LINKAGE

Linking pseudo instructions provides a means for communication between a main program and its subroutines or among several subprograms that are to be run as a single program. These instructions may be used only in a relocatable program.

The Label field of this class is ignored in all cases. The Operand field is usually divided into many subfields, separated by commas. The first space not preceded by a comma or a left parenthesis terminates the entire field.

COM reserves a block of storage locations that may be used in common by several subprograms. Each name identifies a segment of the block for the subprogram in which the COM statement appears. The sizes are the number of words allotted to the related segments. The size is specified as an octal or decimal integer. If the size is omitted, it is assumed to be one.

Any number of COM statements may appear in a subprogram. Storage locations are assigned contiguously; the length of the block is equal to the sum of the lengths of all segments named in all COM statements in the subprogram.

To refer to the common block, other subprograms must also include a COM statement. The segment names and sizes may be the same or they may differ. Regardless of the names and sizes specified in the separate subprograms, there is only one common block for the combined set. It has the same relative origin; the content of the nth word of common storage is the same for all subprograms.

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Organization of common block:

PROG1	PROG2	Common	
name	name	Block	
ADDR1	AAA	(location	1)
		(location	2)
	AAB	(location	3)
		(location	4)
	AAC	(location	5)
ADDR2	AAD	(location	6)
		(location	7)
		(location	8)
		(location	9)
		(location	10)
		(location	11)
		(location	12)
		(location	13)
		(location	14)
		(location	15)
ADDR3		(location	16)
		(location	17)
		(location	18)
		(location	19)
		(location	20)
		(location	21)
		(location	22)
		(location	23)
		(location	24)
		(location	25)

The LDA instructions in the two subprograms each refer to the same location in common storage, location 7.

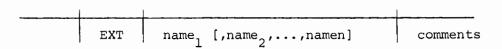
The segment names that appear in the COM statements can be used in the Operand fields of DEF, ABS, EQU, or any memory reference statement; they may not be used as labels elsewhere in the program.

The loader establishes the origin of the common block; the origin cannot be set by the ORG or ORB pseudo instruction. All references to the common area are relocatable.

Two or more subprograms may declare common blocks that differ in size. The subprogram that defines the largest block must be the first submitted for loading.

ENT defines entry points to the program or subprogram. Each name is a symbol that is assigned as a label for some machine operation in the program. Entry points allow another subprogram to refer to this subprogram. All entry points must be defined in the program.

Symbols appearing in an ENT statement may not also appear in EXT or COM statements in the same subprogram.



This instruction designates labels in other subprograms that are referenced in this subprogram. The symbols must be defined as entry points by the other subprograms.

The symbols defined in the EXT statement may appear in memory reference statements, the EQU or DEF pseudo instructions. An external symbol must appear

alone; it may not be in a multiple term expression or be specified as indirect. References to external locations are processed by the BCS loader as indirect addresses linked through the base page.

Symbols appearing in EXT statements may not also appear in ENT or COM statements in the same subprogram. The label field is ignored.

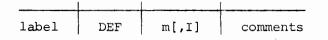
Example:

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PF	2	0	G	Å		N	0	Ρ	Ĩ				Γ	Γ	Γ	Γ			T	Г	Γ	Γ			T						Γ				Π															Τ	Τ	T
Ŧ	+	-	Ť		Н	L	D	Δ		S	Δ	M	D		t	t	t	t	t	s	Δ	N	١ſ)	t	4	N	D		S	Δ	N	D		A R	R	F		R	E	F	E	R	Ε	N	С	Ε	D		r	N	1
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ADDRESS AND SYMBOL DEFINITION

The pseudo operations in this group assign a value or a word location to a symbol which is used as an operand elsewhere in the program.





The address definition statement generates one word of memory as a 15-bit address which may be used as the object of an indirect address found elsewhere in the source program. The symbol appearing in the label is that which is referenced; it appears in the Operand field of a Memory Reference instruction.

The operand field of the DEF statement may be any positive expression in an absolute program; in a relocatable program it may be a relocatable expression or an absolute expression with a value of less than 100₈. Symbols that do appear in the Operand field may appear as operands of EXT or COM statements, in the same subprogram and as entry points in other subprograms.

The expression in the Operand field may itself be indirect and make reference to another DEF statement elsewhere in the source program.

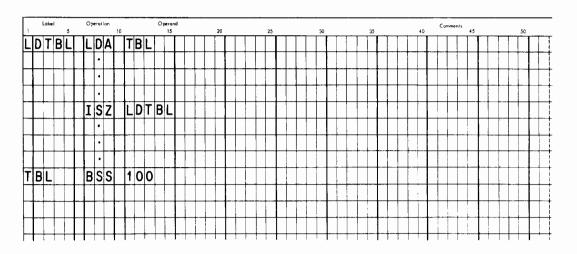
Example:

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The DEF statement provides the necessary flexibility to perform address arithmetic in programs which are to be assembled in relocatable form. Relocatable programs should not modify the operand of a memory reference instruction.

In the example below, if TBL and LDTBL are in different pages, the Loader processes TBL as an indirect address linked through the base page. The ISZ erroneously increments the loader-provided reference to the base page rather than the value of TBL.

Example:



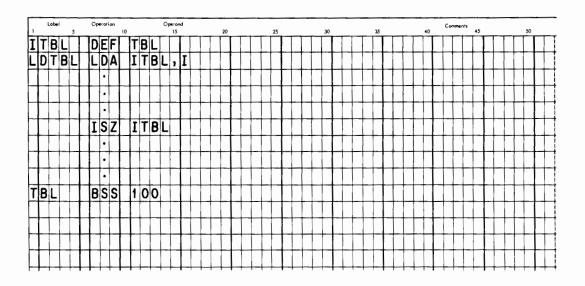
Assuming the loader might assign absolute locations comparable to the following octal values:

Page	Loc	Opcode	Reference
(0)	(700)	DEF	400
(1)	(200)	LDA :	(O) 700(I)
(1)	(300)	ISZ :	(1) 200
(2)	(0)		(TBL)

It can be seen that the ISZ instruction would increment the quantity 700 rather than the address of the table (4000_8) .

The following assures correct address modification during program execution.

Example:



This sequence might be stored by the loader as:

Page	Loc	Opcode	Reference
(1)	(200)	DEF	4000
(1)	(201)		200(1)
(1)	(300)	ISZ :	(1) (200)
(2)	(0)		(TBL)

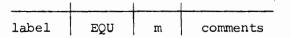
The value of 4000 is incremented; each execution of LDA will access sucessive locations in the table.



ABS defines a 16-bit absolute value to be stored at the location represented by the label. The Operand field, m, may be any absolute expression; a single symbol must be defined as absolute elsewhere in the program.

Example:

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The EQU pseudo operation assigns to a symbol a value other than the one normally assigned by the program location counter. The symbol in the Label field is assigned the value represented by the Operand field. The Operand field may contain any expression. The value of the operand may be common, base page or program relocatable as well as absolute, but it may not be negative. Symbols appearing in the operand must be previously defined in the source program.

The EQU instruction may be used to symbolically equate two locations in memory, or it may be used to give a value to a symbol. The EQU statement does not result in a machine instruction.

Examples:

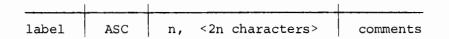
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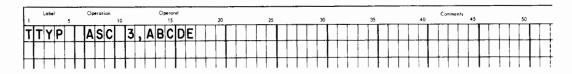
CONSTANT DEFINITION

The pseudo instructions in this class enter a string of one or more constant values into consecutive words of the object program. The statements may be named by labels so that other program statements can refer to the fields generated by them.



ASC generates a string of 2n alphanumeric characters in ASCII code into n consecutive words.[†] One character is right justified in each eight bits; the most significant bit is sero. n may be any expression resulting in an unsigned decimal value in the range 1 through 28. Symbols used in an expression must be previously defined. Anything in the Operand field following 2n characters is treated as comments. If less than 2n characters are detected before the end-of-statement mark, the remaining characters are assumed to be spaces, and are stored as such. The label represents the address of the first two characters.

Example:



causes the following:

	А	LPHABETIC		
<u>15 14</u>		876		0
TTYP	Α	VA	В	
	С	VA	D	
	E		٨	

EQUIVALENT IN OCTAL NOTATION 15 14 8 7 6 0 TTYP 1 0 1 1 0 2 1 0 3 1 0 4

† To enter the code for the ASCII symbols which perform some action (e.g., CR and LF), the OCT pseudo instruction must be used.

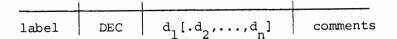
0

4

0

5

0

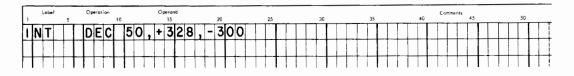


DEC records a string of decimal constants into consecutive words. The constants may be either integer or real (floating point), and positive or negative. If no sign is specified, positive is assumed. The decimal number is converted to its binary equivalent by the Assembler. The label, if given, serves as the address of the first word occupied by the constant.

A decimal integer must be in the range of 0 to 2^{15} -1; it may assume positive, negative, or zero values. It is converted into one binary word and appears as follows:



Example:



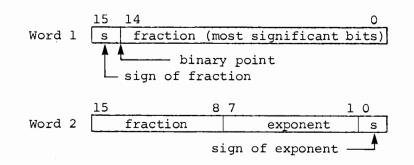
causes the following (octal representation)

	15	14				0
INT	ο	0	0	0	6	2
	0	0	0	5	1	0
	1	7	7	3	2	4

A floating point number has two components, a fraction and an exponent. The exponent specifies the power of 10 by which the fraction is multiplied. The fraction is a signed or unsigned number which may be written with or without a decimal point. The exponent is indicated by the letter E and follows a signed or unsigned decimal integer. The floating point number may have any of the following formats:

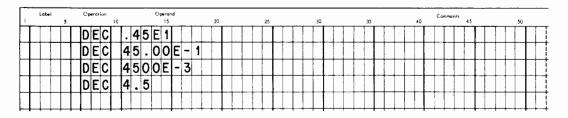
+n.n +n. +n.nE+e +.nE+e +n.E+e +nE+e

The number is converted to binary, normalized (leading bits differ), and stored in two computer words. If either the fraction or the exponent is negative, that part is stored in two's complement form.



The floating point number is made up of a 7-bit exponent with sign and a 23-bit fraction with sign. The number must be in the approximate range of 10^{-38} and zero.

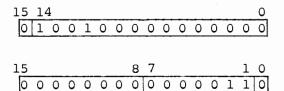
Examples:

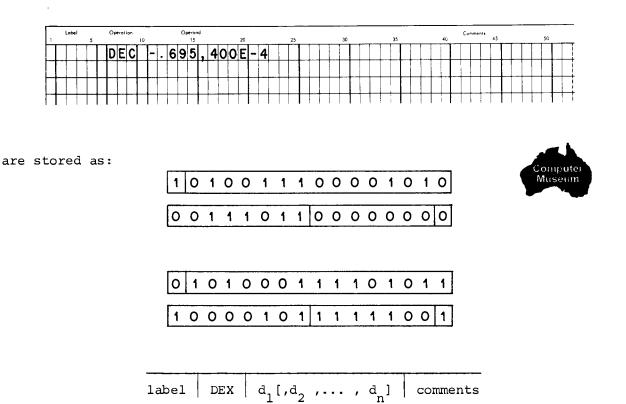


are all equivalent to

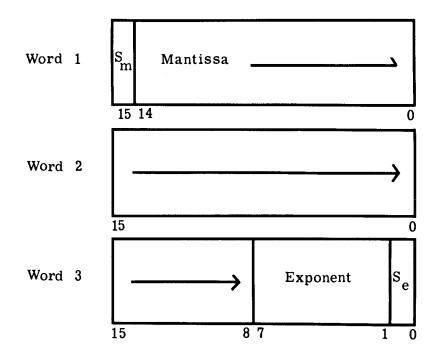
.45x10¹

and are stored in normalized form as:





DEX, for the Extended Assembler, records a string of extended precision decimal constants into consecutive words within a program. Each such extended precision constant occupies three words as shown below:



Legend: S = Sign of the mantissa (fraction)

S = Sign of the Exponent

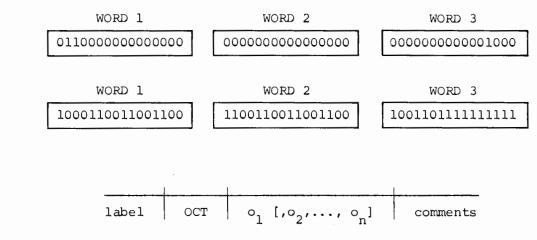
NOTE: A value is entered only if normalizing of the mantissa is needed.

An extended precision floating point number is made up of a 39-bit mantissa (fraction) and sign and a 7-bit exponent and sign. The exponent and sign will be zero if the mantissa does not have to be normalized.

This is the only form used for DEX. All values, whether they be floating point, integer, fraction, or integer and fraction, will be stored in three words as just described. This storage format is basically an extension of that used for DEC, as previously described:

Examples:

are stored as:



OCT stores one or more octal constants in consecutive words of the object program. Each constant consists of one to six octal digits (0 to 17777). If no sign is given, the sign is assumed to be positive. If the sign is negative, the two's complement of the binary equivalent is stored. The constants are separated by commas; the last constant is terminated by a space. If less than six digits are indicated for a constant, the data is right justified in the word. A label, if used, acts as the address of the first

constant in the string. The letter B must not be used after the constant in the Operand field; it is significant only when defining an octal term in an instruction other than OCT.

Examples:

1		Lab	e!		Op	trat	ion	10					Ope	rano 15					20					25					30					35					40		C۰	time:	n1s	45					50		
T			Π	Ń	0	С	T	Ē	+	Q	5	Ι	T	Ī	T	Т	T	Ť	Ĩ					Ĺ	Γ			-	Ĩ	Т	T			32					•				Γ	Γ	Г	T		1	30	Г	Γ
1		T	Π		-	Ċ	T	t	-	1	+-	T	1	1	1	T	T	1	1					F	ſ	t		-	1	1	1											F	F	F	t	1	t		F	F	t
1	U	M	1		0	С	T	Γ	1	7	7	7	,	2	ø	4 (Ø	5	,	-	3	6			T				T															T	t	T	1	T		T	t
T					0	С	Т	Γ	5	1	I,	ŀ	7	7	7	7	7	,	-	1	,	1	ø	1	Ø	1			T															ſ	T	t	T	T		Γ	Ī
T		Γ			0	С	Т	Γ	1	Q	57	1	64	4	2	, '	1	7	7	Ø	7	7							T															ſ	T	T	1		Γ	F	t
T		T			0	С	Т		1	Ś)7	7	6	T										1		Γ	Π			I	L	L	E	G	A	L	;		C	0	N	Т	A	Ι	N	S	5	Τ		Γ	t
I					0	С	Т		-	1	17	7	7	7	7	7		1	1						Γ				T	D	I	G	I	Т		9					-			T	t	T	Τ	T	Γ		t
Ι					0	С	T	ſ	1	7	7	1	B	Τ		T	T		Ţ				Γ							I	L	L	Ε	G	A	L	:		C	0	N	Т	A	Ι	N	S	;	1			Ī
Τ								Γ	Τ			Τ	Τ	Τ		Τ	Τ	Τ	Τ						Γ					C	H	A	R	A	C	Т	E	R		B					Γ	Γ		Γ			Ī
1								Γ	Τ				1	1		T	T		1						Γ				T	1															Γ	Γ					Ī
t							-	T			1	1	1	1	+	-+	+	+	1	- 1			\square	1	t	1		- †	╉	1	+	-1				t	\square	\square	-	-	1			\uparrow	t	+	1	1		1	t

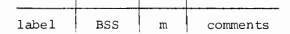
The previous statements are stored as follows:

	1514					0	
	0	0	0	0	0	0	
	1	7	7	7	7	6	
NUM	0	0	0	1	7	7	
	0	2	0	4	0	5	
	1	7	7	7	4	2	
	0	0	0	0	5	1	
	0	7	7	7	7	7	
	1	7	7	7	7	7	
	0	1	0	1	0	1	
	1	0	7	6	4	2	
	1	7	7	0	7	7	
	Х	X	X	х	x	X	ATTEM
	0	0	0	0	0	1	
	х	X	x	x	x	×	CONST PREDIC

THE RESULT OF ATTEMPTING TO DEFINE AN ILLEGAL CONSTANT IS UN-PREDICTABLE

STORAGE ALLOCATION

The storage allocation statement reserves a block of memory for data or for a work area.



The BSS pseudo operation advances the program or base page location counter according to the value of the operand. The Operand field may contain any expression that results in a positive integer. Symbols, if used, must be previously defined in the program. The label, if given, is the name assigned to the storage area and represents the address of the first word. The initial content of the area set aside by the statement is unaltered by the loader.

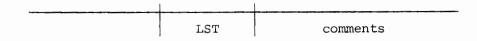
ASSEMBLY LISTING CONTROL

Assembly listing control pseudo instructions allow the user to control the assembly listing Output during pass 2 or 3 of the assembly process. These pseudo instructions may be used only when the source program is translated by the Extended Assembler provided for 8K or larger machines (8,192-word memory or larger).



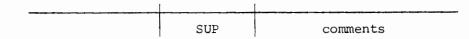
Output is suppressed from the assembly listing, beginning with the UNL pseudo instruction and continuing for all instructions and comments until either an LST or END pseudo instruction is encountered. Diagnostic messages for errors encountered by the Assembler will be printed, however. The source statement sequence numbers (printed in columns 1-4 of the source program listing) are incremented for the instructions skipped.

4-24



The LST pseudo instruction causes the source program listing, terminated by a UNL, to be resumed.

A UNL following a UNL, a LST following a LST, and a LST not preceded by a UNL are not considered errors by the Assembler.



The SUP pseudo instruction suppresses the output of additional code lines from the source program listing. Certain pseudo instructions, because they result in using subroutines, generate more than one line of coding. These additional code lines are suppressed by a SUP instruction until a UNS or the END pseudo instruction is encountered. SUP will suppress additional code lines in the following pseudo instructions:

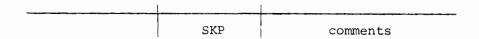
ASC	DIV	FAD	FSB
OCT	DLD	FDV	MPY
DEC	DST	FMP	

The SUP pseudo instruction may also be used to suppress the listing of literals at the end of the source program listing.

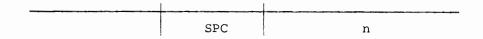


The UNS pseudo instruction causes the printing of additional coding lines, terminated by a SUP, to be resumed.

A SUP preceded by another SUP, UNS preceded by UNS, or UNS not preceded by a SUP are not considered errors by the Assembler.



The SKP pseudo instruction causes the source program listing to be skipped to the top of the next page. The SKP instruction is not listed, but the source statement sequence number is incremented for the SKP.



The SPC pseudo instruction causes the source program listing to be skipped a specified number of lines. The list output is skipped n lines, or to the bottom of the page, whichever occurs first. The n may be any absolute expression. The SPC instruction is not listed but the source statement sequence number is incremented for the SPC.



The HED pseudo instruction allows the programmer to specify a heading to be printed at the top of each page of the source program listing.

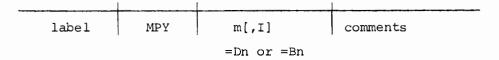
The heading, m, a string of up to 56 ASCII characters, is printed at the top of each page of the source program listing following the occurrence of the HED pseudo instruction. If HED is encountered before the NAM or ORG at the beginning of a program, the heading will be used on the first page of the source program listing. A HED instruction placed elsewhere in the program causes a skip to the top of the next page.

The heading specified in the HED pseudo instruction will be used on every page until it is changed by a succeeding HED instruction.

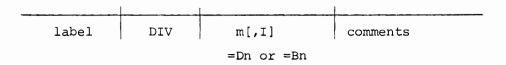
The source statement containing the HED will not be listed, but source statement sequence number will be incremented.

ARITHMETIC SUBROUTINE CALLS

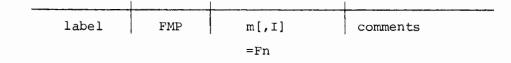
The members of this group of pseudo instructions request the Assembler to gener te calls to arithmetic subroutines* external to the source program. These pseudo instructions may be used in relocatable programs only. The Operand field may contain any relocatable expression or an absolute expression resulting in a value of less than 100_{o} .



Multiply the contents of the A-register by the contents of m or the quantity defined by the literal and store the product in registers B and A. B contains the sign of the product and the 15 most significant bits; A contains the least significant bits.

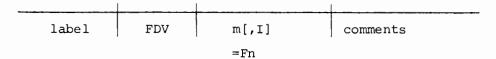


Divide the contents of registers B and A by the contents of m or the quanity defined by the literal. Store the quotient in A and the remainder in B. Initially B contains the sign and the 15 most significant bits of the dividend; A contains the least significant bits.

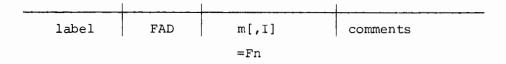


^{*} Not intended for use with DEX formatted numbers. For such numbers JSB's to double precision subroutines must be used. See RELOCATABLE SUBROUTINES Manual.

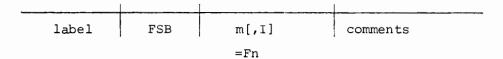
Multiply the two-word floating point quanity in registers A and B by the twoword floating point quantity in locations m and m+l or the quantity defined by the literal. Store the two-word floating point product in registers A and B.



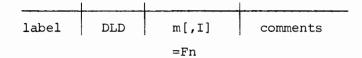
Divide the two-word floating point quantity in registers A and B by the twoword floating point quantity in locations m and m+l or the quantity defined by the literal. Store the two-word floating point quotient in A and B.



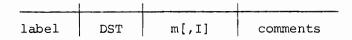
Add the two-word floating point quantity in registers A and B to the twoword floating point quantity in locations m and m+l or the quantity defined by the literal. Store the two-word floating point sum in A and B.



Subtract the two-word floating point quantity in m and m+l or the quantity defined by the literal from the two-word floating point quantity in registers A and B and store the difference in A and B.



Load the contents of locations m and m+l or the quantity defined by the literal into registers A and B respectively.



Store the contents of registers A and B in locations m and m+l respectively.

Each use of a statement from this group generates two words of instructions. Symbolically, they could be represented as follows:

JSB <.arithmetic pseudo operation> DEF m [,I]

An EXT <.arithmetic pseudo operation> is implied preceding the JSB operation.

In the above operations, the overflow bit is set when one of the following conditions occurs:

Integer overflow Floating point overflow or underflow Division by zero.

Execution of any of the subroutines alters the contents of the E-Register.

SECTION V ASSEMBLER INPUT AND OUTPUT

The Assembler accepts as input a paper tape containing a control statement and a source language program. A relocatable source language program may be divided into several subroutines; the designation of these elements is optional. The output produced by the Assembler may include a punched paper tape containing the object program, an object program listing, and diagnostic messages.

CONTROL STATEMENT

The control statement specifies the output to be produced:

 $ASMB, p_1, p_2, \dots, p_n$

"ASMB," is entered in positions 1-5. Following the comma are one or more parameters, in any order, which define the output to be produced. The control statement must be terminated by an end-of-statement mark, CR LF.

The parameters may be any legal combination of the following starting in position 6:

A Absolute: The addresses generated by the Assembler are to be interpreted as absolute locations in memory. The program is a complete entity. It may not include NAM, ORB, COM, ENT, EXT, arithmetic pseudo operation statements or literals. The binary output format is that specified for the Basic Binary loader.

- R Relocatable: The program may be located anywhere in memory. Instruction operands are adjusted as necessary. The binary output format is that specified for the Relocating loader.
- B Binary output: A program is to be punched according to one of the above parameters.
- L List output: A program listing is to be produced either during pass two or pass three (if binary output selected) according to one of the above parameters.
- T Table print: List the symbol table at the end of the first pass. For the Extended Assembler: List the symbol table in alphabetic order in three sections: section 1 for onecharacter symbols, section 2 for two- and three-character symbols, and section 3 for four- and five-character symbols.
- N Include sets of instructions following the IFN pseudo instruction.
- Z Include sets of instructions following the IFZ pseudo instruction.
- F Accepted by the Assembler to provide compatibility with DOS or DOS-M Assembler programs. F causes no action in any other assemblers.

Either A or R must be specified in addition to any combination of B, L, or T.

If a programmer wishes to assemble Pass 1 of a source program to check for errors, he can specify only an A or R to be the sole parameter of the Assembler Control Statement, executing only Pass 1. (This produces Pass 1 error messages without listing the program or providing an object tape). Extended Assembler only.

ASSEMBLER INPUT AND OUTPUT

The Assembler control statement must specifically request pass 2 operations (list or punch) in order for pass 2 to be executed. Lack of pass 2 option information causes processing only of pass 1 errors. If a C option is also provided, an automatic cross-reference symbol table is done after pass 1 when operating in the MTS environment.

The control statement may be on the same tape as the source program, or on a separate tape; or it may be entered via the teleprinter keyboard.



SOURCE PROGRAM

The first statement of the program (other than remarks or a HED statement) must be a NAM statement for a relocatable program or an ORG statement for indicating the origin of an absolute program. The last statement must be an END statement and may contain a transfer address for the start of a relocatable program. Each statement is followed by an end-of-statement mark.

BINARY OUTPUT

The punch output is defined by the ASMB control statement. The punch output includes the instructions translated from the source program. It does not include system subroutines referenced within the source program (arithmetic subroutine calls, .IOC., .DIO., .ENTR, etc.)

LIST OUTPUT

Fields of the object program are listed in the following print columns.

Columns	Content		
1-4	Source statement sequence number generated by the Assembler		
5-6	Blank		
7-11	Location (octal)		
12	Blank		
13-18	Object code word in octal		
19	Relocation or external symbol indicator		
20	Blank		
21-72	First 52 characters of source statement.		

Lines consisting entirely of comments (i.e., * in column 1) are printed as follows:

Columns	Content
1-4	Source statement sequence number
5-72	Up to 68 characters of comments

A Symbol Table listing has the following format:

Columns	Content
1-5	Symbol
6	Blank
7	Relocation of external symbol indicator
8	Blank
9-14	Value of the symbol

ASSEMBLER INPUT AND OUTPUT

The characters that designate an external symbol or type of relocation for the Operand field or the symbol are as follows:

Character	Relocation Base
Blank	Absolute
R	Program relocatable
В	Base page relocatable
С	Common relocatable
х	External symbol

At the end of each pass, the following is printed:

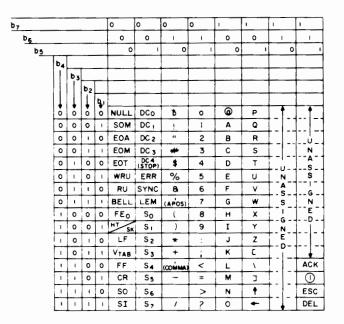
```
**NO ERRORS*
or
**nnnn ERRORS*
```

The value nnnn indicates the number of errors.

Note: For complete operating instructions for the HP Assembler or Extended Assembler, consult SOFTWARE OPERATING PROCEDURES, SIO SUBSYSTEMS module, (5951-1390).

APPENDIX A HP CHARACTER SET

ASCII CHARACTER FORMAT



Standard 7-bit set code positional order and notation are shown below with b_7 the high-order and b_1 the low-order, bit position.

Example: The code for "R" is: $\begin{bmatrix} b_7 & b_6 & b_5 & b_4 & b_3 & b_2 & b_1 \\ 1^7 & 0^6 & 1^5 & 0^4 & 0^3 & 1^2 & 0^1 \end{bmatrix}$

LEGEND

NULL	Null/Idle	DC1-DC3	Device Control
SOM	Start of message	DC ₄ (Stop)	Device control (stop)
EOA	End of address	4 1	
EOM	End of message	ERR	Error
EOT	End of transmission	SYNC	Synchronous idle
WRU	"Who are you?"	LEM	Logical end of media
RU	"Are you?"	s _o -s ₇	Separator (inform- ation
BELL FE	Audible signal Format effector	đ	Word separator (space, normally non-printing)
НТ	Horizontal tabulation	<	Less than
SK	Skip (punched card)	>	Greater than
\mathbf{LF}	Line feed	†	Up arrow (Exponentiation)
V TAB	Vertical tabulation	←	Left arrow (Implies/ Replaced by)
FF	Form feed	\ \	Reverse slant
CR	Crrriage return		
SO	Shift out	ACK	Acknowledge
SI	Shift in	(1)	Unassigned control
DC	Device control reserved for	ESC	Escape
0	data link escape	DEL	Delete/Idle

HP CHARACTER SET

BINARY CODED DECIMAL FORMAT

Symbol	BCD (octal code)	ASCII Equivalent (octal code)	Symbol	BCD (<u>octal code</u>)	ASCII Equivalent (octal code)
(Space)	2Ø	Ø4Ø	A	61	1Ø1
!	52	ø41	В	62	1ø2
#	13	Ø43	c	63	1ø2 1ø3
\$	53	ø44	D	64	1ø4
8	34	ø45	Ē	65	1ø5
&	6Ø	ø46	F	66	1ø6
•	14	ø47	G	67	1Ø7
(34	ø5ø	Н	7Ø	11ø
)	74	Ø51	I	71	111
*	54	Ø52	J	41	112
+	6Ø	Ø53	К	42	113
,	33	Ø54	L	43	114
-	4Ø	Ø55	М	44	115
•	73	Ø56	Ν	45	116
/	21	Ø57	0	46	117
			P	47	12Ø
ø	12	Ø6Ø	Q	5ø	121
1	Øl	Ø61	R	51	122
2	Ø2	Ø62	S	22	123
3	øз	Ø63	Т	23	124
4	Ø4	Ø64	U	24	125
5	Ø5	Ø 65	V	25	126
6	Ø6	Ø66	W	26	127
7	Ø7	Ø6 ⁻ 7	Х	27	13Ø
8	lø	Ø7Ø	Y	ЗØ	131
9	11	Ø71	Z	31	132
:	15	Ø72	[75	133
;	56	Ø73	Ν.	36	134
<	76	Ø74]	55	135
=	13	Ø75			
>	16	Ø76			
?	72	Ø77			
@	14	løø			

Kennedy 1406/1506 ASCII-BCD Conversion

Other symbols which may be represented in ASCII are converted to spaces in BCD (20)

	ASCII	BCD		ASCII	BCD
Symbol	(Octal code)	(Octal code)	Symbol	(Octal code)	(Octal code)
				<u></u>	<u>(()))))))))))))))))))))))))))))))))))</u>
(Space)	4ø	2Ø	А	lØl	61
1	41	52	В	1ø2	62
	42	37	С	lø3	63
#	43	13	D	1ø4	64
\$	44	53	E	1ø5	65
oto	45	34	F	1ø6	66
&	46	60 †	G	lø7	67
,	47	36	Н	llø	7ø
(5Ø	75	I	111	71
)	51	55	J	112	41
*	52	54	К	113	42
+	53	6Ø	L	114	43
,	54	33	М	115	44
-	55	4ø	N	116	45
•	56	73	0	117	46
/	57	21	Р	12Ø	47
			Q	121	5ø
			R	122	51
ø	6Ø	12	S	123	22
1	61	Øl	т	124	23
2	62	Ø2	U	125	24
3	63	øз	V	126	25
4	64	Ø4	W	127	26
5	65	Ø5	Х	13Ø	27
6	66	Ø6	Y	131	ЗØ
7	67	Ø7	Z	132	31
8	7ø	lØ			
9	71	11	[133	75 ‡
]	135	55 ‡
			↑	136	77
:	72	15	*	137	32
;	73	56			
<	74	76			
=	75	35			
>	76	16			
?	77	72			
@	lØØ	14			

HP 2020A/B ASCII-BCD Conversion

BCD code of 60 always converted to ASCII code 53 (+).
BCD code of 75 always converted to ASCII code 50 (() and BCD code of 55 always converted to ASCII code 51 ()).

APPENDIX B ASSEMBLER INSTRUCTIONS

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Symbols	Meaning
label	Symbolic label, 1-5 alphanumeric characters and periods
m	Memory location represented by an expression
I	Indirect addressing indicator
С	Clear flag indicator
(m,m+1)	Two-word floating point value in m and m+1
comments	Optional comments
[]	Optional portion of field
{ }	One of set may be selected
Р	Program Counter
()	Contents of location
^	Logical product
¥	Exclusive "or"
\vee	Inclusive "or"
А	A-register
В	B-register
E	E-register
An	Bit n of A-register
Bn	Bit n of B-register
b	Bit positions in B- and A-register
(A/B)	Complement of contents of register A or B
(AB)	Two-word floating point value in register A and B
sc	Channel select code represented by an expression
d	Decimal constant
0	Octal constant
r	Repeat count
n	Integer constant
lit	Literal value

MACHINE INSTRUCTIONS

MEMORY REFERENCE	
Jump and Increment-Skip	
ISZ m [,I]	(m) + 1 \rightarrow m: then if (m) = 0, execute P + 2 otherwise execute P + 1
JMP m [,1]	Jump to m; $m \rightarrow P$
JSB m [,I]	Jump subroutine to m: $P + 1 \rightarrow m$; m + 1 $\rightarrow P$
Add, Load and Store	
ADA $\left\{ \begin{array}{c} m \ [,I] \\ lit \end{array} \right\}$	$(m) + (A) \rightarrow A$
ADB $\left\{ \begin{array}{c} m \ [,I] \\ lit \end{array} \right\}$	$(m) + (B) \rightarrow B$
LDA $\left\{ \begin{array}{c} m [, I] \\ lit \end{array} \right\}$	$(m) \rightarrow A$
LDB $\left\{ \begin{array}{c} m [, I] \\ lit \end{array} \right\}$	$(m) \rightarrow B$
STA m[,I]	$(A) \rightarrow m$
STB m [,I]	(B) \rightarrow m
Logical	
AND $\left\{ \begin{array}{c} m [, I] \\ lit \end{array} \right\}$	(m) AND (A) \rightarrow A
XOR $\left\{ \begin{array}{c} m \ [, I] \\ lit \end{array} \right\}$	(m) XOR (A) \rightarrow A
IOR $\left\{ \begin{array}{c} m [, I] \\ lit \end{array} \right\}$	(m) IOR (A) \rightarrow A
$CPA \left\{ \begin{array}{c} m [, I] \\ lit \end{array} \right\}$	If (m) \neq (A), execute P + 2, otherwise execute P + 1.
$CPB \left\{ \begin{array}{c} m [, I] \\ lit \end{array} \right\}$	If (m) \neq (B), execute P + 2, otherwise execute P + 1.

MACHINE INSTRUCTIONS (cont.)

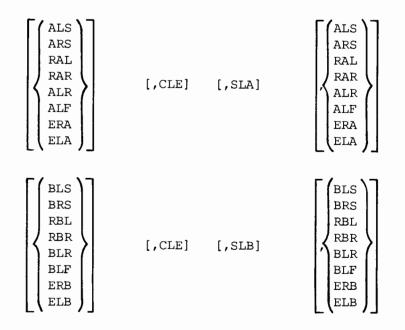
REGISTER REFERENCE

Shift-Rotate

CLE	$0 \rightarrow E$
ALS	Shift (A) left one bit, $0 \rightarrow A_0$, A_{15} unaltered
BLS	Shift (B) left one bit, $0 \rightarrow B_0$, B_{15} unaltered
ARS	Shift (A) right one bit, $(A_{15}) \rightarrow A_{14}$
BRS	Shift (B) right one bit, $(B_{15}) \rightarrow B_{14}$
RAL	Rotate (A) left one bit
RBL	Rotate (B) left one bit
RAR	Rotate (A) right one bit
RBR	Rotate (B) right one bit
ALR	Shift (A) left one bit, $0 \rightarrow A_{15}$
BLR	Shift (B) left one bit, $0 \rightarrow B_{15}$
ERA	Rotate E and A right one bit
ERB	Rotate E and B right one bit
ELA	Rotate E and A left one bit
ELB	Rotate E and B left one bit
ALF	Rotate A left four bits
BLF	Rotate B left four bits
SLA	If $(A_0) = 0$, execute P + 2, otherwise execute P + 1
SLB	If $(B_0) = 0$, execute P + 2, otherwise execute P + 1

MACHINE INSTRUCTIONS (cont)

Shift-Rotate instructions can be combined as follows:



No-operation

NOP Execute P + 1

Alter-Skip

CLA	$0's \rightarrow A$
CLB	O's → B
CMA	$\overline{(A)} \rightarrow A$
CMB	$(B) \rightarrow B$
CCA	$l's \rightarrow A$
CCB	l's → B
CLE	$0 \rightarrow E$
CME	$\overline{(E)} \rightarrow E$
CCE	$1 \rightarrow E$
SEZ	If (E) = 0, execute $P + 2$, otherwise execute $P + 1$
SSA	If $(A_{15}) = 0$, execute P + 2, otherwise execute P + 1
SSB	If $(B_{15}) = 0$, execute P + 2, otherwise execute P + 1

MACHINE INSTRUCTIONS (cont)

(Alter-Skip (cont)

INA	$(A) + 1 \rightarrow A$
INB	$(B) + 1 \rightarrow B$
SZA	If (A) = 0, execute $P + 2$, otherwise execute $P + 1$
SZB	If (B) = 0, execute $P + 2$, otherwise execute $P + 1$
SLA	If $(A_0) = 0$, execute P + 2, otherwise execute P + 1
SLB	If $(B_0) = 0$, execute P + 2, otherwise execute P + 1
RSS	Reverse sense of skip instructions. If no skip instructions precede, execute P + 2

Alter-Skip instructions can be combined as follows:

$$\begin{bmatrix} CLA \\ CMA \\ CCA \end{bmatrix} [,SEZ] \begin{bmatrix} CLE \\ CME \\ CCE \end{bmatrix} [,SSA] [,SLA] [,INA] [,SZA] [,RRS] \\ \begin{bmatrix} CLB \\ CMB \\ CCB \end{bmatrix} [,SEZ] \begin{bmatrix} CLE \\ CME \\ CCE \end{bmatrix} [,SSB] [,SLB] [,INB] [,SZB] [,RSS] \\ \end{bmatrix}$$

INPUT/OUTPUT, OVERFLOW, and HALT

Input/Output

STC	sc	[,C]	Set control bit , enable transfer of one element of data between device and buffer sc sc
CLC	sc	[,C]	Clear control bit . If sc = 0 clear all control bits.
LIA	sc	[,C]	$(buffer) \rightarrow A$
LIB	sc	[,C]	$(buffer) \rightarrow B$
MIA	sc	[,C]	$(buffer)$ (A) \rightarrow A
MIB	sc	[,C]	$(buffer)$ (B) \rightarrow B
OTA	sc	[, C]	(A) \rightarrow buffer sc
OTB	sc	[,C]	(B) $\rightarrow \text{buffer}_{sc}$

MACHINE INSTRUCTIONS (cont)

Input/Output (cont)

STF	sc	Set flag bit . If sc = 0, enable interrupt system. sc ^{sc} l sets overflow bit.
CLF	SC	Clear flag bit . If sc = 0, disable interrupt system. If sc = 1, clear overflow bit.
SFC	SC	If (flag bit) = 0, execute P + 2, otherwise execute P + I. If sc = 1, test overflow bit.
SFS	sc	If (flag bit) = 1, execute P + 2, otherwise execute P + 1. If sc = 1, test overflow bit.

Overflow

CLO		$0 \rightarrow \text{overflow bit}$
STO		l → overflow bit
SOC	[C]	<pre>If (overflow bit) = 0, execute P + 2, otherwise execute P + 1</pre>
SOS	[C]	<pre>If (overflow bit) = 0, execute P + 2, otherwise execute P + 1</pre>

HALT

HLT [sc	[, C]]] Ha	lt	computer
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EXTENDED ARITHMETIC UNIT (requires EAU version of Assembler or Extender Assembler)

МРҮ	m[,I] lit	(A) x (m) \rightarrow (B and A shift) and A shifts (B)
DIV	m[,I] lit	$(B_{\pm msb} \text{ and } A_{sb})/(m) \rightarrow A$, remainder $\rightarrow B$
DLD	m[,I] lit	(m) and (m + 1) \rightarrow A and B
DST	m[,I] lit	(A) and (B) \rightarrow m and m + 1
ASR	b	Arithmetically shift (BA) right b bits, B ₁₅ extended
ASL	b	Arithmetically shift (BA) left b bits, B ₁₅ unaltered, O's to A _{sb}

MACHINE INSTRUCTIONS (cont)

EXTENDED ARITHMETIC UNIT (cont)

RRR	b	Rotate (BA) right b bits
RRL	b	Rotate (BA) left b bits
LSR	b	Logically shift (BA) right b bits, O's to B msb
LSL	b	Logically shift (BA) left b bits, o's to $A_{ sb }$



PSEUDO INSTRUCTIONS

ASSEMBLER CONTROL

NAM	[name]	Specifies relocatable program and its name.
ORG	m	Gives absolute program origin or origin for a segment of relocatable or absolute program.
ORR		Reset main program location counter at value existing when first ORG or ORB of a string was encountered.
ORB		Defines base page portion of relocatable program.
END	[m]	Terminates source language program. Produces transfer to program starting location, m, if given.
REP <statement></statement>	r	Repeat immediately following statement r times.
IFN <statements> XIF</statements>		Include statements in program if control state- ment contains N.
IFZ <statements> XIF</statements>		Include statements in program if control state- ment contains Z.

PSEUDO INSTRUCTIONS (cont)

OBJECT PROGRAM	LINKAGE	
COM	name _l [(siz	e ₁)][,name ₂ [(size ₂)],,name _n [(size _n)]]
		Reserves a block of common storage locations. name, identifies segments of block, each of length size.
ENT	name [,nam	$e_2, \ldots, name_n$]
		Defines entry points, name ₁ , that may be referred to by other programs.
EXT	name [,nam	e ₂ ,,name _n]
		Defines external locations, name ₁ , which are labels of other programs, referenced by this program.

ADDRESS AND SYMBOL DEFINITION

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label	DEF	m[,I]	Generates a 15-bit address which may be refer- enced indirectly through the label.
label	ABS	m	Defines a 16-bit absolute value to be referenced by the label.
label	EQU	m	Equates the value, m, to the label.

PSEUDO INSTRUCTIONS (cont)

CONSTANT DEFINITION

ASC n, <2n characters>	Generates a string of 2n ASCII characters.
DEC d_1 [, d_2 ,, d_n]	Records a string of decimal constants of the form:
	Integer: <u>+</u> n
	<pre>Floating point: +n.n, +n., +.n, +nE+e,</pre>
	+n.nE+e, +n.E+e, +.nE+e
DEX d_1 [, d_2 ,, d_n]	Records a string of extended precision decimals constants of the form
	Floating point: <u>+</u> n, <u>+</u> n.m,
	<u>+</u> n., <u>+</u> .n,
	<pre>+nE+e, +n.nE+e,</pre>
	<u>+n.E</u> +e, +.nE+e
OCT \circ_1 [, \circ_2 ,, \circ_n]	Records a string of octal constants of the form: ± 000000

STORAGE ALLOCATION

BSS	m	Reserves a storage area of length, m.
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PSEUDO INSTRUCTIONS (cont)

ARITHMETIC SUBROUTINE CALLS REQUESTS*

MPY†	${m[,I] \\ lit}$	(A) $x (m) \rightarrow (B_{\pm msb} \text{ and } A sb)$
div†	${m[,I] \\ lit}$	$(B_{+msb} \text{ and } A_{ sb})/(m) \rightarrow A$, remainder $\rightarrow B$
FMP	${m[,I] \\ lit}$	(AB) x (m, m + 1) \rightarrow AB
FDV	${m[,I] \\ lit}$	$(AB)/(m, m + 1) \rightarrow AB$
FAD	${m[,I] \\ lit}$	$(m, m + 1) + (AB) \rightarrow AB$
FSB	${m[,I] \\ lit}$	(AB) - (m, m + 1) \rightarrow AB
DLD†	${m[,I] \\ lit}$	(m) and (m + 1) \rightarrow A and B
DST†	m[,I]	(A) and (B) \rightarrow m and m + 1

⁺For configurations including Extended Arithmetic Unit, these mnemonics generate hardware instructions when the EAU version of the Assembler or Extended Assembler is used.

*Not intended for use with DEX formatted numbers. For such numbers, JSB Machine Instructions must be used.

PSEUDO INSTRUCTIONS (cont)

ASSEMBLY LISTING CONTROL

UNL		Suppress assembly listing output.
LST		Resume assembly listing output.
SKP		Skip listing to top of next page.
SPC	n	Skip n lines on listing.
SUP		Suppress listing of extended code lines (e.g., as produced by subroutine calls).
UNS		Resume listing of extended code lines.
HED	<heading></heading>	Print <heading> at top of each page, where <heading> is up to 56 ASCII characters.</heading></heading>

APPENDIX C ALPHABETIC LIST OF INSTRUCTIONS

ABS	Define absolute value
ADA	Add to A
ADB	Add to B
ALF	Rotate A left 4
ALR	Shift A left 1, clear sign
ALS	Shift A left l
AND	"And" to A
ARS	Shift A right 1, sign carry
ASC	Generate ASCII characters
ASL	Arithmetic long shift left
ASR	Arithmetic long shift right
BLF	Rotate B left 4
BLR	Shift B left 1, clear sign
BLS	Shift B left 1
BRS	Shift B right 1, carry sign
BSS	Reserve block of storage starting at symbol
CCA	Clear and complement A (l's)
CCB	Clear and complement B (l's)
CCE	Clear and complement E (set $E = 1$)
CLA	Clear A
CLB	Clear B
CLC	Clear I/O control bit
CLE	Clear E
CLF	Clear I/O flag
CLO	Clear overflow bit
CMA	Complement A
CMB	Complement B
CME	Complement E
COM	Reserve block of common storage
CPA	Compare to A, skip if unequal
СРВ	Compare to B, skip if unequal

ALPHABETIC LIST OF INSTRUCTIONS (cont)

DEC	Defines decimal constants
DEF	Defines address
DEX	Defines extended precision constants
DIV	Divide
DLD	Double load
DST	Double store
ELA	Rotate E and A left 1
ELB	Rotate E and B left l
END	Terminate program
ENT	Entry point
ERA	Rotate E and A right l
ERB	Rotate E and B right 1
EQU	Equate symbol
EXT	External reference
FAD	Floating add
FDV	Floating divide
FMP	Floating multiply
FSB	Floating subtract
HED	Print heading at top of each page
HLT	Halt
IFN	When N appears in Control Statement, assemble ensuing instructions
IFZ	When Z appears in Control Statement, assemble ensuing instructions
INA	Increment A by 1
INB	Increment B by 1
IOR	Inclusive "or" to A
ISZ	Increment, then skip if zero
JMP	Jump
JSB	Jump to subroutine
LDA	Load into A
LDB	Load into B

ALPHABETIC LIST OF INSTRUCTIONS (cont)

LIA	Load into A from I/O channel
LIB	Load into B from I/O channel
LSL	Logical long shift left
LSR	Logical long shift right
LST	Resume list output (follows a UNL)
MIA	Merge (or) into A from I/O channel
MIB	Merge (or) into B from I/O channel
MPY	Multiply
NAM	Names relocatable program
NOP	No operation
OCT	Defines octal constant
ORB	Establish origin in base page
ORG	Establish program origin
ORR	Reset program location counter
OTA	Output from A to I/O channel
OTB	Output from B to I/O channel
RAL	Rotate A left l
RAR	Rotate A right 1
RRL	Rotate B left l
RBR	Rotate B right 1
REP	Repeat next statement
RRL	Rotate A and B left
RRR	Rotate A and B right
RSS	Reverse skip sense
SEZ	Skip if E = 0
SFC	Skip if I/O flag = 0 (clear)
SFS	Skip if I/O flag = 1 (set)
SKP	Skip to top of next page
SLA	Skip if LSB of $A = 0$
SLB	Skip if LSB of $B = 0$
SOC	Skip if overflow bit = 0 (clear)
SOS	Skip if overflow bit = 1 (set)

ALPHABETIC LIST OF INSTRUCTIONS (cont)

SPC	Space n lines
SSA	Skip if sign $A = 0$
SSB	Skip if sign B = 0
STA	Store A
STB	Store B
STC	Set I/O control bit
STF	Set I/O flag
STO	Set overflow bit
SUP	Suppress list output of additional code lines
SWP	Switch the (A) and (B)
SZA	Skip if $A = 0$
SZB	Skip if $B = 0$
UNL	Suppress list output
UNS	Resume list output of additional code lines
XIF	Terminate an IFN or IFZ group of instructions
XOR	Exclusive "or" to A

APPENDIX D SAMPLE PROGRAM

Following are two sample problems, the second of which implements several options of the Extended Assembler.

PARTS FILE UPDATE

A master file of parts is updated by a parts usage list to produce a new master parts file. A report, consisting of the parts used and their cost, is also produced.

The master file and the parts usage file contain four word records. Each record of the cost report is eleven words long.

The organization of the files is as follows:

Parts Master Files (PRTSM)

Identification Quantity Cost/

Identification field of the Parts Master Files exists in ASCII although the entire record is read and written in binary.

Parts Usage File (PRTSU)

Identification Quantity

The parts usage file has been recorded in ASCII.

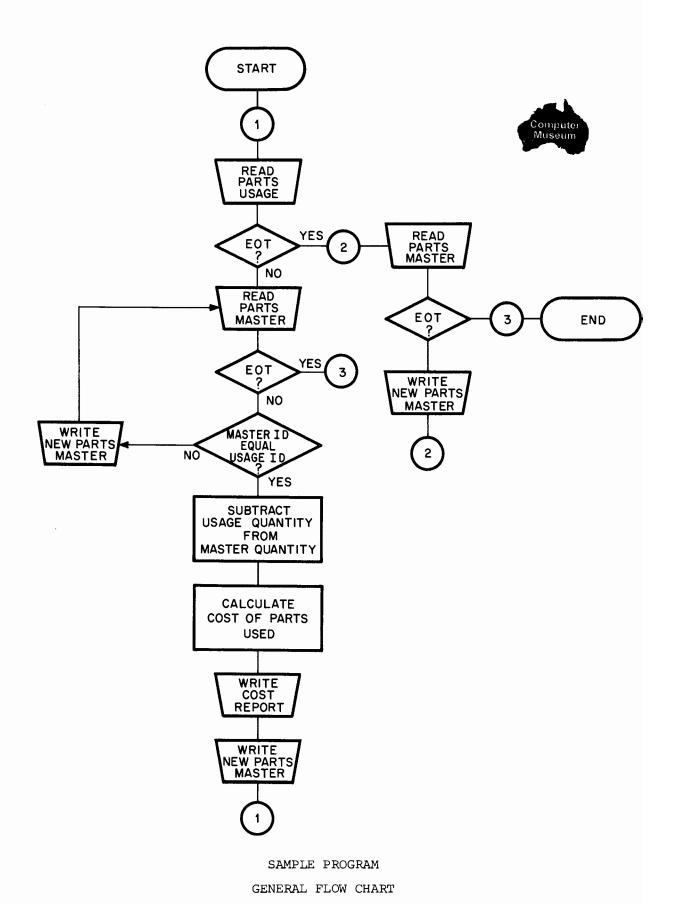
Parts Cost Report (PRTSC)

Identification	Quantity used	#	Cost for Quantity

The Parts Cost Report is recorded in ASCII with spacing and editing for printing.

SAMPLE PROGRAMS (cont)

The sample program reads and writes the files, adjusts the new stock levels, and calculates the cost. External subprograms perform the binary-to-decimal and decimal-to-binary conversions and handle unrecoverable input/output errors, invalid data conditions, and normal program termination. Input/output operations are performed using the Basic Control System input/output subroutine, .IOC.



SAMPLE ASSEMBLER SYMBOL TABLE OUTPUT

PAGE 0001

START R 000000 PRTSM B 000004 PRTSU B 000004 PRTSC B 000023 EOTS1 B 000024 MTEMP B 000025 UTEMP B 000027 SPACS B 000031 DLRSG B 000001 •IOC• X 000002 B 000003 A 000004 B 000003 A 000004 HALT 000005 DTOBI C 000002 BTODI C 000003 BTODO C 000002 SPCFL R 000003 BTODO C 000003 DLD X 000004 MSGU R 000013 CKSTU R 000035 EOTU R 000040 MSGU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTM R 000113 <td< th=""><th></th><th></th><th>0001</th></td<>			0001
PRTSM B 000000 PRTSU B 000004 PRTSC B 000023 EOTS2 B 000024 MTEMP B 000025 UTEMP B 000027 SPACS B 000033 A 000000 B 000001 .10C. X 000001 .10C. X 000003 A 000004 HALT X 000002 DCONV X 000003 BODI C 000003 BTODI C 000003 DTOBI C 000003 DTOBO C 000003 DTODI C 000003 DLD X 00003 DLD X 000040 MSGU R 000013 CKSTU R 00003 DLD X 00003 DLD X 000040 MSGU R 00013 CKSTU R 00013	START	P	
PRTSU B 000004 PRTSC B 000023 EOTS1 B 000024 MTEMP B 000025 UTEMP B 000027 SPACS B 000033 A 000000 B 000001 ·IOC· X 000002 B 000001 ·IOC· X 000002 DCONV X 000003 ABORT X 000003 DTOBI C 000003 BTODI C 000003 DLD X 000003 DLD X 000004 MSGU R 00003 RJCTU R 00003 BCTU R 000040 SCY R 000035 DFEN R <			
PRTSC B 000010 EOTS1 B 000023 EOTS2 B 000024 MTEMP B 000026 SWTMP B 000027 SPACS B 000031 DLRSG B 00000 B 000001 ·IOC· X 000001 ·IOC· X 000003 A 000004 B 000005 DCONV X 000002 BTODI C 000003 BTODI C 000003 BTODI C 000002 BTODI C 000003 BTODO C 000003 BTODI C 000003 BTODI C 000003 DLD X 000003 DLD X 000003 SCTU R 00003 BCDU X 00003 SCTU R 00003 SCTU R 00003 SCTU R		_	
EOTS1 B 000023 EOTS2 B 000024 MTEMP B 000026 SWTMP B 000027 SPACS B 000031 DLRSG B 000033 A 000000 B 00001 ·IOC· X 000002 DCONV X 000002 DCONV X 000003 ABORT X 000004 HALT X 000003 DTOBI C 000003 DTOBI C 000003 BTODO C 000002 BTODI C 000003 BTODO C 000003 BTODO C 000003 DTOBO C 000003 DTOBO C 000003 BTODO C 00003 BTODO C 00002 BTODO C 00003 BTODO			
EOTS2 B 000024 MTEMP B 000025 UTEMP B 000027 SPACS B 000031 DLRSG B 000033 A 000000 B 00001 ·IOC· X 000002 DCONV X 000002 DCONV X 000003 ABORT X 000004 HALT X 000005 DTOBI C 000003 BTODO C 000002 BTODI C 000003 BTODO C 000005 OPEN R 000002 SPCFL R 000003 DLD X 000006 DST X 000007 READU R 00003 DLD X 000007 READU R 000013 CKSTU R 000035 EOTU R 000013 CKSTM R 00013 CKSTM R 00013 CKSTM R 00013 CKSTM R 00013 CKSTM R 00015 EOTM R 000105 EOTM R 000117 HLTSW R 000117 HLTSW R 000137 COMPR R 000117 HLTSW R 000137 CONVM R 000137 CONVC R 000246 MRITC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301 CKSTN R 000301			
MTEMP B 000025 UTEMP B 000027 SWTMP B 000027 SPACS B 000031 DLRSG B 000001 ILCS X 000001 IOC. X 000002 DCONV X 000003 ABORT X 000004 HALT X 000002 DTOBI C 000003 BTODI C 000003 BTODO C 000002 SPCFL R 000003 DLD X 000003 SCTU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTM R 000010 CMMR R 000105 EOTU R 000105 <td></td> <td></td> <td></td>			
UTEMP B 000026 SWTMP B 000027 SPACS B 000031 DLRSG B 000033 A 000000 B 00001 .IOC. X 000002 DCONV X 000003 ABORT X 000004 HALT X 000005 DTOBI C 000000 DTOBO C 000002 BTODI C 000003 BTODO C 000002 SPCFL R 000003 DLD X 000006 DST X 000007 READU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTM R 00013 CKSTM R 00013 CKSTM R 00013 CKSTM R 000105 EOTM R 000105 EOTM R 000107 RJCTM R 000117 HLTSW R 000117 HLTSW R 000137 COMPR R 000117 HLTSW R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000117 HLTSW R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000246 MFY X 000010 CONVM R 000246 MFITC R 000266 RJCTC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000301 CKSTN R 000301		_	
SWTMP B 000027 SPACS B 000033 A 000000 B 000001 ·IOC·X 000002 DCONV X 000003 ABORT X 000004 HALT X 000002 DTONV X 000003 ABORT X 000003 DTOBI C 000002 DTOBO C 000003 BTODO C 000003 BTODO C 000003 DLD X 000002 SPCFL R 000003 DLD X 000003 DLD X 000007 READU R 000013 CKSTU R 000013 CKSTU R 000013 RGU R 000013 RGU R 000010 CMSGU R 000105 EOTM R 000105 EOTM R 000117 HLTSW R			
SPACS B 000031 DLRSG B 000033 A 000001 IOC.X 000001 BCONV X 000002 DCONV X 000003 ABORT X 000004 HALT X 000002 DTOBI C 000002 DTOBI C 000003 BTODI C 000002 SPCFL R 000003 DLD X 000006 DST X 000007 READU R 000013 CKSTU R 00003 DLD X 000007 READU R 000013 CKSTU R 000013 CKSTU R 000010 RJCTU R 000010 RSGU R 000105 EOTM R 000105 EOTM R 000117 HLTSW R 000105 EOTM R 000117 HLTSW			
DLRSG B 000033 A 000001 IOC·X 000001 BCONV X 000002 DCONV X 000003 ABORT X 000004 HALT X 000002 DTOBI C 000002 BTODI C 000002 BTODI C 000003 BTODO C 000002 SPCFL R 000003 DLD X 000006 DST X 000007 READU R 000013 CKSTU R 000035 DUT R 000013 CKSTU R 000013 CKSTU R 000013 CKSTU R 000010 RJCTU R 000105 EOTU R 000105 EOTU R 000105 EOTM R 000105 EOTM R 000105 EOTM R 000117 HLTSW			
A 000000 B 000001 .IOC. X 000001 BCONV X 000002 DCONV X 000003 ABORT X 000004 HALT X 000002 DTOBI C 000003 BTODI C 000002 BTODI C 000003 BTODO C 000002 SPCFL R 000003 DLD X 000007 READU R 000013 CKSTU R 000020 RJCTU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTU R 000010 NSGU R 000105 POTM R 000105 EOTM R 000110 MSGM R 000117 HLTSW R 000137 COMPR R 000137 CONVM R 000213 CONVM R 000140 PROCC R 000140 PROCM R 000157 PROCC R 000165 MPY X 000010 CONVM R 000224 <tr< td=""><td> +</td><td></td><td></td></tr<>	+		
B 000001 .IOC. X 000001 BCONV X 000002 DCONV X 000003 ABORT X 000004 HALT X 000002 DTOBI C 000002 BTODI C 000003 BTODO C 000002 SPCFL R 000003 DLD X 000007 READU R 000013 CKSTU R 000020 RJCTU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTW R 000105 DTM R 000105 COTM R 000105 EOTM R 000110 MSGM R 000117 HLTSW R 000137 COMPR R 000140 PROCM R 000157 PROCC R 000140 CONVM R 000213 CONVM R 000213 CONUL R 000224 CONUL R 000224 CONUL R 000223		-	
• IOC • X 000001 BCONV X 000002 DCONV X 000003 ABORT X 000004 HALT X 000002 DTOBI C 000002 BTODI C 000003 BTODO C 000003 BTODO C 000003 DTOBO C 000003 DTOBO C 000003 DTOBO C 000003 DTOBO C 000003 DLD X 000006 DST X 000007 READU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTM R 000105 EOTU R 000105 EOTM R 000105 EOTM R 000117 HLTSW R 000117 HLTSW R 000117 HLTSW R 000137 COMPR R 000117 HLTSW R 000137 COMPR R 000117 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000246 WRITC R 000266 RJCTC R 000266 RJCTC R 000266 RJCTC R 000301 CKSTN R 000301 CKSTN R 000301			
BCONV X 000002 DCONV X 000003 ABORT X 000004 HALT X 000005 DTOBI C 000002 BTODI C 000003 BTODO C 000003 BTODO C 000002 SPCFL R 000002 SPCFL R 000007 READU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTU R 000051 READM R 00013 CKSTM R 000105 EOTU R 000105 EOTM R 000105 EOTM R 000117 HLTSW R 000137 COMPR R 000137 COMPR R 000137 COMPR R 000137 COMPR R 000157 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000246 WRITC R 000266 RJCTC R 000266 RJCTC R 000266 RJCTC R 000306 RJCTN R 000316		x	
DCONV X 000003 ABORT X 000004 HALT X 000005 DTOBI C 000002 BTODI C 000003 BTODO C 000003 BTODO C 000003 DTOBO C 000003 DTOBO X 000006 DST X 000006 DST X 000007 READU R 000013 CKSTU R 000013 CKSTU R 000013 CKSTU R 000051 READM R 000051 READM R 000051 READM R 000105 EOTU R 000105 EOTM R 000105 EOTM R 000117 HLTSW R 000137 COMPR R 000137 COMPR R 000117 HLTSW R 000137 COMPR R 000157 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000246 WRITC R 000266 RJCTC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000301 CKSTN R 000301			
ABORT X 000004 HALT X 000005 DTOBI C 000002 BTODI C 000003 BTODO C 000003 BTODO C 000003 DTOBO C 000002 SPCFL R 000003 DLD X 000007 READU R 000013 CKSTU R 000035 EOTU R 000013 CKSTU R 000013 CKSTU R 000035 EOTU R 000013 SGU R 000013 RGU R 000010 CKSTM R 000010 RJCTM R 000117 HLTSW R 000117 HLTSW R 000117 HLTSW R 0001157 PROCM R 0001157 PROCC R 000165 MPY X 000010 CONVM R 000224			
HALT X 000005 DTOBI C 000002 BTODI C 000003 BTODO C 000002 STODO C 000003 DTOBO C 000003 BTODO C 000002 SPCFL R 000007 READU X 000013 CKSTU R 000013 SQU R 000010 CKSTM R 000010 SGU R 000110 MSGU R 000117 RLTSW R 000137 COMPR R 000137 COMPR R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000140 CONVM R 000224 CONU2 R 000235 </td <td></td> <td></td> <td></td>			
DTOBI C 00000 DTOBO C 000002 BTODI C 000003 BTODO C 000003 OPEN R 000002 SPCFL R 000003 DLD X 000007 READU R 000013 CKSTU R 000013 CKSTU R 000035 EOTU R 000035 EOTU R 000035 EOTU R 000051 READM R 000051 READM R 000105 EOTM R 000105 EOTM R 000105 EOTM R 000117 HLTSW R 000117 HLTSW R 000137 COMPR R 000137 COMPR R 000157 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 00024 WRITC R 000266 RJCTC R 000266 RJCTC R 000301 CKSTN R 000301 CKSTN R 000306 RJCTN R 000316			
DTOBO C 000002 BTODI C 000003 BTODO C 000005 OPEN R 000002 SPCFL R 000003 DLD X 000006 DST X 000007 READU R 000013 CKSTU R 000035 EOTU R 000035 EOTU R 000035 EOTU R 000051 READM R 000051 READM R 000105 EOTM R 000105 EOTM R 000105 EOTM R 000117 HLTSW R 000137 COMPR R 000137 COMPR R 000157 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 00024 CONU2 R 00024 WRITC R 00026 RJCTC R 00026 RJCTC R 00030 CKSTN R 000306 RJCTN R 000316			+ +
BTODI C 00003 BTODO C 000003 OPEN R 000002 SPCFL R 000003 DLD X 000007 READU R 000013 CKSTU R 000020 RJCTU R 000035 EOTU R 000035 EOTU R 000051 READM R 000051 READM R 000105 EOTM R 000105 EOTM R 000105 EOTM R 000117 HLTSW R 000137 COMPR R 000137 COMPR R 000137 COMPR R 000157 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 00024 CONU2 R 00024 WRITC R 000266 RJCTC R 000266 RJCTC R 000301 CKSTN R 000301 CKSTN R 000306 RJCTN R 000316			
BTODO C 000005 OPEN R 000002 SPCFL R 000003 DLD X 000006 DST X 000007 READU R 000013 CKSTU R 000020 RJCTU R 000035 EOTU R 000035 EOTU R 000051 READM R 000051 READM R 000105 EOTM R 000105 EOTM R 000110 MSGM R 000117 HLTSW R 000137 COMPR R 000137 COMPR R 000137 COMPR R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000266 RJCTC R 000266 RJCTC R 000306 RJCTN R 000316			
OPEN R 000002 SPCFL R 000003 DLD X 000006 DST X 000007 READU R 000013 CKSTU R 000035 EOTU R 000035 EOTU R 000035 EOTU R 000051 READM R 000051 READM R 000105 EOTM R 000105 EOTM R 000117 HLTSW R 000117 HLTSW R 000137 COMPR R 000137 COMPR R 000140 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000266 RJCTC R 000266 RJCTC R 000306 RJCTN R 000316			
SPCFL R 000003 DLD X 000006 DST X 000007 READU R 000013 CKSTU R 000035 EOTU R 000040 MSGU R 000051 READM R 000070 RSGU R 000105 EOTM R 000105 EOTM R 000117 HLTSW R 000137 COMPR R 000137 COMPR R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000137 CONVM R 000140 CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000266 RJCTC R 000301 CKSTN R 000306 RJCTN R 000306			
DLD X 000006 DST X 000007 READU R 000013 CKSTU R 000035 EOTU R 000040 MSGU R 000051 READM R 000070 RJCTM R 000105 EOTM R 000110 MSGM R 000117 HLTSW R 000137 COMPR R 000157 PROCC R 000157 PROCC R 000165 MOV R 000213 CONU1 R 000224 CONU2 R 000235 CONU2 R 000243 CONU2 R 000246 WRITC R 000246 RJCTC R 000276 WRITN R 000301 CKSTN R 000301 CKSTN R 000306 RJCTN R 000306 RJCTN R 0003016			
DST X 000007 READU R 000013 CKSTU R 000020 RJCTU R 000035 EOTU R 000040 MSGU R 000051 READM R 000070 RJCTM R 000105 EOTM R 000105 EOTM R 000110 MSGM R 000117 HLTSW R 000137 COMPR R 000137 COMPR R 000157 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316			
READU R 000013 CKSTU R 000020 RJCTU R 000035 EOTU R 000040 MSGU R 000051 READM R 000070 RJCTM R 000105 EOTM R 000110 MSGM R 000117 HLTSW R 000137 COMPR R 000157 PROCC R 000165 MPY X 000213 CONVM R 000224 CONU1 R 000235 CONU2 R 000246 WRITC R 000266 RJCTC R 000301 CKSTC R 000266 RJCTC R 000301 CKSTN R 000306 RJCTN R 000306			
CKSTU R 000020 RJCTU R 000035 EOTU R 000040 MSGU R 000051 READM R 000070 RJCTM R 000105 EOTM R 000110 MSGM R 000117 HLTSW R 000137 COMPR R 000137 COMPR R 000137 COMPR R 000157 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316			
RJCTU R 000035 EOTU R 000040 MSGU R 000051 READM R 000051 READM R 000070 RJCTM R 000105 EOTM R 000110 MSGM R 000137 COMPR R 000137 COMPR R 000137 PROCM R 000137 CONPR R 000137 PROCC R 000137 CONVR R 000137 CONVR R 000137 CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000261 CKSTC R 000266 RJCTC R 000301 CKSTN R 000301 CKSTN R 000306 RJCTN R 000316			
EOTU R 000040 MSGU R 000051 READM R 000063 CKSTM R 000105 EOTM R 000105 EOTM R 000110 MSGM R 000117 HLTSW R 000137 COMPR R 000137 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000266 RJCTC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316			
MSGU R 000051 READM R 000063 CKSTM R 000105 EOTM R 000105 EOTM R 000110 MSGM R 000117 HLTSW R 000137 COMPR R 000137 COMPR R 000140 PROCM R 000157 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000246 WRITC R 000266 RJCTC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316		R	000040
CKSTM R 000070 RJCTM R 000105 EOTM R 000110 MSGM R 000117 HLTSW R 000137 COMPR R 000137 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000266 RJCTC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316			000051
RJCTM R 000105 EOTM R 000110 MSGM R 000117 HLTSW R 000137 COMPR R 000157 PROCM R 000165 MPY X 00010 CONVM R 000213 CONU1 R 000224 CONU2 R 000246 WRITC R 000266 RJCTC R 000301 CKSTN R 000306 RJCTN R 000316	READM	R	000063
RJCTM R 000105 EOTM R 000110 MSGM R 000117 HLTSW R 000137 COMPR R 000157 PROCM R 000165 MPY X 00010 CONVM R 000213 CONU1 R 000224 CONU2 R 000246 WRITC R 000266 RJCTC R 000301 CKSTN R 000306 RJCTN R 000316	CKSTM	R	000070
EOTM R 000110 MSGM R 000117 HLTSW R 000137 COMPR R 000140 PROCM R 000157 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000266 RJCTC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316		R	000105
MSGM R 000117 HLTSW R 000137 COMPR R 000140 PROCM R 000157 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316		R	
COMPR R 000140 PROCM R 000157 PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000266 RJCTC R 000266 RJCTC R 000301 CKSTN R 000306 RJCTN R 000316	MSGM	R	000117
PROCM R 000157 PROCC R 000165 MPY X 000213 CONVM R 000224 CONU1 R 000235 CONU2 R 000246 WRITC R 000266 RJCTC R 000301 CKSTN R 000306 RJCTN R 000316	HLTSW	R	000137
PROCC R 000165 MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONU2 R 000246 WRITC R 000266 RJCTC R 000301 CKSTN R 000306 RJCTN R 000316	COMPR	R	000140
MPY X 000010 CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316	PROCM	R	
CONVM R 000213 CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000261 CKSTC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316	PROCC	R	000165
CONU1 R 000224 CONU2 R 000235 CONVC R 000246 WRITC R 000261 CKSTC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316	MPY	х	000010
CONU2 R 000235 CONVC R 000246 WRITC R 000261 CKSTC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316	CONVM	R	000213
CONVC R 000246 WRITC R 000261 CKSTC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316	CONUI	R	000224
WRITC R 000261 CKSTC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316	CONU2	R	000235
CKSTC R 000266 RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316	CONVC	R	000246
RJCTC R 000276 WRITN R 000301 CKSTN R 000306 RJCTN R 000316		R	
WRITN R 000301 CKSTN R 000306 RJCTN R 000316			
CKSTN R 000306 Rjctn r 000316			
RJCTN R 000316			
** NO ERRORS*			
	** N	0	ERRORS*

ASMB, R, B, L, T

				NAM		DTE	
0001	00000		CTADT		u۳		
0002		000000	START		00		
0003		026002R		JMP	06		ASSIGN STORAGE & CONSTANTS TO BP
0004	00000		DDTCM	ORB			MASTER PARTS FILE - BINARY.
0005		000000	PRTSM				PARTS USAGE LIST - ASCII.
0006		000000	PRTSU				
0007		000000	PRTSC				PARTS COST REPORT - ASCII.
0008		Ø26063R					
0009		Ø263Ø1R				ITN	
0010		000000	MTEMP				
0011		000000	UTEMP				
0012		000000	SWTMP				
0013		020040	SPACS	ASC	2,		
		020040					
001 4	00033	020044	DLRSG			\$	
0015	00000		Α	EQU	0		
0016	00001		В	EQU	1		
0017				EXT	• I	0C.	PERFORM I/O OPERATIONS USING BCS
0018*							I/O CONTROL ROUTINE.
0019				EXT	BC	ONV	ENTRY POINT FOR DECIMAL(ASCII)
0020*							TO BINARY CONVERSION SUBPROGRAM.
0021				EXT	DC	ONV	ENTRY POINT FOR BINARY TO
0022*							DECIMAL(ASCII) CONVERSION SUB-
0023*							PROGRAM.
0024				EXT	AB	ORT	ENTRY POINT FOR SUBPROGRAM WHICH
0025*							HANDLES UNRECOVERABLE I/O ERRORS
0026*							OR INVALID DATA.
0027				EXT	ΗА	Т	END OF PROGRAM SUBROUTINE.
0028							DTOBO, BTODI(2), BTODO(2)
0029*				000		001(2))	COMMON STORAGE LOCATIONS USED TO
0030*							PASS DATA BETWEEN MAIN PROGRAM
0031*							AND CONVERSION SUBPROGRAMS.
0032	00002			ORR			RESETS PLC AFTER USE OF ORB AT
0033*				UKK			BEGINNING OF PROGRAM.
0034		000000	OPEN	NOP			BEGINATING OF FROGRAM
0035		016006X			CD	1005	STORES EDITING CHARACTERS IN
0055		000031B			35	ACS	STORES EDITING CHARACTERS IN
0036		016007X		DOT	00	TSC+2	AUTRUT AREA FOR BARTS COST
0000		000012B		031	FR	13672	OUTPUT AREA FOR PARTS COST
0037		016007X		DET	PP	TSCAL	PERAPT
0031		000016B		021	PR	TSC+6	REPORT.
0020				1.04	рі	DCC	
0038 0039		060033B		LDA			
		070020B				TSC+8	READ ONE RECORD FROM USAGE LEGT
00 40 0 0 0 1		016001X	RLADU				READ ONE RECORD FROM USAGE LIST
0041		010001				001	LOCATED ON STANDARD UNIT 1
00 42		▲ 026035R				ICTU	(TELEPRINTER INPUT). PRTSU IS
0043		000004B				TSU	ADDRESS OF STORAGE AREA; AREA IS
0044		000004	0.4 C T	DEC		20	4 WORDS LONG.
0045		016001X	CKSTU				CHECK STATUS OF UNIT 1.
0046		040001		OCT		001	
00 47		002020		SSA			
0048		026020R			CK	STU	IF BUSY, LOOP UNTIL FREE.
00 49		001200		RAL			
0050		002020		SSA			
0051		026030R		JMP			
0052		Ø26063R		JMP	RE	ADM	IF COMPLETE, TRANSFER TO SECTION
0053*							WHICH READS MASTER FILE RECORD.

0054 0055 0056	00031	001727 001200 002020		ALF, RAL SSA	ALF	TEST END OF TAPE STATUS BIT (ORIGINAL BIT 05).
0057 0058 0059+	00033	026040R 026004X		JMP	EOTU ABORT	IF SET, GO TO EOT PROCEDURE. IF NOT SET, SOME ERROR CONDITION (UNRECOVERABLE) EXISTS.
0060 0061 0062 0063 0064 0065 0066 0065 0066 0068 0069	00036 00037 00040 00041 00042 00043 00043 00045	006020 026013R 026004X 060023B 072002R 060024B 072140R 016001X 020002 026044R		JMP JMP LDA STA LDA STA JSB OCT	READU ABORT EOTS1 OPEN EOTS2 COMPR • IOC • 20002 EOTU+ 4	CHECK CAUSE OF REJECT. IF UNIT BUSY LOOP UNTIL FREE. ANY OTHER CAUSE IS UNRECOVERABLE ERROR. IF END OF USAGE FILE, ALTER PROGRAM SEQUENCE TO BYPASS SECTIONS THAT READ AND PROCESS USAGE FILE. PRINT MESSAGE ON TELEPRINTER INDICATING EOT.
0070 0071	00047 00050	000051R 000011		DEF DEC	MSGU 9	
0072	00052 00053 00054 00055 00056 00057 00060	0 42 51 6 0 420 40 0 47 50 6 0 20 1 2 5 0 51 50 1 0 43 50 5 0 20 1 0 6 0 44 51 4 0 42 440	MSGU	ASC	9,END OF	USAGE FILE
0073		Ø26063R		_	READM	DEAD A DECORD FROM MACTED BARTS
0074 0075		016001X 010105	READM		•10C• 10105	READ A RECORD FROM MASTER PARTS File on standard unit 05(punched
0076		Ø26105R			RJCTM	TAPE READER). PRTSM IS ADDRESS
0077		000000B			PRTSM	OF STORAGE AREA; AREA IS 4 WORDS
0078 0079		000004 016001X	CKSTM	DEC		LONG. RECORD IS IN BINARY FORMAT CHECK STATUS OF UNIT 5.
0019		040005	CRAIM		40005	CHECK STATUS OF UNIT 5.
0081		002020		SSA		
0082		026070R		JMP	CKSTM	IF BUSY, LOOP UNTIL FREE.
0083		001200		RAL		
0084		002020		SSA	* + 0	
0085 0086		026100R 026140R			COMPR	IF COMPLETE, TRANSFER TO EITHER
0087		001727			ALF	PROCESSING OR WRITE OUTPUT
0088	00101	001200		RAL		DEPENDING ON SETTING OF COMPR.
0089		002020		SSA		TEST FOR END OF TAPE.
0090		026110R 026004X			EOTM ABORT	IF END, GO TO EOT PROCEDURE. If Not, an unrecoverable error
0091 0092*		9200947		Jun	ABORI	EXISTS.
0093		006020	RJCTM	SSB		CHECK CONTENTS OF B FOR CAUSE OF
0094		026063R			READM	REJECT. IF UNIT BUSY, LOOP UNTIL
0095		026004X			ABORT	FREE, OTHERWISE I/O ERROR EXISTS
0096		Ø62137R			HLTSW	ALTER PROGRAM SEQUENCE TO HALT
0097 0098		072315R 016001X			CKSTN+7 •10C•	EXECUTION AFTER LAST RECORD IS WRITTEN PRINT MESSAGE
0098 0099		020002			20002	INDICATING END OF MASTER INPUT.
0100		026112R			EOTM+2	
0101		000117R			MSGM	
0102		000017		DEC		
0103	00117	042516	MSGM	ASC	15, END OF	MASTER PARTS FILE INPUT

	00120	042040				
	00121	047506				
	00122	020115				
	00123	040523				
	00124	052105				
	00125	051040				
	00126	050101				
	00127					
	00130					
	00131					
		046105				
	00133					
	00134					
		0 52 52 4				
0104		026140R		IMP	COMPR	
0105		026005X		-		END OF PROGRAM SUBROUTINE.
			COMPR	-		END OF TROUMAN SUBROUTINES
0106		000000	COMPR		CONUS	CONVERT ID NUMBER FIELDS OF
0107		Ø16224R		-	CONU1	MASTER AND USAGE FILES TO BIN.
0108		016213R			CONVM	LOAD THESE FIELDS FROM TEMPORARY
0109		060026B				STORAGE.
0110		0640258			MTEMP	
0111	00145			CPA	-	COMPARE
0112		Ø26157R		-	PROCM	IF EQUAL, JUMP TO PROCESSING
0113		007004			INB	IF ID NUMBER OF MASTER GREATER
0114		040001		ADA	8	THAN ID NUMBER OF USAGE, DATA IN
0115		002020		SSA		USAGE FILE ERRONEOUS. TERMINATE
0116		026004X		_	ABORT	RUN•
0117		Ø62156R			*+3	IF ID MASTER LESS THAN ID USAGE,
0118		Ø72315R			CKSTN+7	ALTER SEQUENCE: READ NEXT MASTER
0119	00155	Ø263Ø1R		JMP	WRITN	RECORD IMMEDIATELY AFTER WRITING
0120	00156	Ø26063R		JMP	READM	CURRENT MASTER RECORD.
0121	00157	Ø16235R	PROCM	JSB	CONU2	CONVERT QUANTITY FIELD OF USAGE
0122	00160	060002B		LDA	PRTSM+2	FILE TO BINARY AND SUBTRACT FROM
0123	00161	064027B		LDB	UTEMP+1	QUANTITY FIELD OF MASTER AND
0124	00162	007004		CMB.	INB	STORE RESULT.
0125	00163	040001		ADA	В	
0126	00164	070002B		STA	PRTSM+2	
0127	00165	016006X	PROCC	DLÐ	PRTSU	STORE ID OF PARTS USED IN REPORT
	00166	000004B				
0128	00167	016007X		DST	PRTSC	FILE STORAGE AREA.
	00170	000010B				
0129		016006X		DLD	PRTSU+2	STORE QUANTITY OF PARTS USED IN
		000006B				
0130		016007X		DST	PRTSC+4	REPORT FILE STORAGE AREA.
		000014B				
0131		060003B		LDA	PRTSM+3	COMPUTE COST OF PARTS USED.
0132		016010X			UTEMP+1	
		000027B		••••		
0133		070030B		STA	SWTMP+1	
		074027B			SWTMP	
0135		016246R			CONVC	CONVERT RESULT TO DECIMAL
0135		016006X			SWTMP	CONVERT REJUET TO DECIMAL
0130		000027B		040	Swith P	
0137				Det	PPTSC+0	STORE IN REPORT FILE AREA
0137		016007X		021	PRTSC+9	STORE IN REPORT FILE AREA.
a1 22		000021B			+ + 2	ALTER SEQUENCES DEAD NEXT HOASE
0138		Ø62212R		LUA	*+3	ALTER SEQUENCE: READ NEXT USAGE
		072315R			CKSTN+7	
0140	00211	Ø26261R		JWP	WRITC	MASTER RECORD.

a.					0.04.04	
0141		026013R	~~~~~	_	READU	
01 42		000000	CONVM			
0143		016006X		DLD	PRTSM	STORE ID FIELDS IN COMMON
		000000B				
0144	00216	016007X		DST	DTOBI	LOCATIONS TO BE PROCESSED BY
	00217	000000C				
0145	00220	016002X				CONVERSION SUBPROGRAM. ON
0146	00221	062002C		LDA	DTOBO	COMPLETION, STORE RESULTS IN
0147	00222	Ø7ØØ25B		STA	MTEMP	LOCATIONS USED BY PROCESSING
0148	00223	126213R		JMP	CONVM, I	SECTIONS. CONVM APPLIES TO ID OF
0149	00224	000000	CONUL	NOP		MASTER PARTS FILE; CONUL, TO ID
0150	00225	016006X		DLD	PRTSU	OF USAGE; CONU2, TO QUANTITY OF
	00226	000004B				
0151		016007X		DST	DTOBI	USAGE; AND CONVC, TO COST OF
		000000C				
0152		016002X		JSB	BCONV	PARTS(THIS IS A BINARY TO
		062002C			DTOBO	DECIMAL CONVERSION).
		070026B			UTEMP	
		126224R			CONUL	
		000000		-	00110171	
0158		016006X			PRTSU+2	
0121		000006B			11130+2	
A1 5 0				DOT	DTOBI	
0158		016007X		021	DIOBI	
21 F 2		000000C			DODUU	
		016002X			BCONV	
0160		062002C			DTOBO	
0161		070027B			UTEMP+1	
		126235R			CONU2, I	
		000000				
0164		016006X		DLÐ	SWTMP	
		000027B				
0165	00251	016007X		DST	BTODI	
	ØØ2 52	000003C				
0166	00253	016003X		JSB	DCONV	
0167	00254	016006X		DLD	BTODO	
	00255	000005C				
0168	00256	016007X		DST	SWTMP	
	00257	000027B				
0169	00260	126246R		JMP	CONVC, I	
0170	00261	016001X	WRITC	JSB	•10C•	WRITE ONE RECORD OF PARTS COST
0171		020102			20102	REPORT ON STANDARD UNIT 2
0172		026276R			RJCTC	(TELEPRINTER OUTPUT). PRTSC IS
0173		000010B			PRTSC	ADDRESS IN STORAGE AREA; AREA IS
0174		000013		DEC		11 WORDS LONG. RECORD IS IN ASCI
0175		016001X	CKSTC			CHECK STATUS OF UNIT 2.
0176		040002			40002	
0177		002020		SSA		
0178		Ø26266R			CKSTC	IF BUSY, LOOP UNTIL FREE.
0179		001200		RAL		
0180		002020		SSA		
0181		026004X			ABORT	TERMINATE IF ANY I/O ERROR.
0182		026301R			WRITN	IF COMPLETE, TRANSFER TO WRITN.
0182		006020				IF BUSY, LOOP UNTIL FREE.
0183		026261R			WRITC	TERMINATE ON ANY OTHER REJECT
		026261K			ABORT	CONDITION.
0185 0186		026004A 016001X				WRITE ONE RECORD (BINARY) OF
			MICTIN		20104	NEW MASTER PARTS LIST ON UNIT 4
0187		020104 026316P			RJCTN	(TAPE PUNCH). PRTSM (INPUT AREA)
0188	00303	Ø26316R		0.47	1.0011	CHARLE FORGHT FRIJE CIMPUL AKEAJ

0189 00304 000000B DEF PRTSM 0190 00305 000004 DEC 4 0191 00306 016001X CKSTN JSB . IOC. 0192 00307 040004 OCT 40004 0193 00310 002020 SSA 0194 00311 026306R JMP CKSTN 0195 00312 001200 RAL 0196 00313 002020 SSA 0197 00314 026004X JMP ABORT JMP READU 0198 00315 026013R 0199 00316 006020 RJCTN SSB 0200 00317 026301R JMP WRITN 0201 00320 026004X JMP ABORT 0202 END START ** NO ERRORS*

IS ALSO USED AS OUTPUT AREA.

CHECK STATUS OF UNIT 4.

IF BUSY, LOOP UNTIL FREE.

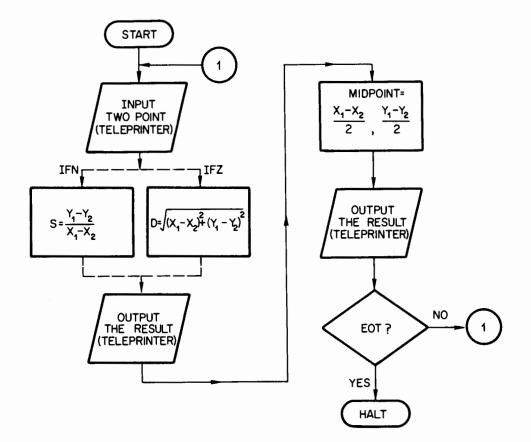
IF BUSY, LOOP UNTIL FREE, OTHER-WISE TERMINATE.

CALCULATING DISTANCE

Program "Line" will either calculate the distance between two points or find the slope of the line connecting the points; then the point equidistant from each point (the mid-point) is calculated.

Data is input using the formatter library routine four n-digit real numbers at a time. The first quantity is the X coordinate of the first point; the second quantity is the Y coordinate of the first point; the third and fourth quantities are the X and Y coordinates of the second point.

The result is output to the teleprinter by the formatter library routine; each quantity cannot be more than an eight-digit real number.



GENERAL FLOW CHART

Below is the source program as it is typed up on the teleprinter. After it are the assembler listings. The first listing results from including the Z option in the control statement. In the second listing the N option has been included in the control statement.

```
NOTE: When the complete data tape has been read and the tape
reader encounters 10 blank feed frames, an EQT message
is typed on the teleprinter and the computer halts.
Thus no halt instruction is needed in the program.
```

HED LINE FORMULI: DISTANCE, SLOPE, MID-POINT PROGRAM LINE WILL EITHER CALCULATE THE DISTANCE BETWEEN TWO POINTS OR FIND THE SLOPE OF THE LINE CONNECTING THE POINTS; THEN THE POINT EQUIDISTANT FROM EACH POINT (THE MID-POINT) IS CALCULATED. DATA IS INPUT USING THE FORMATTER LIBRARY ROUTINE FOUR N-DIGIT REAL NUMBERS AT A TIME. THE FIRST QUANTITY IS THE X COORDINATE OF THE FIRST POINT; THE SECOND QUANTITY IS THE Y COORDINATE OF THE FIRST POINT; * THE THIRD AND FOURTH QUANTITIES ARE THE X AND Y COORDINATES OF THE SECOND POINT. THE RESULT IS OUTPUT TO THE TELEPRINTER BY THE FORMATTER LIBRARY ROUTINE; EACH QUANTITY CANNOT BE MORE THAN AN EIGHT DIGIT REAL NUMBER. NAM LINE START NOP JMP INPUT EXT . IOC., FLOAT, IFIX, SQRT EXT .DIO.,.IOI.,.DTA.,.RAR. EXT .IOR.,.IAR. .DATA DEF DATA •PRIN DEF PRINT DATA BSS 4 FMT ASC 3,(F8.3) FMT2 ASC 8, (F8.3,",",F8.3/) FMT3 ASC 3, (412) SKP * INPUT THE FIRST TWO POINTS; FOUR DATA WORDS INPUT NOP LDA =B5 CLB, INB JSB .DIO. DEF FMT3 DEF *+4 LDA = B4LDB .DATA JSB .IAR. SPC 3

SAMPLE PROGRAMS

* THE	DISTANCE BETWEEN THE TWO POINTS:
	IFZ
	LDA DATA+2
	CMA, INA
	ADA DATA
	SPC 1
	JMP *+5
PRINT	REP 4
	NOP
	SPC 1
	STA PRINT
	SUP
	MPY PRINT
	STA PRINT
	SPC 1
	LDA DATA+3
	CMA, INA
	ADA DATA+1
	STA PRINT+1 MPY PRINT+1
	ADA PRINT
	SPC 1
	JSB FLOAT
	JSB SQRT
	DST PRINT
	XIF
	SPC 3
* F1	ND THE SLOPE OF THE LINE
	IFN
	LDA DATA+2
	CMA, INA
	ADA DATA
	JMP *+5
PRINT	
	NOP
	STA PRINT
	SPC 1
	LDA DATA+3
	CMA, INA
	ADA DATA+1
	CLB
	DIV PRINT
	DST PRINT
	XIF
	SPC 3
* 0U	TPUT THE RESULT
	LDA =B2
	CLB
	JSB .DIO.
	DEF FMT
	DEF *+4
	DLD PRINT
	JSB .IOR.
	JSB .DTA.
	SPC 3

SAMPLE PROGRAMS

* FIND THE MID-POINT OF THE LINE SEGMENT: LDA DATA ADA DATA+2 CLB JSB FLOAT FMP =F.5 DST PRINT SPC 1 LDA DATA+1 ADA DATA+3 CLB JSB FLOAT $FMP = F \cdot 5$ DST PRINT+2 SPC 1 UNL LDA =B2 CLB JSB .DIO. DEF FMT2 DEF *+5 LDA =B2 LDB .PRIN JSB .RAR. JSB .DTA. LST SPC 3 UNS JMP INPUT END START

0001		
START	R	000000
.IOC.	Х	000001
FLOAT	х	000002
IFIX	Х	000003
SORT	х	000004
.DIO.	X	000005
.101.	X	000006
.DTA.	X	000007
.RAR.	x	000010
.IOR.	X	000011
.IAR.	, х	000012
.DATA	R	000002
.PRIN	N R	000003
DATA	R	000004
FMT	R	000010
FMT2	R	000013
FMT 3	R	000023
INPU1	R	000026
PRINT	R	000043
•MPY	х	000013
•DST	х	000014
.DLD	х	000015
.FMP	х	000016
** !	NO 1	ERRORS*

ASMB, R, L, T, Z

PAGE	0002 #01 LIN	E FORMULI	I: DISTANCE, SLOPE, MID-POINT
0002* 0003* 0004*	TWO POINTS O THE POINTS;	R FIND TH THEN THE	THER CALCULATE THE DISTANCE BETWEEN HE SLOPE OF THE LINE CONNECTING POINT EQUIDISTANT FROM EACH
0005*) IS CALCULATED.
0006*			NG THE FORMATTER LIBRARY ROUTINE
0007*			MBERS AT A TIME. THE FIRST
0008*			ORDINATE OF THE FIRST POINT; THE
0009*			HE Y COORDINATE OF THE FIRST POINT;
0010*			QUANTITIES ARE THE X AND Y COORDINATES
0011*			
0012* 0013*			PUT TO THE TELEPRINTER BY THE
0013*			JTINE; EACH QUANTITY CANNOT BE MORE
0014+	00000		AM LINE
0015	00000 000000		
0017	00001 026026R		MP INPUT
0018	00001 0200200		XT ·IOC·,FLOAT,IFIX,SQRT
0019			Computer
0020			XT •DIO•,•IOI•,•DTA•,•RAR• Museum
0021	00002 000004R		
0022	00003 000043R		
0023	00004 000000	DATA B	
0024	00010 024106		SC 3, (F8.3)
	00011 034056		
	00012 031451		
0025	00013 024106	FMT2 A	SC 8, (F8,3,",",F8,3/)
	00014 034056		
	00015 031454		
	00016 021054		
	00017 021054		
	00020 043070		
	00021 027063		
	00022 027451		
0026	00023 024064	FMT3 A	SC 3,(412)
	00024 044462		
	00025 024440		

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PAGE 0003 #01 LINE FORMULI: DISTANCE, SLOPE, MID-POINT

0028* INPUT THE FIRST TWO POINTS; FOUR DATA WORDS 0029 00026 000000 INPUT NOP 0030 00027 062131R LDA =B5 0031 00030 006404 CLB, INB 0032 00031 016005X JSB .DIO. 0033 00032 000023R DEF FMT3 0034 00033 000037R DEF *+4 0035 00034 062132R LDA =B4 0036 00035 066002R LDB .DATA 0037 00036 016012X JSB .IAR.

0039*	THE DIST	ANCE BETWE	EEN THE TWO POINTS
0040			IFZ
0041	00037 062	2006R	LDA DATA+2
0042	00040 003	3004	CMA, INA
0043	00041 042	2004R	ADA DATA
0045	00042 026	047R	JMP *+5
0046	00040 000		T REP 4
0047	00043 000		NOP
0047	00044 000		NOP
	00045 000		NOP
0047	00046 000		NOP
0049	00047 072	2043R	STA PRINT
0050			SUP
0051	00050 016	5013X	MPY PRINT
0052	00052 072	2043R	STA PRINT
0054	00053 062	2007R	LDA DATA+3
0055	00054 003	3004	CMA, INA
0056	00055 042	2005R	ADA DATA+1
0057	00056 072	2044R	STA PRINT+1
0058	00057 016	5013X	MPY PRINT+1
0059	00061 042	2043R	ADA PRINT
		6002X	JSB FLOAT
0062			JSB SQRT
0063	00064 016	5014X	DST PRINT
0064			XIF

0066*	FIND	THE	SLOPE (DF T	HE L	INE
0067					IFN	
0068					LDA	DATA+2
0069					СМА,	INA
0070					ADA	DATA
0071					JMP	*+5
0072			PR	INT	REP	4
0073					NOP	
0074					STA	PRINT
0075					SPC	1
0076					LDA	DATA+3
0077					CMA,	INA
0078					ADA	DATA+1

SAMPLE PROGRAMS

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PAGE 0004 #01 LINE FORMULI: DISTANCE, SLOPE, MID-POINT

0079 0080	CLB DIV PRINT
0081	DST PRINT
0082	XIF

0084*	OUTPL	JT THE	RESULT		
0085	00066	062133	3R	LDA	=B2
0086	00067	006400	9	CLB	
0087	00070	016005	5X	JSB	.DIO.
0088	00071	000010)R	DEF	FMT
0089	00072	000076	SR	DEF	*+4
0090	00073	016015	5X	DLD	PRINT
0091	00075	016011	X	JSB	•IOR•
0092	00076	016007	XX	JSB	•DTA•

0094*	FIND	THE MID-POINT	OF THE LINE	SEGMENT:
0095	00077	062004R	LDA DATA	
0096	00100	Ø42006R	ADA DATA+2	
0097	00101	006400	CLB	
0098	00102	Ø16002X	JSB FLOAT	
0099	00103	Ø16016X	FMP =F.5	
0100	00105	016014X	DST PRINT	
0102	00107	062005R	LDA DATA+1	
0103	00110	042007R	ADA DATA+3	
0104	00111	006400	CLB	
0105	00112	Ø16002X	JSB FLOAT	
0106	00113	Ø16Ø16X	FMP =F.5	
0107	00115	Ø16014X	DST PRINT+2	
0119			LST	

0121			UNS	
0122	00130	Ø26Ø26R	JMP	INPUT
	00131	000005		
	00132	000004		
	00133	000002		
	00134	040000		
	00135	000000		
0123			END	START
** N	O ERROI	RS*		

0001		
START	R	000000
.IOC.	Х	000001
FLOAT	Х	000002
IFIX	Х	000003
SQRT	Х	000004
•DIO•	Х	000005
.101.	Х	000006
•DTA•	х	000007
•RAR•	х	000010
.IOR.	х	000011
•IAR•	х	000012
•DATA	R	000002
•PRIN	R	000003
DATA	R	000004
FMT	R	000010
FMT2	R	000013
FMT3	R	000023
INPUT	R	000026
PRINT	R	000043
•DIV	х	000013
•DST	х	000014
.DLD	х	000015
. FMP	х	000016
** N	D I	ERRORS*

ASMB, R, L, T, N

SAMPLE PROGRAMS

PAGE	0002 #01 LIN	IE FORML	JLI:	DISTANCE, SLOPE, MID-POINT
0002*				ER CALCULATE THE DISTANCE BETWEEN
0003*	TWO POINTS C	R FIND	THE	SLOPE OF THE LINE CONNECTING
0004*	THE POINTS;	THEN TH	E PO	DINT EQUIDISTANT FROM EACH
0005*				IS CALCULATED.
0006*				THE FORMATTER LIBRARY ROUTINE
0007*				ERS AT A TIME. THE FIRST
0008*				DINATE OF THE FIRST POINT; THE
0009*				Y COORDINATE OF THE FIRST POINT;
0010*				UANTITIES ARE THE X AND Y COORDINATES
0011*				
0012*				T TO THE TELEPRINTER BY THE
0013*				INE; EACH QUANTITY CANNOT BE MORE
0014*		T DIGI1		
0015	00000			LINE
0016	00000 000000			
0017	00001 026026	8		INPUT
0018				• IOC • , FLOAT, IFIX, SORT
0019				•DIO•,•IOI•,•DTA•,•RAR•
0020				·IOR., ·IAR.
0021	00002 000004F			
0022	00003 000043			
0023	00004 000000		BSS	
0024	00010 024106	FMT	ASC	3,(F8+3)
	00011 034056			
_	00012 031451			
0025	00013 024106	FMT2	ASC	8,(F8.3,",",F8.3/)
	00014 034056			
	00015 031454			
	00016 021054			
	00017 021054			
	00020 043070			
	00021 027063			
	00022 027451			
0026		FMT3	ASC	3,(412)
	00024 044462			
	00025 024440			

PAGE	0003 #01 LINE FORM	ULI: DISTANCE, SLOPE, MID-POINT
0029 0030 0031 0032 0033 0034 0035	00026 00000 INPUT 00027 062123R 00030 006404 00031 016005X 00032 000023R 00033 000037R 00034 062124R 00035 066002R	O POINTS; FOUR DATA WORDS NOP LDA =B5 CLB,INB JSB .DIO. DEF FMT3 DEF *+4 LDA =B4 LDB .DATA JSB .IAR.
0039* 0040 0041 0042 0043 0044 0045 0046 0047 0048 0047 0048 0049 0050 0051 0052 0053 0055 0056 0055 0056 0057 0058 0057 0058 0057 0058 0057 0058 0059 0060 0061 0062 0063 0064	THE DISTANCE BETWE	EN THE TWO POINTS: IFZ LDA DATA+2 CMA, INA ADA DATA SPC 1 JMP *+5 REP 4 NOP SPC 1 STA PRINT SUP MPY PRINT STA PRINT SPC 1 LDA DATA+3 CMA, INA ADA DATA+1 STA PRINT+1 MPY PRINT+1 ADA PRINT SPC 1 JSB FLOAT JSB SQRT DST PRINT XIF
0066* 0067 0068 0069 0070 0071 0072 0073 0073 0073 0073 0074 0076 0078	FIND THE SLOPE OF 00037 062006R 00040 003004 00041 042004R 00042 026047R PRINT 00043 000000 00044 000000 00045 000000 00045 000000 00045 000000 00045 002007R 00050 062007R 00051 003004 00052 042005R	IFN LDA DATA+2 CMA,INA ADA DATA JMP *+5

PAGE 0004 #01 LINE FORMULI: DISTANCE, SLOPE, MID-POINT

0079	00053	006400	CLB	
0080	00054	Ø16Ø13X	DIV	PRINT
	00055	000043R		
0081	00056	016014X	DST	PRINT
	00057	000043R		
0082			XIF	

0084*	OUTPL	IT THE RESULT		
0085	00060	Ø62125R	LDA	=B2
0086	00061	006400	CLB	
0087	00062	016005X	JSB	•DIO•
0088	00063	000010R	DEF	FMT
0089	00064	000070R	DEF	*+4
0090	00065	Ø16Ø15X	DLD	PRINT
	00066	000043R		
0091	00067	Ø16Ø11X	JSB	.IOR.
0092	00070	016007X	JSB	•DTA•

0094*	FIND	THE MID-POINT	OF THE LINE SEGMENT:
0095	00071	062004R	LDA DATA
0096	00072	042006R	ADA DATA+2
0097	00073	006400	CLB
0098	00074	Ø16002X	JSB FLOAT
0099	00075	Ø16Ø16X	FMP =F.5
	00076	000126R	
0100	00077	Ø16014X	DST PRINT
	00100	000043R	
0102	00101	062005R	LDA DATA+1
0103	00102	042007R	ADA DATA+3
0104	00103	006400	CLB
0105	00104	Ø16002X	JSB FLOAT
0106	00105	Ø16016X	FMP =F.5
	00106	000126R	
0107	00107	Ø16014X	DST PRINT+2
	00110	000045R	
0119			LST
0121			UNS
0122		026026R	JMP INPUT
		000005	
		000004	
		000002	
		040000	
	00127	000000	
0123			END START
NIC) [L T	

** NO ERRORS*

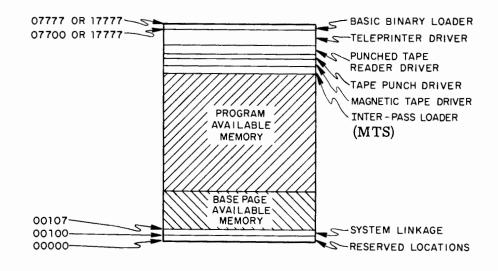
APPENDIX E SYSTEM INPUT/OUTPUT SUBROUTINES

The System Input/Output (SIO) subroutines may be used to perform basic input/ output operations for programs in absolute form.⁺

MEMORY ALLOCATION

These drivers are stored in high memory immediately preceding the Basic Binary Loader. The Teleprinter driver must be loaded first; it is stored in the highest portion of this area. The drivers for the Punched Tape Reader (or Marked Card Reader), the Tape Punch, and the Magnetic Tape Unit may then be loaded. The sequence of loading must fall within this order, depending on your equipment configuration: Line Printer Driver, Punched Tape Reader Driver (or Marked Card Reader), Tape Punch Driver, Magnetic Tape Driver, and if needed, the MTS Boot.

The drivers are accessed through 15-bit absolute addresses which are stored in the System Linkage area starting at location 101₈. The allocation of memory is as follows:



⁺The SIO subroutines are designed for use with FORTRAN, Assembler, Symbolic Editor, etc.; however, they may be used with any absolute object program.

OPERATION AND CALLING SEQUENCE: PAPER TAPE DEVICES

All data transmission is accomplished without interrupt control, and therefore, operations are not buffered by the drivers. Control is not returned to the calling program unitl an operation is completed. Date is transferred to and from buffer storage areas specified in the user program.

The general form of the paper tape input/output calling sequence is:

LDA <buffer length> (words or characters) LDB <buffer address> JSB l0fB,I (f is Input/Output function) <normal return>

Register Contents

When the JSB is performed, the A-Register must contain the length of the buffer storage area and the B-Register, the address of the buffer. Control returns to the location following the JSB. After an input request is completed, the A-Register contains a positive integer indicating the number of characters or words transmitted, or zeros, if an end-of-tape condition occurred.

The digit supplied for f in the JSB instruction determines the paper tape input/output function to be performed. The value of the operand address is the location in the system linkage that contains the absolute address of the driver entry point. The following are available:

- 101 Input
- 102 List Output
- 103 Punch Output
- 104 Keyboard Input-ASCII data is read from teleprinter and printed as it is received.

E-2

If the Teleprinter driver alone is loaded, these locations point to entry points of this driver. If Punched Tape Reader and Tape Punch drivers are in memory, location 101 points to the Punched Tape Reader driver and location 103, to the Tape Punch driver. If the latter are to be used, they must be loaded after the Teleprinter driver.

OPERATION AND CALLING SEQUENCE: MAGNETIC TAPE DRIVER

As with the Paper Tape SIO drivers, all data transmission is accomplished without interrupt control. Control is not returned to the calling program until an operation is completed. (Rewind and Rewind/Standby are the only exceptions to this. In these cases return is made as soon as the command is accepted.)

The general form of the calling sequence is:

LDA <buffer length> or <file count> LDB <buffer address> or <record count> JSB 107B,I OCT <command code> <EOF/EOT/SOT return> <error return> <normal return>

NOTE: Location 107 must contain the address of the magnetic tape driver.

Register Contents

Before initiating read or write operations, the A-Register must contain the buffer length. This will be a positive integer if length is defined in characters and a negative integer if length is defined in words. The

B-Register must contain the buffer address.

Before initiating tape positioning operations, the A-Register must contain the number of files that are to be spaced. A positive integer indicates forward spacing; a negative integer indicates backward spacing. The B-Register contains the number of records that are to be spaced. A positive integer indicates forward spacing; a negative integer indicates backward spacing. The positioning may be defined in terms of any combination of forword or backward spacing of files and records (e.g., space forward two files then backspace three records). If files only or records only are to be spaced, the contents of the other register should be zeros.

The registers are not used when entering the subroutine to perform one of the following operations:

Write end-of-filé Rewind/Standby Write file gap Status Rewind

Linkage Address

107B is the system linkage word that contains the absolute address of the entry point for the Magnetic Tape driver.

On return from a read operation, the A-Register contains a positive value indicating the number of words or characters transmitted.

On return from all operations except Rewind and Rewind/Standby, the B-Register contains status of the operation. (See Status.)

MAGNETIC TAPE OPERATIONS

The magnetic tape driver will perform the following operations. The pertinent operation is specified by the command code which appears after the OCT in the calling sequence.

Operation	Command Code
Read	0
Write	1
Write End-of-File	2
Rewind (Auto mode)	3
Position	4
Rewind/Standby (Local mode)	5
Gap	6
Status	7

Read

One tape record is read into the buffer. The number of characters or words read is stored in the A-Register. The value will be equal to the buffer length except when the data on tape is less than the length of the buffer. One tape record is read to transfer the number of characters specified into the buffer. The number of characters in that record (not the number transferred) will be stored in the A-Register. If the tape record exceeds the buffer length, the data will be read into the buffer until the buffer is filled, the remainder of the record will be skipped. If the length of an input buffer is an odd number of characters, a read operation will result in the overlaying of the character following the last character of the buffer; the subroutine actually transmits full words only.

Three attempts are made to read the record before returning control to the parity error address.

If an EOT condition exists at the time of entry, the command will be ignored and control will be returned to the EOT/EOF address.

If the buffer length specified is 0, control will return to the normal address without any tape movement.

The input buffer storage area can be as large or as small as needed. The number of characters in the tape record will be stored in the A-Register.

Write

The contents of the buffer is written on tape preceded by the record length. Since a minimum of 7 tape characters (12 on 3030) may be written, short records are padded.

If the end-of-tape is detected during the write operation, the normal return is used. The next write operation, however, results in a return of control of the EOF/EOT location; no data is written. If an EOT condition exists at the time of entry, the command will be ignored and control will be returned to the EOT/EOF address.

Write End-of-File

A standard EOF character $(17_8 \text{ for } 2020, 23_8 \text{ for } 3030)$ is written on tape. Control return to the normal location with the EOF status on the B-Register. No gap is written.

If the end of tape was reached on a previous write command, control returns to the EOF/EOT location; the character is written.

Rewind

This command initiates a rewind operation and then immediately returns control to the normal location.

The calling sequence for a Rewind operation consists of:

JSB 107B,I OCT 3 <normal return>



The user need not test status on the rewind operation before issuing the next call.

Position

This is the general command to move the tape. Both file and record operations may be defined in the same operation. Either may be specified for forward or backward spacing. At the completion of the operation the tape will be positioned ready for reading or writing.

An attempt to space beyond the end-of-tape or start-of-tape will terminate the positioning operation and return control to the EOF/EOT/SOT location.

Rewind/Standby

This causes the tape to be positioned at load point and switches the device to local status. Control returns to the normal location immediately after the operation is initiated. The calling sequence for a Rewind/Standby operation consists of:

JSB 107B,I OCT 5 <normal return>

An attempt to issue another call on this device results in a halt (102044). The device must be switched to AUTO before the program can continue.

Gap

This command causes a three-inch gap to be written on the tape.

If the end-of-tape was reached on a previous write command, control returns to the EOF/EOT location; the gap is not written.

Status

This command returns certain status bits in the B-Register. The driver performs a clear command whenever it is entered and as a result the only bits that are valid indicators are:

Start-of-Tape End-of-Tape Write Not Enabled

All other commands (except Rewind and Rewind/Standby) provide valid status replies on return to the program.

The status reply consists only of bits 8-0 and has the following significance:

Bits 8-0	Condition
lxxxxxxx	Local - The device is in local status
xlxxxxxx	EOF - An End-of-File character (17 ₈ for 7 track, 23 ₈ for 9) has been detected while reading, forward spacing, or backspacing.
xxlxxxxx	SOT - The Start-of-Tape marker is under the photo sense head.
xxxlxxxxx	EOT - The End-of-Tape reflective marker is sensed while the tape is moving forward. The bit remains set until a rewind command is given.
xxxxlxxxx	Timing - A character was lost.
xxxxxlxxx	Reject - a) Tape motion is required and the unit is busy. b) Backward tape motion is required and the tape is at load point. c) A write command is given and the tape reel does not have a write enable ring.
xxxxxlxx	Write not enabled - Tape reel does not have write enable ring or tape unit is rewinding.
xxxxxxlx	Parity error - A vertical or logitudinal parity error occurred during reading or writing. (Parity is not checked during forward or backward spacing operations.)
xxxxxxxl	Busy - The tape is in motion or the device is in local status.

Following is a table summarizing the tape commands:

	Command	Call		Reti	ırn
Operation	Code	A	В	A	В
Read	ø	Buffer Length	Buffer Address	Buffer or Record Length	Status
Write	1	Buffer Length	Buffer Address	Buffer Length	Status
Write EOF	2	-	-	_	Status
Rewind (Auto mode)	3	-	-	-	-
Position	4	Number of Files, Direc- tion	Number of Records, Direction	-	Status
Rewind/ Standby (Local mode)	5	-	_	_	_ ,
Gap	6	-	-	-	Status
Status	7	-	-	-	Status

Additional Linkage Addresses

Other locations in the system linkage area contain the following:

- 100 $_{\mbox{8}}$ Used by the standard software system to store a JMP to the transfer address.
- 105 $_{8}$ First word address of available memory.
- 106_8 Last word address of available memory.

The latter two locations may be accessed by an absolute program. The user may store the first word of available memory in 105 by performing the following:

ORG 105B ABS <last location of user program +1>

The last word of available memory is established by the drivers; it is the location immediately preceding the first location used by the last driver loaded.

BUFFER STORAGE AREA

The Buffer Address is the location of the first word of data to be written on an output device or the first word of a block reserved for storage of data read from an input device. The length of the buffer area is specified in the A-Register in terms of ASCII input or output characters or binary output words. For binary input, the length of the buffer is the length of the record which is specified in the first character of the record. ASCII and binary input record lengths are given as positive integers. The length of a binary output record is specified as the two's complement of the number of words in the record.

In addition to describing the buffer area in the calling sequence (for first word of binary input record), the area must also be specifically defined in the program, for example with a BSS instruction.

RECORD FORMATS

ASCII Records (Paper Tape)

An ASCII record is a group of characters terminated by an end-of-record mark which consists of a carriage return, (CR), and a line feed, (LF).

E-11

For an input operation, the length of the record transmitted to the buffer is the number of characters designated in the A-Register, or less if an end-ofrecord mark is encountered before the character count is exhausted. The codes for \overrightarrow{CR} and \overrightarrow{LF} are not transmitted to the buffer. An end-of-record mark preceding the first data character is ignored.

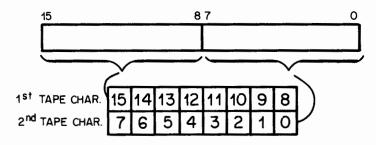
For an output operation, the length of the record is determined by the number of characters designated in the request. An end-of-record mark is supplied at the end of each output operation by the driver.

If a RUB OUT code followed by a (CR) (LF) is encountered on input from the teleprinter or punched tape reader, the current record is ignored (deleted) and the next record transmitted.[†]

If less than ten feed frames (all zeros) are encountered before the first data character from the punched tape reader, they are ignored. Ten feed frames are interpreted as an end-of-tape condition.

Binary Records (Paper Tape)

A binary record is transmitted exactly as it appears in memory or on 8-level paper tape. Each computer word is translated into two tape "characters" (and vice versa) as follows:



⁺ RUB OUT which appears on the teleprinter keyboard is synonymous with the ASCII symbol DEL .

For an output operation, the record length is the number of words designated by the value in the A-Register (the value is the two's complement of the number of words). For input operations, the first word of the record contains a positive integer in bits 15-8 specifying the length (in words) of the record including the first word.

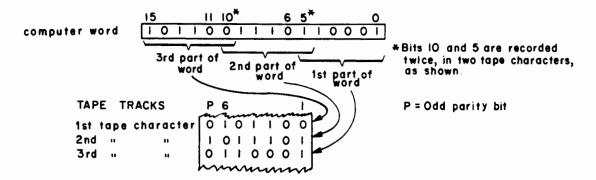
On input operations if less than ten feed frames precede the first data character, they are ignored; ten feedframes are interpreted as an end-of-tape condition. On output, the driver writes four feed frames to serve as a physical record separator.

Binary Records (Magnetic Tape)

The Magnetic Tape subroutine reads and writes binary (odd parity) records only. A record count is supplied by the driver as the first word of the record. This allows automatic padding of short records to the minimum record length with automatic removal of the padded portion of the record on read.

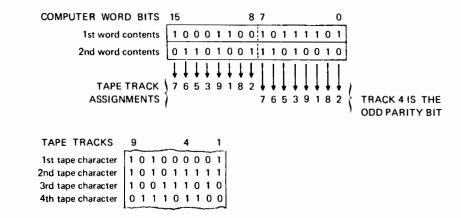
2020 7-LEVEL TAPE

Each computer word is translated into three tape "characters" (and vice versa) as follows:



3030 9-LEVEL TAPE

Each computer word is translated into two tape "characters" by repositioning the bits in the following scheme:



OPERATION AND CALLING SEQUENCE: MARK SENSE CARD READER

The SIO Mark Sense Card Reader Driver overlays the Punched Tape Reader Driver exactly, therefore, only one or the other of these two drivers may be used in any one SIO System configuration. Further, the driver has no binary read capability; if this ability is needed, the BCS Mark Sense Card Reader Driver will have to be used.

All data transmission is accomplished without interrupt control. Execution control is not returned to the calling program until a complete card has been read.

The general form of the calling sequence is:

LDA <character count> (positive) LDB <buffer address> JSB <101B,I> <normal return>

Register Contents

Before the JSB is executed, the A-Register must contain the character count (the buffer length) and the B-Register must contain the buffer address. Control returns to the location following the JSB; then the A-Register will contain the number of characters transmitted not including trailing blanks, or, if a transmission error was detected, it will contain all zeroes.

APPENDIX F CONSOLIDATED CODING SHEET

15	14	13 12	11	10	9	8	7	6	5	4	3	2	1	0
D/I D/I D/I D/I D/I D/I D/I D/I D/I	AND XOR IOR JSB JMP ISZ AD* CP* LD* ST*	001 010 011 001 010 011 100 101 110 111	0 0 1 1 A/B A/B A/B A/B	Z/C Z/C Z/C Z/C Z/C Z/C Z/C Z/C Z/C	•				Memory	y Addre	ss			
15	14	13 12	11	10	9	8	7	6	5	4	3	2	1	0
0	SRG	000	A/B NOP	0	D/E	*LS *RS R*L R*R *LR ER* EL* *LF	000 001 010 011 100 101 110 111 000		CLE	D/E	SL*	*LS *RS R*L R*R *LR ER* EL* *LF	S 001 L 010 R 011 R 100 * 101 * 110	
15	14	13 12	11	10	9	8	7	6	5	4	3	2	1	0
0	ASG	000	A/B	1	CL* CM* CC*	01 10 11	CLE CME CCE	01 10 11	SEZ	SS*	SL*	IN*	SZ*	RSS
15	14	13 12	11	10	9	8	7	6	5	4	3	2	1	0
1	IOG	000	A/B 0 1	1 1 1 1 1 1 1 1 1 1 1 1 1	H/C 0 1 0 0 H/C H/C H/C H/C H/C 1 H/C H/C	HLT STF CLF SFC SFS MI* LI* OT* STC CLC STO CLC SOC SOC	0 0 0 1 1 1 1 1 1 0 0 0 0 0	00 01 01 10 11 00 01 11 11 01 01 10 11	← Select Code → 000 001 000 001 000 001 000 001					
15	14	13 12	11	10	9	8	7	6	5	4	3	2	1	0
1	EAU	000	MPY** DIV** DLD** ASR ASL LSR LSL RRR RRL	τ.	000 010 000 100 100 010 100 100 001 000 001 000 001 000 000 000 001 001 000 000 001 001 000 001				000 000 000 000				f —	
No	r	= A or B. D/I, A/B, Z/C *Second word				/1.								

APPENDIX G

ASSEMBLER ERROR MESSAGES

Errors detected in the source program are indicated by a 1- or 2- letter mnemonic followed by the sequence number and the first 62 characters of the statement in error. The messages are printed on the list output device during the passes indicated:

For Extended Assembler, error listings produced during Pass 1 are preceded by a number which identifies the source input file where the error was found. Pass 2 and 3 error messages are preceded by a reference to the previous page of the listing where an error message was written. The first error will refer to page "O".

Computer

Error		* Museum
Code	Pass	Description
CS	1	Control statement error:
		 The control statement contained a parameter other than the legal set.
		b) Neither A nor R, or both A and R were specified.
		c) There was no output parameter (B, T, or L.)
DD	l	Doubly defined symbol: A name defined in the
		symbol table appears more than once as:
		a) A label of a machine instruction.
		b) A label of one of the pseudo operations:
		BSS EQU ASC ABS
		DEC OC T
		DEF Arithmetic subroutine call DEX
		c) A name in the Operand field of a COM or EXT
		statement.
		d) A label in an instruction following a REP
		pseudo operation.

Error <u>Code</u>	Pass	Description
		e) Any combination of the above.
		An arithmetic subroutine call symbol appears in a
		program both as a pseudo instruction and as a label.
EN	1	The symbol specified in an ENT statement has already
		been defined in an EXT or COM statement.
enøøø <	symbol> 2	The entry point specified in an ENT statement does
		not appear in the label field of a machine or BSS
		instruction. The entry point has been defined in
		the Operand field of an EXT or COM statement, or
		has been equated to an absolute value.
IF	1	An IFZ or an IFN follows either an IFZ or an IFN
		without an intervening XIF. The second pseudo
		instruction is ignored.
IL	1	Illegal instruction:
		a) Instruction mnemonic cannot be used with type of
		assembly requested in control statement. The
		following are illegal in an absolute assembly:
		NAM EXT
		ENT COM ORB Arithmetic subroutine calls
		b) The ASMB statement has an R parameter, and NAM
		has been detected after the first valid Opcode.
IL	2 or 3	Illegal character: A numeric term used in the
		Operand field contains an illegal character (e.g.
		an octal constant contains other than +, -, or \emptyset -7).
		Illegal instruction: ORB in an absolute assembly.
м	1, 2 or 3	Illegal operand:
		a) Operand is missing for an Opcode requiring one.
		b) Operands are optional and omitted but comments
		are included for:
		END
		HLT

Error			
Code	Pass		Description
М	1, 2 or 3	c)	An absolute expression in one of the following instructions from a relocatable program is greater than 77 ₈ .
			Memory Reference
			DEF
			Arithmetic subroutine calls
		d)	A negative operand is used with an Opcode field other than ABS, DEX, DEC, and OCT.
		e)	A character other than I follows a comma in one
			of the following statements:
			ISZ ADA AND DEF JMP ADB XOR Arithmetic JSB LDA IOR subroutine LDB CPA calls STA CPB STB
		f)	A character other than C follows a comma in one
			of the following statements: STC MIB CLC OTA LIA OTB LIB HLT MIA
		g)	A relocatable expression in the operand field of
			one of the following:
			ABS ASR RRL REP ASL LSR SPC RRR LSL
		h)	An illegal operator appears in an Operand field
			(e.g. + or - as the last character).
		i)	An ORG statement appearing in a relocatable pro- gram includes an expression that is base page or common relocatable or absolute.
		j)	A relocatable expression contains a mixture of program, base page, and common relocatable terms.

Error

Code	Pass	Description
		k) An external symbol appears in an operand expres- sion or is followed by a common and the letter I.
		 The literal or type of literal is illegal for the operation code used (e.g., STA = B7).
		<pre>m) An illegal literal code has been used (e.g., LDA = 077).</pre>
		n) An integer expression in one of the following in- structions does not meet the condition $1 \le n \le 16$. The integer is evaluated modulo 2^4 .
		ASR RRR LSR ASL RRL LSL
		o) The value of an 'L' type literal is relocatable.
NO	1, 2, 3	No origin definition: The first statement in the as- sembly containing a valid opcode following the ASMB control statement (and remarks and/or HED, if present) is neither an ORG nor a NAM statement. If the A par- ameter was given on the ASMB statement, the program is assembled starting at 2000; if an R parameter was given, the program is assembled starting at zero. Illegal Opcode preceding first valid Opcode. The statement being processed does not contain an aster- isk in position one. The statement is assumed to
		contain an illegal Opcode; it is treated as a remarks statement.
OP	1,2, or 3	Illegal Opcode: A mnemonic appears in the Opcode field which is not valid for the hardware configura- tion or assembler being used. A word is generated in the object program.
OV	1,2, or 3	Numeric operand overflow: The numeric value of a term or expression has overflowed its limit:

G-4

Error		
Code	Pass	Description
		l>N>16 Shift-Rotate Set
		2 ⁶ -1 Input/Output, Overflow, Halt
		2 ¹⁰ -1 Memory Reference (in absolute assembly)
		2 ¹⁵ -1 DEF and ABS operands; data generated by DEC; or DEX: expressions concerned with program location counter.
		2 ¹⁶ -1 OCT
R?	Before 1	An attempt is made to assemble a relocatable program following the assembly of an absolute program.
SO		There are more symbols defined in the program than the symbol table can handle.
SY	1,2,3	Illegal Symbol: A Label field contains an illegal character or is greater than 5 characters. A label with illegal characters may result in an erroneous assembly if not corrected. A long label is trun- cated on the right to 5 characters.
SY	2 or 3	Illegal Symbol: A symbolic term in the Operand field is greater than five characters; the symbol is truncated on the right to 5 characters.
		Too many control statements: A control statement has been input both on the teleprinter and the source tape or the source tape contains more than one con- trol statement. The Assembler assumes that the source tape control statement is a label, since it begins in column 1. Thus, the commas are considered as illegal characters and the "label" is too long. The binary object tape is not affected by this error, and the control statement entered via the teleprinter is the one used by the Assembler.
TP	1,2, or 3	An error has occurred while reading magnetic tape.
UN	1,2, or 3	Undefined Symbol:

G-5

Error		
Code	Pass	Description
		a) A symbolic term in an Operand field is not de-
		fined in the Label field of an instruction or
		is not defined in the Operand field of a COM
		or EXT statement.
		b) A symbol appearing in the Operand field of on

b) A symbol appearing in the Operand field of one of the following pseudo operations was not defined previously in the source program:

BSS ASC EQU ORG END

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