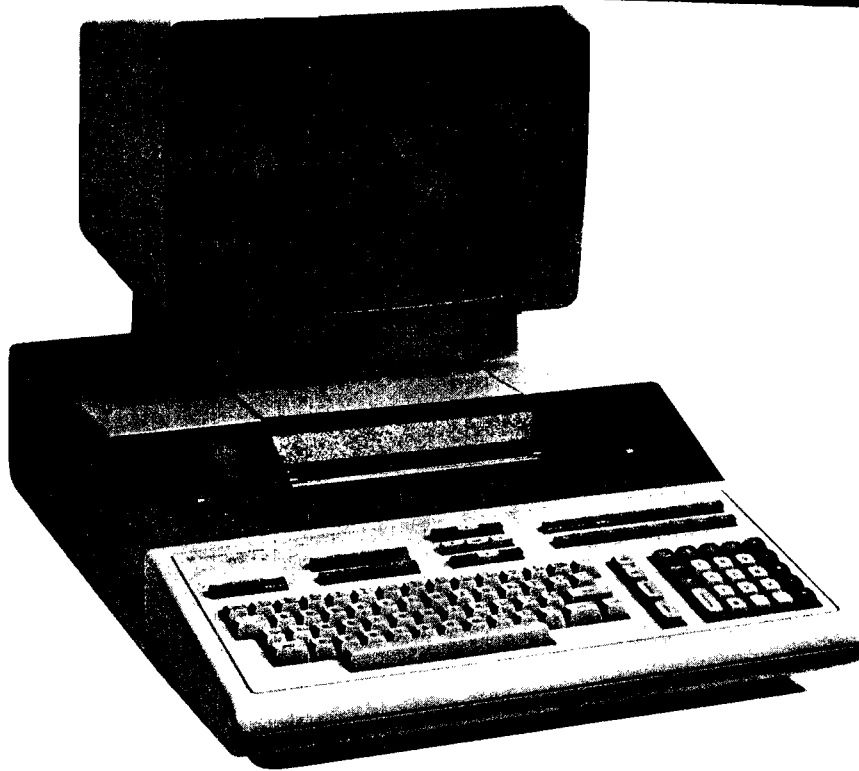


# Advanced Programming ROM

TECHNICAL COMPUTER  
GROUP - MELBOURNE

**LIBRARY COPY**



HP System 45B Desktop Computer



Hewlett-Packard Desktop Computer Division  
3404 East Harmony Road, Fort Collins, Colorado 80525  
(For World-wide Sales and Service Offices see back of manual.)  
Copyright by Hewlett-Packard Company 1979



# Table of Contents

<b>Chapter 1: General Information</b>	
Introduction .....	2
ROM Installation .....	2
Lexical Tables Cartridge .....	3
Manual Requirements .....	3
Error Messages .....	3
Manual Syntax .....	4
<b>Chapter 2: Data Manipulation</b>	
Introduction .....	6
Array Structure and Terminology .....	6
MAT SORT Statement .....	11
Sorting Numeric Data .....	11
Sorting String Data .....	18
MAT REORDER Statement .....	27
MAT SEARCH Statement .....	30
Searching Numeric Arrays .....	30
Searching String Arrays .....	34
<b>Chapter 3: Extended Character Sets</b>	
Introduction .....	38
LEXICAL ORDER IS Statement .....	38
The LEX Function .....	39
Uppercase and Lowercase Functions .....	40
<b>Appendix A: Statement Summary</b> .....	43
<b>Appendix B: User-Defined Lexical Order</b> .....	45
Collating Sequences .....	47
Collating Section .....	48
Uppercase/Lowercase Section .....	50
Mode Section .....	51
Accent Priority .....	51
1 For 2 Character Replacement .....	53
2 For 1 Character Replacement .....	57
"Don't Care" Characters .....	58
ASCII Character Codes .....	59
Roman Extension Character Codes .....	60
ASCII Table .....	62
French Table .....	67

German Table .....	72
Spanish Table .....	77
Swedish Table .....	82
<b>Appendix C: Error Messages .....</b>	<b>87</b>
Mainframe Errors .....	87
I/O Device Errors .....	90
Advanced Programming ROM Error Messages .....	93
Mass Storage ROM Errors .....	91
Graphics ROM Errors .....	91
I/O ROM Errors .....	92
<b>Subject Index .....</b>	<b>94</b>
<b>Sales and Service Offices .....</b>	<b>97</b>

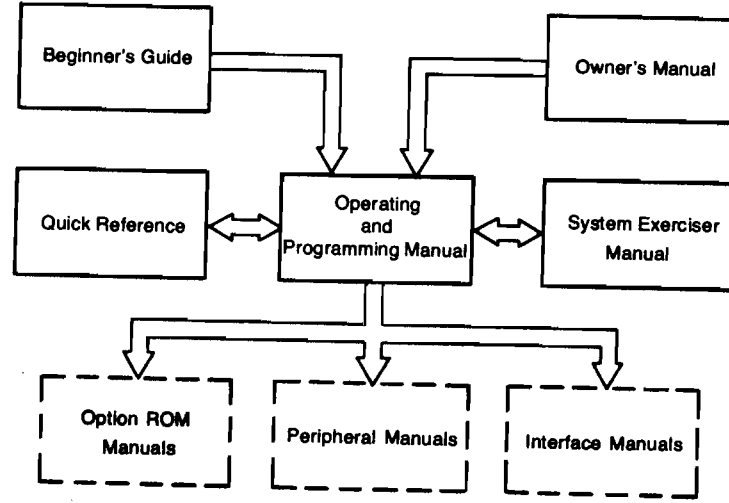
## Printing History

Periodically, this manual is updated. Each new edition of this manual incorporates all material updated since the previous edition. Each new or revised page is indicated by a revision date.

The date on the back cover changes only when each new edition is published.

First Printing...May 1, 1979

## Manual Map



# General Information

## Overview

Your HP System 45B Advanced Programming ROM provides extended programming capabilities such as sorting list data in numerical or lexical order, defining a lexical order and searching lists for conditions which you specify. The ASCII, French, German, Spanish and Swedish lexical orders are provided.

You should be familiar with the basic operation of your System 45B before attempting to use the Advanced Programming ROM and this manual. Refer to your System 45B Operating and Programming manual for this information.

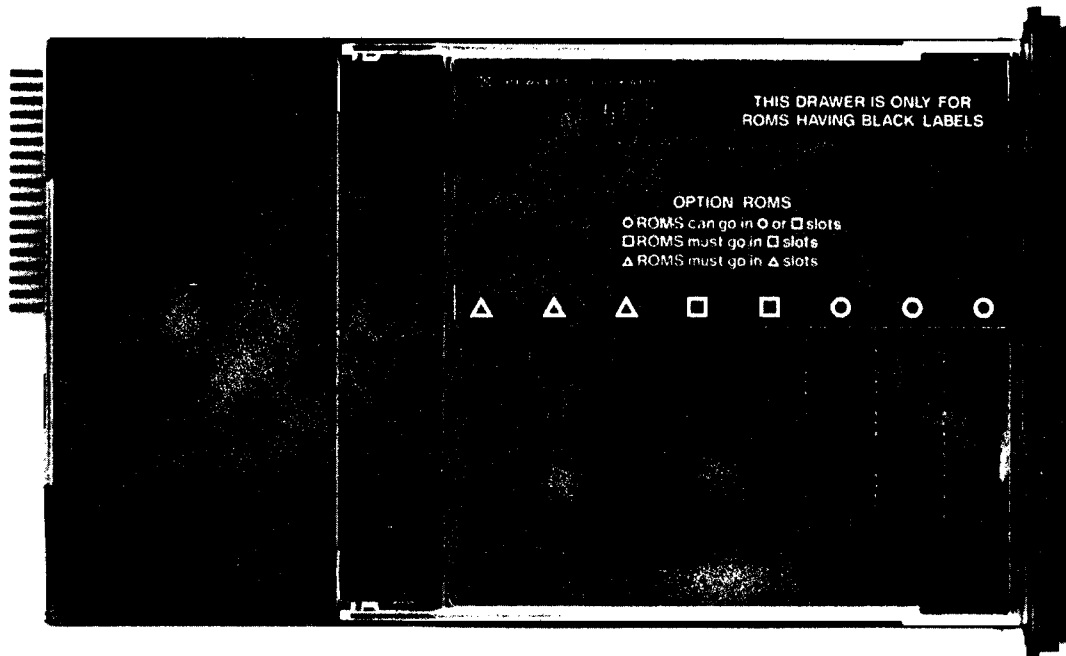
# Chapter 1

The System 45B Advanced Programming ROM provides extended computing capabilities for your System 45B Desktop Computer. This ROM enables you to perform such functions as ordering list data in numerical or lexical order and searching lists for conditions which you specify. These features can be very useful in the areas of mathematics, statistics, and information processing. This manual explains and demonstrates the programming features provided by the Advanced Programming ROM.

### ROM INSTALLATION

The Advanced Programming ROM (HP Part Number 09845-81314) is plugged into the right ROM drawer (black-labeled ROMs). The installation procedure is as follows:

1. Turn off computer power.
2. Slide the right ROM drawer all the way out (see photo).



3. Open the plastic cover by squeezing the sides of the cover and lifting to gain access to the ROM connectors.
4. Press the Advanced Programming ROM onto any available  $\circ$  drawer connector so that it seats all the way down.
5. Close the plastic cover and slide the ROM drawer back into the computer until the drawer is flush with the computer housing.
6. Turn the computer on.

---

**CAUTION**

POWER TO THE COMPUTER MUST BE TURNED OFF WHILE INSTALLING ROMS OR DAMAGE TO THE ROM OR TO THE COMPUTER MAY RESULT.

---

## Lexical Tables Cartridge

An Advanced Programming ROM Lexical Tables Cartridge (P/N 09845-90448) is provided for use with the Advanced Programming ROM. The cartridge contains ASCII and local language collating tables which can be modified for particular collating applications. Descriptions of the tables and instructions for their use are found in Appendix B.

## Manual Requirements

Before using this manual or the Advanced Programming ROM, you should be familiar with the basic operating procedures of the System 45B Desktop Computer as explained in its Operating and Programming Manual.

## Error Messages

The Advanced Programming ROM adds 15 additional error messages (330 through 344) to the System 45B mainframe error message list. Explanations of Advanced Programming error messages as well as mainframe error messages are found in Appendix C of this manual.

## Manual Syntax

The following conventions apply to the syntax used in this manual:

- Dot Matrix - All items in dot matrix must be entered as shown.
- [ ] - All items in brackets are optional.
- ... - Three dots indicate that successive parameters are allowed, when each is separated by a comma.



# Data Manipulation

- page 11 ● **MAT SORT** (orders data records within an array)
- page 24 ● **MAT REORDER** (orders an array according to the contents of an existing pointer array)
- page 30 ● **MAT SEARCH** (provides information about user-defined conditions within an array)

## Terms

- **Record** – represents data which is manipulated as a unit within an array.
- **Key** – a data item within a record used to identify the record for sorting purposes.
- **Key specifier** – a format used in a syntax to specify primary and/or secondary keys within a record.
- **Pointer array** – a one-dimensional numeric array which contains the sorting order of specified records.
- **Location specifier** – a format used within a **MAT SEARCH** statement syntax to specify the locations to be searched.

## Statement Syntax

**MAT SORT** source array (key specifier) [ [substring specifier] ]  
[DES] [, (secondary key specifier) [ [substring specifier] ] [DES]... ] [TO pointer array]

**MAT REORDER** object array BY pointer array [, dimension specifier]

**MAT SEARCH** source array (location specifier), condition; variable [, starting address]

condition:

LOC (relational operator-expression)  
#LOC (relational operator-expression)  
MAX  
MIN

# Chapter 2

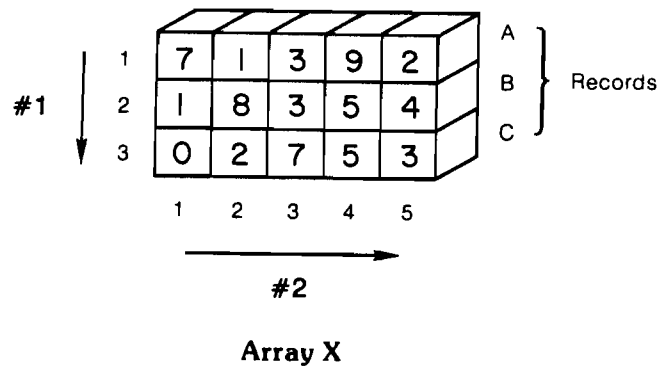
## Introduction

The statements described in this chapter are used as programming aids to manipulate data. They concern the sorting and searching of numeric and string arrays. Each of these statements can be executed from within a program or from the keyboard. Examples of each statement are included to demonstrate its use.

### Array Structure and Terminology

For a better understanding of the sorting and searching processes introduced in this manual, a brief description of array structure and terminology is helpful.

The following illustration represents a two-dimensional array containing numeric data.



It is dimensioned 3 rows by 5 columns as follows:

```
OPTION BASE 1
DIM X(3,5)
```

The subscripts within parentheses (3,5) are written in the order in which the array is dimensioned (i.e., in the order of the numbered arrows). Throughout this manual dimensioning assumes OPTION BASE 1 (i.e., numbering begins with 1, not 0). The dimension statement specifies that the array name is X and that it is three rows by five columns.

An array **record** represents data (string or numeric) which is manipulated as a unit within an array. For example, if you arrange a list of names in a specified order, each name is considered a record within that list. Similarly, an array might contain a list of social security numbers. If you arrange them in a particular order, each social security number is considered a record within that array.

In the previous example, if array X is to be rearranged by rows, each row is considered a record and ordering is performed along the first dimension. That is, the positions of the rows in relationship to each other along the first dimension are rearranged. Likewise, if the array is to be reordered by column, each column is considered a record and reordering is performed along the second dimension.

In the previous illustration, assume that each row is a record (A,B, and C). If the records are to be sorted, a **key** is needed to determine how each row is to be ordered in relationship to the others. Assume that the records are to be sorted in ascending order according to the value of the first number in each record. The first column, then, contains the keys for this sort. The sorting process sorts the keys in ascending order along the dimension in which they lie. In so doing, the records in which the keys lie are rearranged.

The keys selected are described by the following format:

$$(*, 1)$$

The asterisk indicates a varying subscript. The first dimension subscript is varied over its range of values (1 to 2 to 3) to designate a key in all three records (rows). The second dimension subscript is fixed at 1 to indicate that the keys lie in the first column of each record.

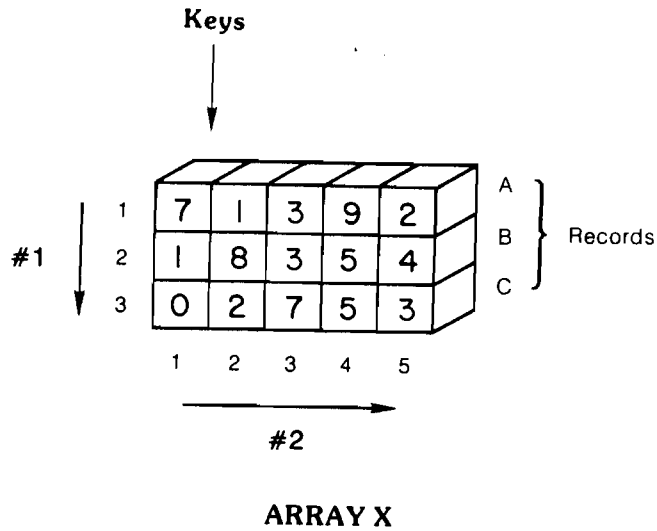
The format  $(*, 1)$  is called a **key specifier**. The asterisk replaces the subscript which corresponds to the dimension along which the records lie, in this case, the rows. The number 1 indicates where the key is located within each record. In this way, the key specifier indicates how the array is partitioned into records, and where the keys are located within those records.

If all the combinations of the key specifier in the previous example are listed, a description of the individual keys is obtained as follows.

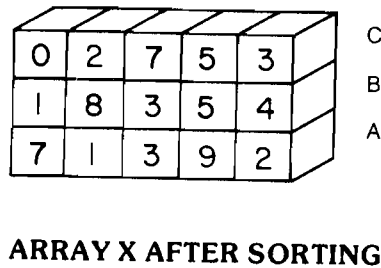
1st Subscript	2nd Subscript	Key
1	1	1st
2	1	2nd
3	1	3rd

## 8 Data Manipulation

The shaded cells in the next figure represent the keys described by the key specifier.



Upon execution, the sorting process arranges the records in ascending order along the first dimension according to the values of their respective keys. The sorted array is shown next.



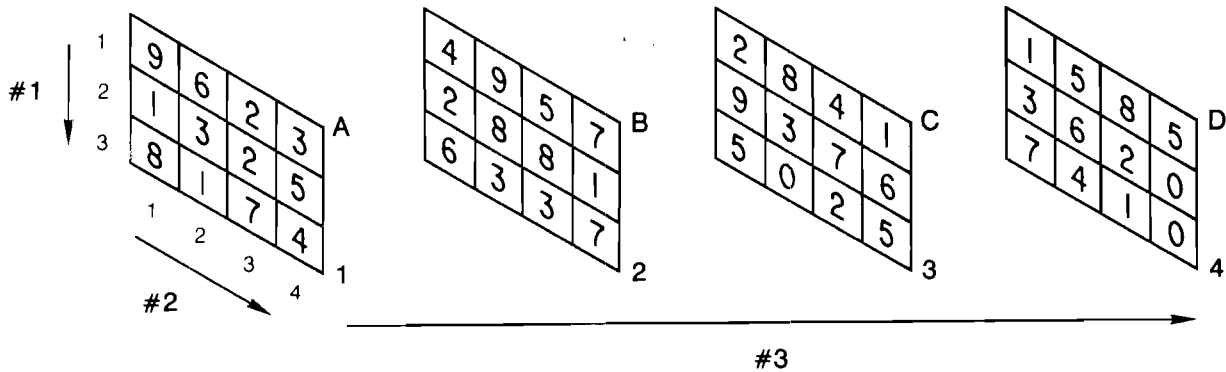
The records, or rows, are repositioned in relationship to each other according to the value of their keys. The contents of the records are left unchanged.

Note that any corresponding numbers within the rows could be selected as keys. The column containing the keys would be described by the key specifier to reflect the proper keys.

**HP Computer Museum**  
**[www.hpmuseum.net](http://www.hpmuseum.net)**

**For research and education purposes only.**

As another example, assume array Y is a three dimensional numeric array as shown.



ARRAY Y WITH KEYS

It is dimensioned in the order of the numbered arrows (3 rows by 4 columns by 4 planes):

```
OPTION BASE 1
DIM Y(3,4,4)
```

Assume that the array is divided into records (A,B,C, & D) along the third dimension, that is, each plane is considered a record. The keys selected must also lie along the third dimension since there must be one for each record. If the upper right number in each record is designated as a key, the key specifier is

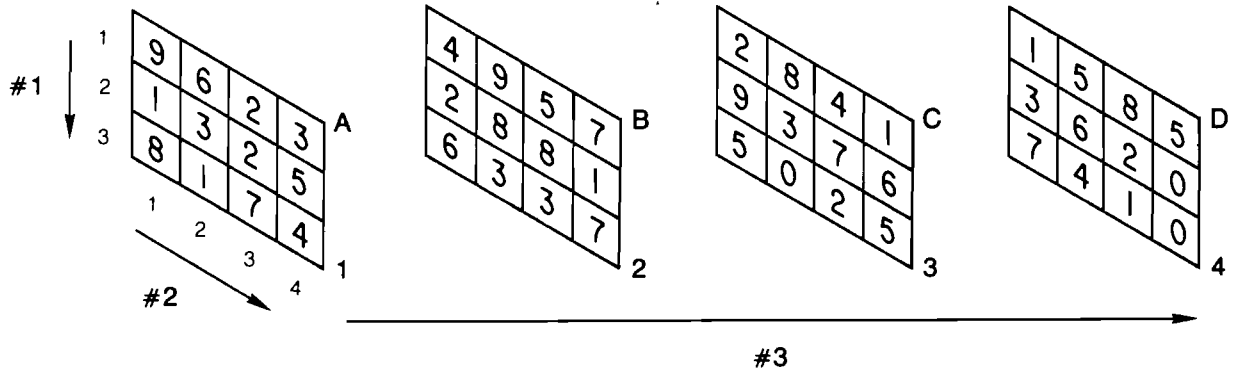
```
(1,4,*)
```

where the asterisk indicates that the third subscript is varied over its entire range of values (1 through 4). In this way, a total of four keys is described as shown.

1st Subscript	2nd Subscript	3rd Subscript	Key
1	4	1	1st
1	4	2	2nd
1	4	3	3rd
1	4	4	4th

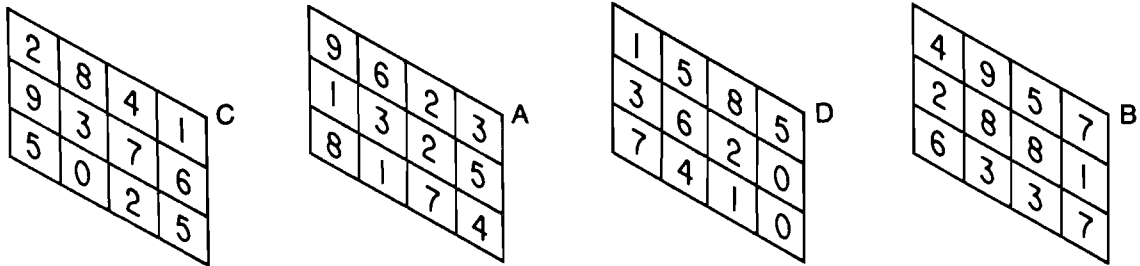
10 Data Manipulation

The shaded cells in the following illustration represent the four keys described by the key specifier (1,4,\*).



ARRAY Y WITH KEYS

If the array records are sorted in ascending order, they are rearranged along the third dimension according to the values of their keys. The following illustration represents the array Y after it is sorted.



ARRAY Y AFTER SORTING

Note that the records have been rearranged according to the ascending values of their keys.

## MAT SORT Statement

### Sorting Numeric Data

The MAT SORT statement is used to order data records in an array. If the data is numeric, it can be sorted in either ascending or descending order. If the array contains string data, it can be sorted in either lexical (alphabetical) or reverse lexical order.

Syntax:

```
MAT SORT numeric source array (key specifier)
```

The source array represents the array in which sorting is performed. The key specifier allows you to specify the keys by which the records of data are sorted. If the source array is one-dimensional, no key specifier need be included.

Example:

```
MAT SORT A(1,*,3)
```

In this case, A is a numeric array in which sorting is performed. The key specifier indicates that A is a three dimensional array and that its records lie along the second dimension. The position of the asterisk within the key specifier determines which subscript is varied, and therefore, how the array is divided into records. If the first subscript were to identify records instead of the second, MAT SORT A(\*,2,3) would be entered. The array in which sorting is performed can have no more than six dimensions. Numeric sorting comparisons are performed in the INTEGER mode for integer-precision arrays and in the REAL mode for short or real-precision arrays.

The default order of sorting is ascending. In the previous example statement, the records are sorted in ascending order. If descending order is desired, the correct entry is

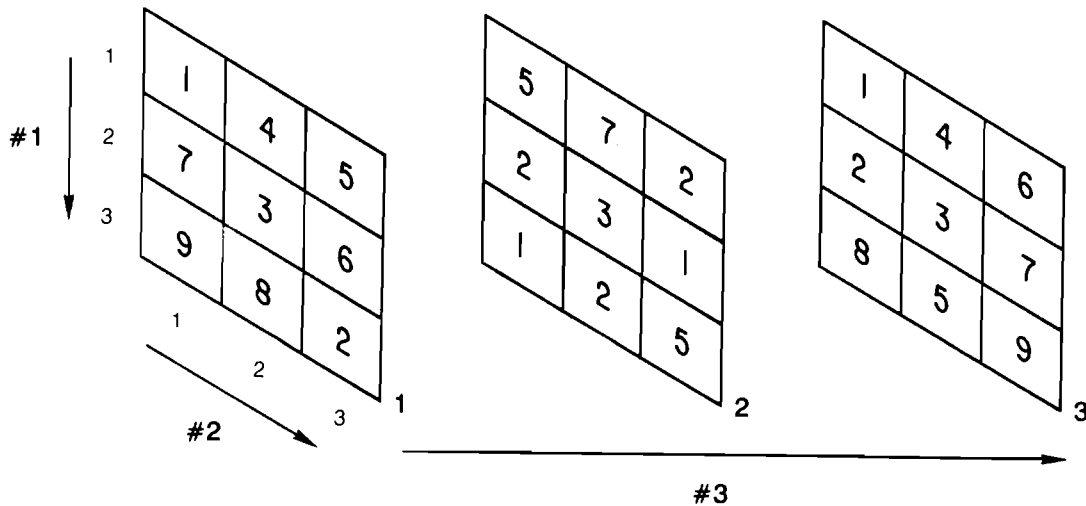
```
MAT SORT A(1,*,3) DES
```

where the letters DES included after the key specifier indicate descending order.



## 12 Data Manipulation

The following example serves to explain the MAT SORT process. Assume array Data is a three dimensional numeric array containing random numbers.



**ARRAY Data BEFORE SORTING**

This illustration is a graphical representation of the array and its contents. Each record is numbered as are its rows and columns. The numbered arrows represent the order in which the array was originally dimensioned.

Assume that the records are to be sorted in ascending order and that the numbers in the upper right location in each record are designated as keys. Since the records all lie along the third dimension, the third subscript is replaced by an asterisk in the key specifier. The proper sorting statement is

```
MAT SORT Data(1,3,*)
```

where the first two numbers indicate the upper right location of each record. The asterisk indicates that the records are selected by varying the third subscript over its range of values (1 to 2 to 3). After the sort is performed, the array is rearranged as shown next.

1	4	5
7	3	6
9	8	2

1

5	7	2
2	3	1
1	2	5

2

1	4	6
2	3	7
8	5	9

3

### ARRAY Data AFTER SORTING

Note that the order of the records is changed according to the value of their keys.

A pointer array can be specified in a sorting statement to maintain a record of how the source array should be rearranged.

Syntax:

```
MAT SORT source array (key specifier) TO pointer array
```

When a pointer array is included in a MAT SORT statement, the sorting process does **not** rearrange the source array. Instead, it fills the pointer array with a series of numbers representing the order of the source array records as if they were sorted. In this way, the source array is not disturbed, but the order in which its records should be sorted is maintained for future use.

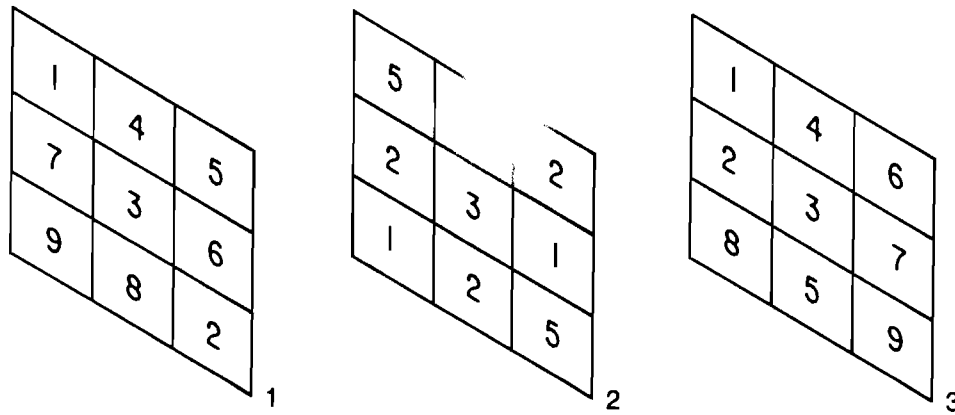
The pointer array must be a one-dimensional numeric array and it must be the same length as the range of records to be sorted. The pointer array does not contain the contents of the records, but rather, a series of numbers representing the order in which the records would be sorted.

For example, assume that the sorting statement in the previous example is modified to include the pointer array Point. The pointer array must be dimensioned to be three elements in length to accommodate the number of records in the source array ( DIM Point(3) ). The sorting statement is modified as follows:

```
MAT SORT Data (1,3,*) TO Point
```

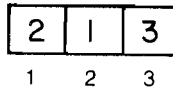
14 Data Manipulation

Upon execution, the source array remains unchanged as shown.



ARRAY Data AFTER SORTING

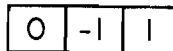
However, the pointer array now contains the order of the records as if they were rearranged.



POINTER ARRAY Point AFTER SORTING

Note that the numbers in the pointer array correspond to the sorted order of the records in the previous example. However, the source array is not rearranged.

The numbers in the pointer array are the values that the varied subscript would assume in designating each record. For example, if Data was dimensioned DIM Data (3,3,-1:1), then Point would contain



This allows the pointer array to be used for indirect reference into the original array. Thus, after sorting this example, Data (3,1, Point(2) ) = 9.

If two items described by the key specifier are identical, a secondary key specifier can be used to complete the sort. The sorting process utilizes the order of the data described by the secondary key specifier to arrange the records containing identical primary keys. Should further identical data be described by the secondary key specifier, additional key specifiers can be included. The asterisks must appear in the same respective positions in all related key specifiers.

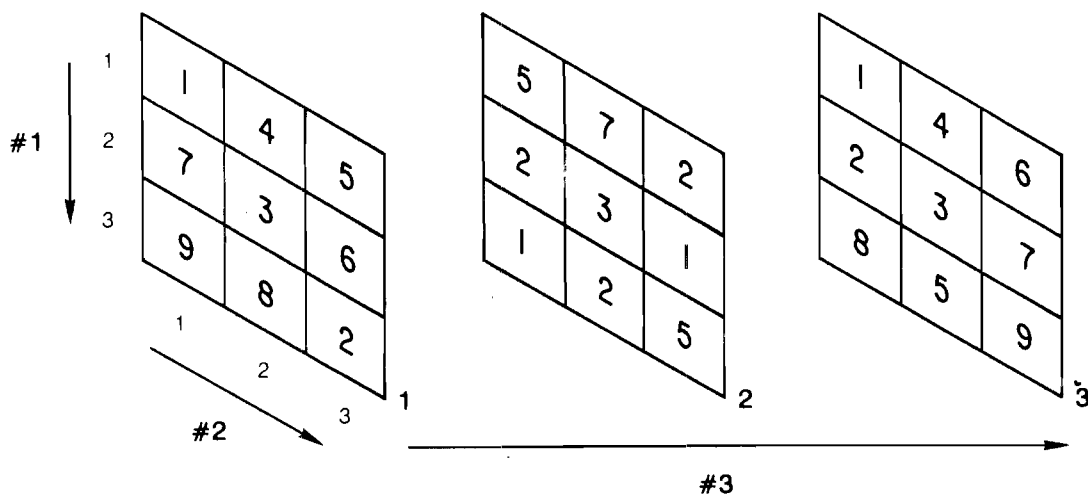
Syntax:

```
MAT SORT source array (key specifier) [, (secondary key specifier) ...]
```

Note that commas are used to separate individual key specifiers. All key specifiers in a given MAT SORT statement must partition the array into records in the same way. That is, the asterisk must always appear in the same position.

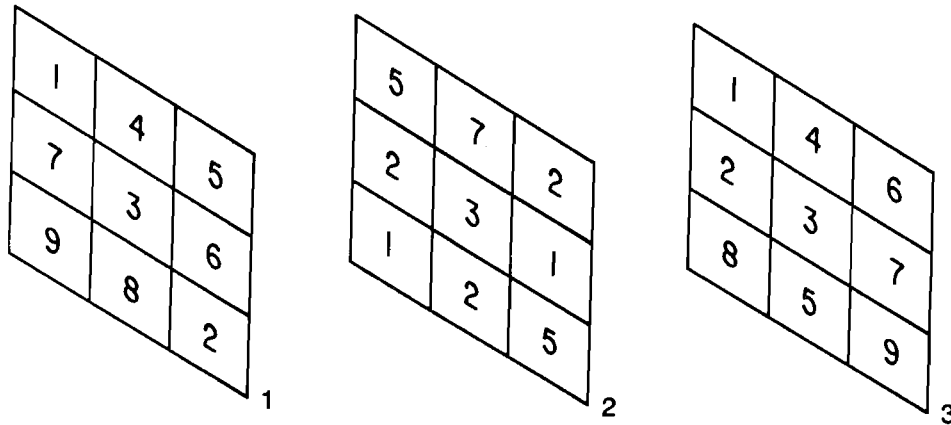
Referring to the previous array example, assume that the records are to be sorted using their upper left location contents as keys. The correct sorting statement is `MAT SORT Data(1,1,*)`, where `(1,1,*)` specifies the upper left numbers as keys. An ambiguity would develop, however, since the number "1" appears in the keys of both records one and three. The values of another series of keys can be used to determine which "1" should be ordered first. Assume that the numbers in the lower right locations are described by a secondary key specifier. The proper sorting statement is:

```
MAT SORT Data (1,1,*), (3,3,*)
```



ARRAY Data WITH PRIMARY & SECONDARY KEY SPECIFIERS

Since identical numbers have been encountered in records one and three, the sorting process examines the lower right keys of the records (as described by the secondary key specifier(3,3,\*) ) to complete the sort. Since the "2" of record one precedes the "9" of record three, record one is ordered before record three. Upon execution of the sorting statement, array Data is rearranged as shown.



ARRAY Data AFTER SORTING

If identical items cannot be differentiated, then the relative order of the sorted records is indeterminate (due to the nature of the sort).

The program shown next demonstrates a use of numeric sorting. A list of salespeople, their districts, and their sales volumes is sorted according to district and volume.

```

10 ! *****
20 ! ***      This program demonstrates the use of the      ***
30 ! ***                MAT SORT statement.                ***
40 ! *****
50
60 OPTION BASE 1
70 DIM Salespeople$(9)[20],Sales(12,2,9),Rank(9) ! Dimension
80                                     ! source and pointer arrays.
90
100 ! *****
110 ! *** Data represents district, sales volume, salespeople. ***
120 ! *****
130
140 DATA 3,3000,1,6000,2,1900,1,1200,3
150 DATA 1400,1,2500,3,900,2,2200,2,1700
160 DATA JILL TANDY,DON DEEDS,RICK HILL,SAM SPADE,NICK DANGER
170 DATA JOE JONES,HARRY WHITE,BARB SMITH,H.J. STEED
180
190 ! *****
200 ! ***                Print headings and sort.                ***
210 ! *****
220
230 PRINT TAB(10);" DATA FOR MONTH";Month;LIN(1)
240 PRINT "SALESPEOPLE";TAB(17);"DISTRICT";TAB(33);"SALES";LIN(1)
250

```

```

260     Month=6                ! Select month of sales survey.
270                                     !
280     FOR I=1 TO 9
290         READ Sales(Month,1,I),Sales(Month,2,I) ! Input district
300     NEXT I                    ! and sales volume.
310                                     !
320     MAT READ Salespeople$     ! Input salespeople.
330                                     !
340     FOR I=1 TO 9
350         PRINT Salespeople$(I);TAB(19);Sales(Month,1,I);TAB(32);
360         PRINT Sales(Month,2,I)
370     NEXT I                    ! Print the unsorted sales data.
380                                     !
390     WAIT 5000                ! Wait 5 seconds before continuing.
400     PRINT PAGE               ! Clear the screen.
410                                     !
420     PRINT TAB(10);" RESULTS FOR MONTH";Month;LIN(1) ! Print
430                                     ! a new heading.
440     MAT SORT Sales(Month,1,*), (Month,2,*) DES TO Rank ! Sort
450                                     ! the sales by district; use
460                                     ! sales volume as a secondary key.
470                                     !
480 ! *****
490 ! *** Establish FOR/NEXT loop for printing sales results. ***
500 ! *****
510                                     !
520     FOR I=1 TO 9
530         IF District=Sales(Month,1,Rank(I)) THEN Skip ! Test: Has
540                                     ! district changed?
550                                     ! If so, skip print routine.
560                                     ! If not, continue.
570         District=Sales(Month,1,Rank(I)) ! Set District to new value.
580                                     !
590         PRINT TAB(10);" DISTRICT";District;"RESULTS" ! Print heading.
600 Skip: PRINT Salespeople$(Rank(I));TAB(30);" ";Sales(Month,2,Rank(I))
610     NEXT I                    ! Go on to next salesperson.
620     END

```

The program first prints the unsorted list. The sort is performed by district using sales volume as the secondary key specifier. If the program is run, the results shown next are obtained.

## DATA FOR MONTH 0

SALESPEOPLE	DISTRICT	SALES
JILL TANDY	3	3000
DON DEEDS	1	6000
RICK HILL	2	1900
SAM SPADE	1	1200
NICK DANGER	3	1400
JOE JONES	1	2500
HARRY WHITE	3	900
BARB SMITH	2	2200
H.J. STEED	2	1700

```

RESULTS FOR MONTH 6
DISTRICT 1 RESULTS
DON DEEDS          6000
JOE JONES          2500
SAM SPADE          1200
DISTRICT 2 RESULTS
BARB SMITH         2200
RICK HILL          1900
H.J. STEED         1700
DISTRICT 3 RESULTS
JILL TANDY         3000
NICK DANGER        1400
HARRY WHITE        900

```

## Sorting String Data

The sorting process for string data follows that for numeric data except that ordering is lexical rather than numeric.

Syntax:

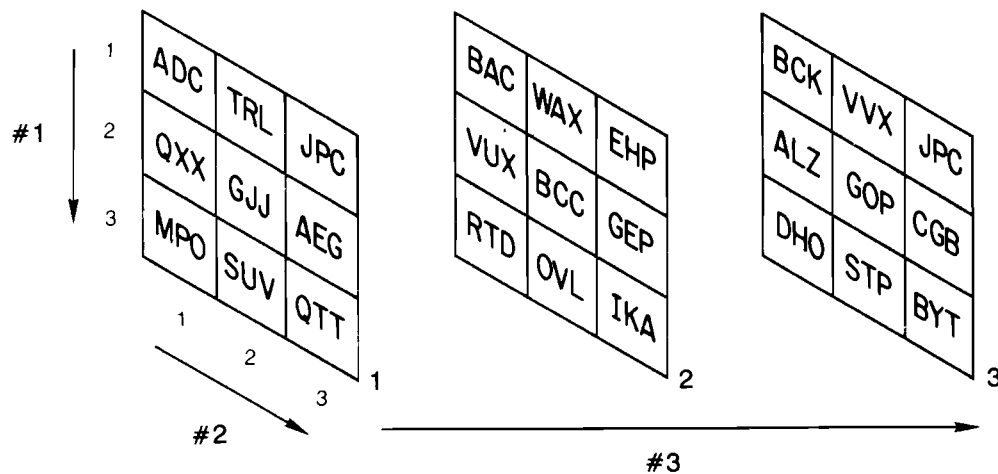
```
MAT SORT string source array (key specifier) [TO pointer array]
```

The source array in this case contains string data. The key specifier specifies which locations are designated as keys. The optional pointer array is a one-dimensional numeric array which contains the order in which the records should be sorted. As with numeric arrays, the source array is not actually rearranged when a pointer array is used.

String sorting uses the LEX function on the string data to perform the ordering (see Chapter 3). The default order is lexical (ascending), but DES can be used to specify reverse lexical (descending) order. The order can also be determined by a user-defined table as described in the LEXICAL ORDER IS section.

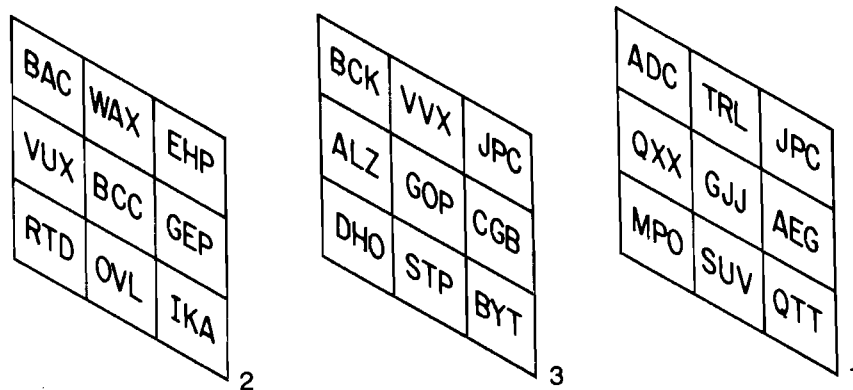
For example, assume array A\$ is dimensioned to be a three-dimensional string array containing three characters per element (DIM A\$(3,3,3) [3]). Assume that the middle location in the right-hand column of each record (plane) is designated as a key. If the records are to be sorted along the third dimension in reverse lexical order, the correct sorting statement is:

```
MAT SORT A$(2,3,*) DES
```



ARRAY A\$ WITH KEYS

This illustration represents the source array. The sorting process sorts the records according to the ASCII values of the data in the key locations. The sorting process begins comparing key characters until dissimilar characters are found. Since the first characters of all the data in the keys are different, sorting can be performed without comparing further characters. When the sort is complete, the array is rearranged as shown.



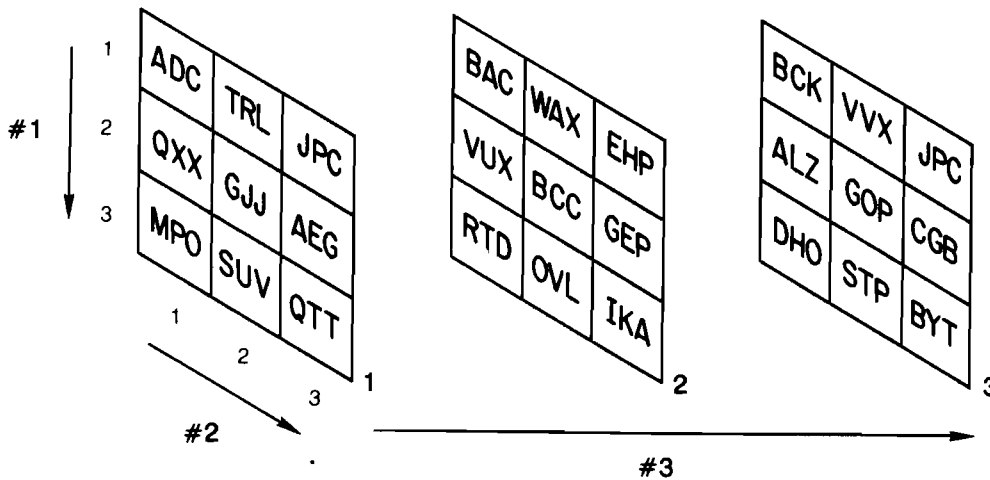
Note that the string data within each memory location has not been rearranged. The records, however, have been rearranged to reflect the new order of the keys of string data. It is important to note that strings of unequal length present no problem. Any unused dimensioned character spaces are filled with the null string and are sorted accordingly. (Example: AB precedes ABC.)



As with numeric sorting, a secondary key specifier can be used to order records which contain identical keys. In the previous array example, assume that the characters in the upper right locations of the records are designated as keys and that the records are to be sorted along the third dimension in lexical order. The sorting process begins comparing characters until dissimilar characters are found. After comparing all the characters in the primary keys of records one and three, the sort recognizes identical data. The sort process could then utilize a secondary key specifier to perform the sort. (All characters must match in order to be considered identical data).

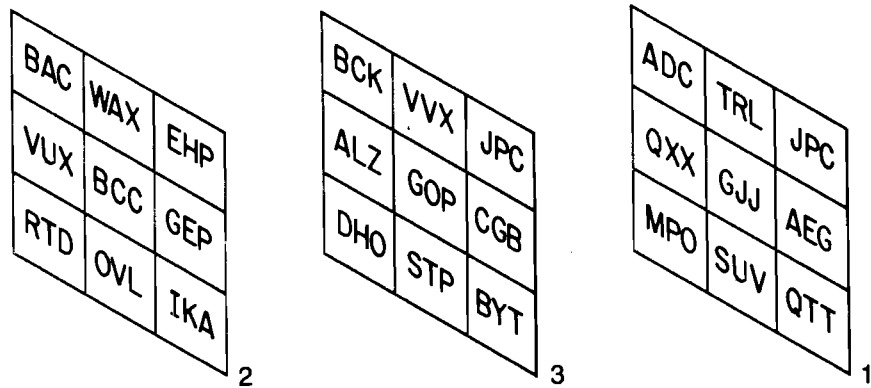
Assume that the data in the lower right locations is described by a secondary key specifier. The correct sorting statement is:

```
MAT SORT A$(1,3,*),(3,3,*)
```



Comparing the data described by the secondary key specifier, the sorting process would order record three before record one since the letter "B" precedes the letter "Q" in the STANDARD lexical order.

Upon execution of the sort statement, the array A\$ is rearranged as shown.



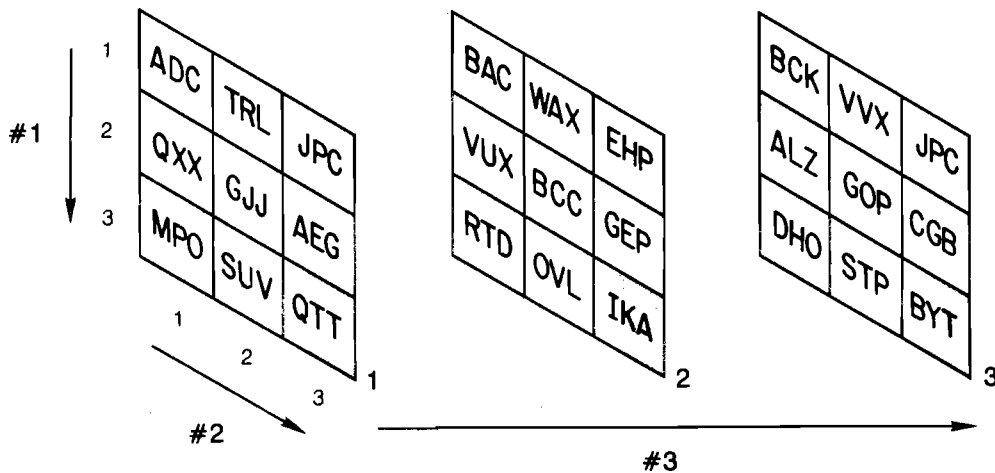
ARRAY A\$ AFTER SORTING

When sorting string data, a substring specifier can be used to describe partial key specifier strings.

Syntax:

```
MAT SORT string source array (key specifier) [ [substring specifier] ]
```

The optional substring specifier indicates which portion of the data string in the specified keys is used to order the records. Referring to the previous example, assume the second and third characters in the strings located in the lower right memory locations are designated as keys.



A\$ KEYS WITH SUBSTRING SPECIFIER

The correct sorting statement is

```
MAT SORT A$(3,3,*) [2,3]
```

where [2,3] specifies that portion of each key string which begins at the second character and ends with the third. In this example, the sorting process begins sorting the records according to the values of the second and third characters in each key string. Sorting can be completed at the second characters since none of them are identical. After execution of the MAT SORT statement, the rearranged record order is 2-1-3.

A substring specifier can also be written in the form [2;3], where the substring described in this case starts at the second character and continues for three characters.

A substring specifier can be used with a secondary key specifier also. The following is a typical entry.

```
MAT SORT A$(1,*,2), (4,*,6) [4;4]
```

The substring specifier ( [4;4] ) indicates that the secondary key specifier (4,\*,6) describes a substring which starts at the fourth character and extends for four characters.

Substring specifiers can be used with primary and secondary key specifiers simultaneously, as can descending specifiers. Example:

```
MAT SORT A$(1,*,3) [4;4] DES, (4,*,6) [3;3] DES
```

When used with multiple key specifiers, substring specifiers need not be of equal length.

A pointer array can also be included:

```
MAT SORT A$(1,*,3) [4;4] DES, (4,*,6) [3;3] DES TO B
```

Note that no punctuation is included between a key specifier and its related specifiers. A comma is included, however, between complete individual specifier entries.

The program shown next demonstrates string sorting. A list of names and phone numbers is read into a string array and then sorted according to last name.

```

10 ! *****
20 ! ***      This program demonstrates the use of      ***
30 ! ***      the MAT SORT statement.                  ***
40 ! *****
50
60      OPTION BASE 1                                ! Select OPTION BASE.
70      DIM Phones$(7,2) [20]                       ! Dimension source array.
80
90      DATA RICHARD HILL,818-8086
100     DATA MICHAEL HILL,916-6502
110     DATA BARB SMITH,917-6802 ! Name and phone
120     DATA HARRY RULE,818-8048 ! number data.
130     DATA JULIE TANDY,818-8090
140     DATA SAMUAL SPADE,818-6809
150     DATA MICHAEL SMITH,916-6800
160
170     MAT READ Phones$ ! Read data into source array.
180
190     MAT SORT Phones$(*,1) [10],(*,1) [1,9] ! Sort data by last name;
200     ! use first name as secondary key.
210
220 ! *****
230 ! ***      Print heading and sorted directory.      ***
240 ! *****
250
260     PRINT TAB(4);"THE SORTED DIRECTORY:";LIN(1) ! Print heading.
270     FOR I=1 TO 7
280         PRINT Phones$(I,1) [10];", ";Phones$(I,1) [1,9];TAB(20);
290         PRINT Phones$(I,2) ! Print the sorted directory.
300     NEXT I
310     END

```

Notice that the data are sorted by last names using the first names as a secondary key. If the program is run, the results shown next are obtained.

#### THE SORTED DIRECTORY:

HILL,MICHAEL	916-6502
HILL,RICHARD	818-8086
RULE,HARRY	818-8048
SMITH,BARB	917-6802
SMITH,MICHAEL	916-6800
SPADE,SAMUAL	818-6809
TANDY,JULIE	818-8090

## MAT REORDER Statement

The MAT REORDER statement is used to order an array according to the contents of an existing pointer array.

Syntax:

```
MAT REORDER object array BY pointer array [, dimension specifier]
```

The object array in this syntax is rearranged in the order specified by the contents of the pointer array. The maximum allowable number of dimensions of the object array is six. The optional dimension specifier selects the dimension of the object array along which records are ordered. The dimension specifier is a number from 1 to 6 or an expression which represents this number. If it is not specified, a value of 1 is assumed by the computer. The dimension specifier describes the object array dimensions in the manner shown next.

```
DIM Object_array(A,B,...,F)
                ↑  ↑  ↑
Dimension Specifier: 1 2 6
```

For example, assume that the array Point is a pointer array containing a range of values determined by a previous sorting process as shown.

4	2	1	3
---	---	---	---

It is dimensioned as follows:

```
OPTION BASE 1
DIM Point (4)
```

Also, assume that array Object is a two-dimensional object array which contains numeric data and is dimensioned as shown.

```
OPTION BASE 1
DIM Object (4,5)
```

The following illustration represents array Object.

1	9	3	8	6	2	A
2	1	4	5	1	1	B
3	7	3	7	2	0	C
4	5	1	8	6	3	D
	1	2	3	4	5	

#2

**ARRAY Object**

If array Object is to be rearranged along the first dimension according to the contents of array Point, the correct reordering statement is

```
MAT REORDER Object BY Point,1
```

where Object is the object array, Point is the pointer array, and 1 specifies that the object array is reordered along its first dimension (i.e., each row is considered a record). Upon execution of the reorder statement, array Object is rearranged as shown next.

5	1	8	6	3	D
1	4	5	1	1	C
9	3	8	6	2	B
7	3	7	2	0	A

} Reordered Records

**ARRAY Object AFTER REORDERING**

The records are rearranged in the order specified by the pointer array. Note that they are not necessarily rearranged in ascending or descending order. There is no necessary relationship among the records. They are merely rearranged according to the pointer array.

The pointer array must be dimensioned the same size as the dimension of the object array along which reordering is performed. In the previous example, the pointer array Point is dimensioned four elements in length as is the first dimension of array Object as shown next.

```
DIM Object(4,5), Point(4)
```

---

**NOTE**

If the pointer array contains duplicate numbers, unpredictable results occur. Also, if the pointer array contains numbers which are out of range for the specified dimension of the object array, an error message results when reordering is attempted.

---

As you may recall, a MAT SORT statement which contains a pointer array does not rearrange the source array upon execution. It merely fills the pointer array with the order of the records as if they were rearranged. The MAT REORDER statement can be used to rearrange that source array at a later time. For example,

```
MAT SORT A(1,*,3) TO B
```

fills the pointer array B with the proper sorted order of records, but it does not actually rearrange the source array A. In order to rearrange the source array, a MAT REORDER statement can be used. Example:

```
100 MAT SORT A(1,*,3) TO B
```

```
200 MAT REORDER A BY B, 2
```

The execution of line 100 fills the pointer array but does not rearrange the source array A. Line 200, however, does rearrange the source array according to the order defined earlier by the MAT SORT statement.

A program is now presented that demonstrates a possible use of the MAT SORT and MAT REORDER statements. Original\$ is a one-dimensional array list of names and grade point averages.

The program fills the array list and sorts the data both by name and grade point average.

```

10 ! *****
20 ! ***          This program demonstrates the use          ***
30 ! ***          of the MAT SORT statement.                ***
40 ! *****
50
60     OPTION BASE 1                                     ! Select OPTION BASE.
70     DIM Original$(6)[21],B(6),C(6),D(6) ! Dimension source
80                                     ! and pointer arrays.
90     MAT READ Original$                               ! Read data into source array;
100                                     ! data in Lines 330-380.
101
110     PRINT "THE ORIGINAL SEQUENCE IS:",LIN(1)
120     PRINT Original$(*)                             ! Print the unsorted list.
130
140     MAT SORT Original$(*)[1,16] TO B               ! Sort list by last names.
150     MAT SORT Original$(*)[18,21] DES, (*)[1,16] TO D ! Sort list by
160                                     ! grade points; use names
170                                     ! as secondary keys.
180     FOR I=1 TO 6
190         C(B(I))=I                                 ! Compute the inverse permutation.
200     NEXT I
210
220     MAT REORDER Original$ BY B                       ! Reorder list by name.
230     PRINT "REORDERED BY NAME:",LIN(1)
240     PRINT Original$(*)                             ! Print the reordered list.
250
260     MAT REORDER C BY D                               ! Combine permutations C and D
270                                     ! into C.
280     MAT REORDER Original$ BY C                       ! Reorder list by grade point
290                                     ! using the combined permutation.
300     PRINT "REORDERED BY GPA:",LIN(1)
310     PRINT Original$(*)                             ! Print reordered list.
320
330     DATA SMITH BILL H           3.75
340     DATA JONES BOB R            2.00
350     DATA BROWN MARY A           2.00
360     DATA SMITH GLEN C           2.90
370     DATA TAYLOR RALPH E         3.50
380     DATA JONES BONNIE R         3.50
390
400     END

```

Notice that the inverse permutation of pointer array B is computed. The use of this permutation allows two sorts to be performed on the original object array thereby eliminating the need for a duplicate array. If the program is run, the results shown next are obtained.



THE ORIGINAL SEQUENCE IS:

SMITH BILL H	3.75
JONES BOB R	2.00
BROWN MARY A	2.00
SMITH GLEN C	2.90
TAYLOR RALPH E	3.50
JONES BONNIE R	3.50

REORDERED BY NAME:

BROWN MARY A	2.00
JONES BOB R	2.00
JONES BONNIE R	3.50
SMITH BILL H	3.75
SMITH GLEN C	2.90
TAYLOR RALPH E	3.50

REORDERED BY GPA:

SMITH BILL H	3.75
JONES BONNIE R	3.50
TAYLOR RALPH E	3.50
SMITH GLEN C	2.90
BROWN MARY A	2.00
JONES BOB R	2.00

It is important to note that several arrays can be reordered by the same pointer array. The program shown next utilizes this feature. Several "parallel" arrays are reordered according to the contents of a single pointer array.

```

10 | *****
20 | ***      This program demonstrates the use      ***
30 | ***      of the MAT REORDER statement.        ***
40 | *****
50 |
60 | OPTION BASE 1
70 | DIM Id_no(4),Jobs$(4),Clock_no(4),Order(4) | Dimension source and
80 |                                     | pointer arrays.
90 | MAT READ Id_no | Read appropriate data into arrays.
100 | MAT READ Jobs$ | Data is in lines 230,240,250.
110 | MAT READ Clock_no
120 |
130 | *****
140 | ***      Print the heading and reorder the arrays.      ***
150 | *****
160 |
170 | PRINT TAB(5);"ID NUMBER";TAB(20);"JOB";TAB(30);"CLOCK NUMBER";LIN(2);
180 | CALL Print(Id_no(*),Jobs$(*),Clock_no(*)) | Call the Print subroutine
190 |                                     | using the array data as parameters.
200 | PRINT LIN(1)
210 |

```

```

220     MAT SORT Id_no(*) TO Order ! Sort ID numbers to pointer array.
230     !
240     MAT REORDER Jobs# BY Order
250     MAT REORDER Id_no BY Order ! Reorder all arrays by the same pointer.
260     MAT REORDER Clock_no BY Order
270     !
280     PRINT " ----- ";LIN(1)
290     CALL Print(Id_no(*),Jobs#(*),Clock_no(*)) ! Call subroutine
300     ! for printing the reordered arrays.
310     DATA 2325,4157,9020,1025
320     DATA MANAGER,GUARD,TYPIST,NURSE ! Data is entered in the proper order
330     DATA 307,118,476,534 ! to create "parallel" arrays.
340     !
350     END
360     !
370     ! *****
380     ! *** This subroutine prints data in the specified order. ***
390     ! *****
400     !
410     SUB Print(Id_no(*),Jobs#(*),Clock_no(*))
420     FOR I=1 TO 4
430     PRINT TAB(6);Id_no(I);TAB(19);Jobs#(I);TAB(33);Clock_no(I);
440     NEXT I
450     SUBEND

```

Notice that the individual object arrays are reordered by the same pointer array. If the program is run, the results shown next are obtained.

ID NUMBER	JOB	CLOCK NUMBER
2325	MANAGER	307
4157	GUARD	118
9020	TYPIST	476
1025	NURSE	534
-----		
1025	NURSE	534
2325	MANAGER	307
4157	GUARD	118
9020	TYPIST	476



## MAT SEARCH Statement

### Searching Numeric Arrays

The purpose of the MAT SEARCH statement is to provide information about user-defined conditions within an array. This information is returned to a variable for recall and examination.

Syntax:

```
MAT SEARCH source array (location specifier) , condition ; variable
[ , starting address]
```

The source array represents the array in which searching is performed. The location specifier defines the locations within the source array which are searched. No location specifier need be included if the source array is one-dimensional. The condition is that value or location for which you are searching. When the condition is satisfied, the memory content or location which satisfies it is returned to the variable. The optional starting address specifier allows you to select a location within the range defined by the location specifier at which you want searching to begin. Numeric searching comparisons are performed in the INTEGER mode for integer-precision arrays and in the REAL mode for short or real-precision arrays.

The conditions available in the MAT SEARCH process are entered in the following form:

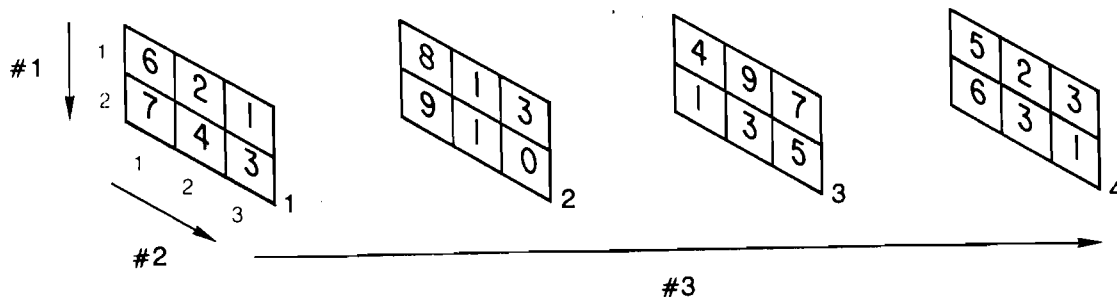
```
LOC (relational operator - expression)
#LOC (relational operator - expression)
MAX
MIN
```

The list of relational operators includes: >, <, =, #, >=, <=, <>. For an integer search, the expression is rounded before the comparison is performed. A LOC condition causes the search process to scan the specified locations until it finds the first value which satisfies the condition. The default value for a relational operator is =.

A #LOC condition causes the search process to scan the specified locations and return the total number of locations whose contents satisfy the condition.

MAX and MIN conditions do not include an argument in parentheses because they automatically cause the search process to scan all specified locations to find the maximum and minimum values. To return the result of a MAX or MIN search, a string variable is used with string data, and a numeric variable with numeric data.

To demonstrate the MAT SEARCH process, assume that array A is a three dimensional numeric array containing random numbers as shown below.



ARRAY A WITH SEARCH LOCATIONS

Assume that the shaded locations are to be searched until one is found which contains a value greater than 5. Let B represent the variable to which the result is returned.

The shaded memory locations must be described by a location specifier. Since they lie along the third dimension, the records in which they lie are defined accordingly. The correct location specifier is:

(1,3,\*)

Note that the location specifier is written in the same format as a key specifier for a MAT SORT statement. The subscripts are written in the same order as the array dimensions (the numbered arrows). The first two subscripts define the upper right location, and the asterisk indicates that the third subscript is varied over its range of values to describe all three records.

The correct search statement for this example is

```
MAT SEARCH A(1,3,*),LOC(>5);B
```

where A is the source array, (1,3,\*) defines the locations to be searched, LOC(>5) is the condition, and B is the variable to which the result is returned.

After execution of the search statement, the number 3 is returned to the variable B. This is the first record whose specified location satisfies the condition (i.e., it contains a value greater than 5). Searching begins at the location in that record described by the smallest number in the range of varied subscripts and continues to the largest.

If a condition is not satisfied upon completion of the searching process, a value one greater than the upper limit of the varied subscript is returned to the numeric variable. For example, if the specified locations in the previous example are searched for a value greater than 8, none is found. Therefore, a value of 5 representing a number one greater than the record containing the last searched location is returned to B. Note that if the upper limit of the varied subscript is 32 767, the value returned by an unsuccessful search is - 32 768.

The location specifier initially defines the range of locations to be searched. If you do not wish to search the entire range, a starting address specifier can be used to designate where the search is to begin. In the previous example, assume that the search process is to scan only the last three records. The correct search statement is

```
MAT SEARCH A(1,3,*) ,LOC(>5);B,2
```

where the number 2 directs the search to begin at the specified location in the second record and proceed to the last record. Upon execution, the number 3 indicating the third record is returned to the variable B because the content of its specified location is the first to satisfy the condition.

The following program is included to demonstrate a search for maximum and minimum values and number of occurrences. The array Numbers is filled with random numbers from which the maximum and minimum values are selected.

```
MAXIMUM = 9 MINIMUM = 1
```

If this program is run, the following results are obtained.

```
10  ! *****
20  ! ***      This program demonstrates the use of the      ***
30  ! ***      MAT SEARCH statement to find maximum and    ***
40  ! ***      minimum values within an array.            ***
50  ! *****
60
70  OPTION BASE 1
80  DIM Numbers(11)
90
100 FOR I=1 TO 11
110     DISP "NUMBER";I;
120     INPUT Numbers(I)
130     NEXT I
140
150 ! *****
160 ! *** The data used in this example: 6,1,9,2,8,3,8,9,1,7,5 ***
170 ! *****
180
190     MAT SEARCH Numbers(*),MAX;Max
200     MAT SEARCH Numbers(*),MIN;Min
210
220     PRINT "MAXIMUM =" ;Max, "MINIMUM =" ;Min
230     END
```



```

130 ! *****
140 ! *** Search for locations of numbers greater than 2. ***
150 ! *****
160 !
170 Search: MAT SEARCH B(1,2,*),LOC(>2);C,C+1
180 !
190     IF C+1>6 THEN 260     ! Test: Have all locations been searched?
200 !                         ! If yes, stop.
210     DISP C;              ! Display most recent search results.
220     WAIT 500             ! Wait slightly before further display.
230     IF C<5 THEN Search   ! Test: Has search reached last location?
240 !                         ! If not, continue.
250     DATA 7,4,0,1,6,2,3,4,5,6,2,1,8,7,0,9,6,3,0,4
260     END

```

The variable *C* is initially set to zero. Line 90 begins the search routine. Since *C* was set to zero, the starting address specifier (*C* + 1) directs the search to begin at record one. Line 110 tests to see if the specified locations have all been searched. (Remember, if all locations have been searched, a number one greater than the last record searched is returned to the variable. In this case, that number is six). If all locations have not yet been searched, line 120 displays the contents of *C* (it now contains a satisfactory value). If *C* contains a value less than 5, the search is not finished and line 140 directs the program to search again. Due to the nature of the starting address specifier, the search begins this time at the next memory location beyond that which satisfied the condition previously. In other words, the search resumes where it left off. If you run this program, the following should be displayed:

```

2 3 4 5

```

## Searching String Arrays

Searching string arrays follows the same format as searching numeric arrays.

Syntax:

```

MAT SEARCH source array (location specifier), condition; variable
      [, starting address]

```

The source array, location specifier, condition, variable, and starting address specifier serve the same functions for string arrays as for numeric arrays. If a relational operator is included, the LEXICAL ORDER IS statement determines the ordering of the string data. The order can be STANDARD (according to the ASCII table) or user-defined (see Appendix B).

For example, assume array List\$ contains a list of names and dollar amounts. The program shown next inputs the data to the source array. It then searches for a particular name and outputs the corresponding dollar amount.

```

10 | *****
20 | ***      This program demonstrates the use of      ***
30 | ***      the MAT SEARCH statement.                ***
40 | *****
50 |
60 | OPTION BASE 1      | Select OPTION BASE.
70 | DIM List$(4)(20)  | Dimension source array.
80 |
90 |   FOR I=1 TO 4    | Establish loop for data entry.
100 |     READ List$(I) | Read data into List$.
110 |     NEXT I        | Read next data item.
120 |
130 |   MAT PRINT List$ | Output the original list.
140 |
150 | *****
160 | ***      Search List$ for a particular name.      ***
170 | *****
180 |
190 |   MAT SEARCH List$(*)[1,5],LOC("HAYS ");B          | Search the proper
200 |                                                     | portion of each
210 |                                                     | string.
220 |
230 |   PRINT List$(B)[1,5];": ";List$(B)[15,18]        | Output specified
240 |                                                     | name and dollar
250 |                                                     | amount
260 |
270 |   DATA BROWN FRANK $100 ,SMITH BOB $75
280 |   DATA HAYS SUE $150 ,JONES ED $200
290 | END

```



In this program a MAT SEARCH is used to find the string which contains the required name. Once that string is found, the portion of it containing the dollar amount is displayed. Note that the substring specifier is used in the search and display statements. If you run this program, the following results are obtained.

\$150

```

BROWN FRANK $100
SMITH BOB $75
HAYS SUE $150
JONES ED $200
HAYS : $150

```

MAX and MIN values can also be obtained from a string search as demonstrated by the program shown next.

```

10 ! *****
20 ! *** This program demonstrates use of the MAT SEARCH ***
30 ! *** statement in conjunction with strings. ***
40 ! *****
50
60 OPTION BASE 1 | Select the OPTION BASE.
70 DIM Stringsearch$(3)(3) | Dimension the string array.
80
90 MAT READ Stringsearch$ | Read data into the array.
100 DATA "ABC", "BCD", "DEF"
110
120 MAT SEARCH Stringsearch$(*),MAX;A$
130 | Search string array for
140 | maximum string value.
150
160 MAT SEARCH Stringsearch$(*),MIN;B$
170 | Search string array for
180 | minimum string value.
190
200 PRINT "MAXIMUM = ";A$,"MINIMUM = ";B$
210 | Output the results.
220 END

```

The array Stringsearch\$ is filled with string data. MAT SEARCH statements are then used to find the maximum and minimum values of the data. Note that the search statements use the current lexical order of the data to perform the search comparisons (see LEXICAL ORDER IS statement). If this program is run, the results shown next are obtained.

```

MAXIMUM = DEF MINIMUM = ABC

```

# Extended Character Sets

- page 38 • **LEXICAL ORDER IS** (allows you to select the collating sequence for string sorts and string comparisons)
- page 38 • **LEXICAL ORDER IS STANDARD** (selects ASCII collating sequence)
- page 39 • **LEX** (allows string comparisons)
- page 40 • **LWC\$** (returns a string with all uppercase letters converted to lowercase)
- page 40 • **UPC\$** (returns a string with all lowercase letters to uppercase)

## Statement Syntax

**LEXICAL ORDER IS** lexical order designator

lexical order designator: an integer array containing a local language collating table or the word **STANDARD**

**LEX** (string expression 1, string expression 2)

Dot Matrix - All items in dot matrix must be entered as shown.

[ ] - All items in brackets are optional.

... - Three dots indicate that successive parameters are allowed, when each is separated by a comma.

# Chapter 3

## Introduction

The statements described in this chapter are used as programming aids to define and access comparison and case functions on the ASCII and Roman Extension Character Sets. Each of these statements or functions can be executed from within a program or from the keyboard. Examples of each statement are included to demonstrate its use.

### LEXICAL ORDER IS Statement

The LEXICAL ORDER IS statement allows you to select the collating sequence for string sorts and determine the results of string comparisons. It also determines the results of UPC\$ and LWC\$ transformations for Roman Extension characters.

Syntax:

```
LEXICAL ORDER IS lexical order designator
```

The lexical order designator must be an integer array containing a local language collating table, or the word STANDARD which indicates the standard ASCII character sequence. The following local language collating tables are available on the Lexical Tables Cartridge for use with the Advanced Programming ROM.

```
FRENCH   SWEDSH (Swedish)  
GERMAN   SPANSH (Spanish)
```

Refer to Appendix B for information on the lexical order collating tables.

The STANDARD lexical order designator selects the ASCII table as the collating sequence (i.e., string sorting is performed according to ASCII values). STANDARD is the default lexical order designator at power ON and after SCRATCH A. As such, you need not include a LEXICAL ORDER IS statement if you wish to use ASCII as the collating sequence. RESET does not clear the current lexical order. However, if you wish to change to STANDARD from some other collating sequence, you can do so by entering:

```
LEXICAL ORDER IS STANDARD
```

The program shown next demonstrates how a local language collating sequence is selected.

```
10 INTEGER A(400)
20 ASSIGN #1 TO "FRENCH"
30 MAT READ #1;A
40 LEXICAL ORDER IS A(*)
```

Line 10 dimensions an integer array to hold the specified collating table. Here, A is chosen as the array and is dimensioned to a length of 400 which accommodates any of the local language character collating tables listed previously.

Line 20 assigns #1 to the specified file on the cartridge; here, FRENCH.

Line 30 reads the contents of the file (the specified collating table) into the array.

Finally, line 40 establishes the contents of the array (which now contains the collating table) as the lexical order.

Note that when a LEXICAL ORDER IS statement is executed, the information in the array is transferred into an internal buffer. This allows you to use the array for something else. Care must be taken, however, to insure that enough memory exists for the buffer when the statement is executed. Otherwise, an error occurs.

## The LEX Function

The LEX function allows you to compare two strings and determine which precedes the other according to the current lexical order. The LEX function can be executed from the keyboard or from within a program.

Syntax:

```
LEX (string expression 1, string expression 2)
```

The LEX function begins comparing string characters until a difference is found or until one or both strings terminates. The function returns the number -1 if string 1 precedes string 2. If the current lexical order cannot distinguish between the two strings, 0 is returned. The number 1 is returned if string 2 precedes string 1.

The program shown next demonstrates the LEX function. Two strings are compared and the result is displayed.

```
10 A=LEX("XEF","XBC")
20 DISP A
30 END
```

Notice that the program utilizes the default STANDARD lexical order. The LEX function compares the strings up to their second characters since these are the first respective characters that differ. Upon execution, the number 1 is displayed because B (string 2) precedes E (string 1) according to the current lexical order (ASCII).

The previous example program compared strings of equal length. The example program shown next demonstrates a comparison of strings of unequal length.

```
10 A=LEX("A","A ")
20 B=LEX("A","AB")
30 DISP A,B
40 END
```

In this case, both A and B are -1. In both comparisons, the first string terminates before the second.

## Uppercase and Lowercase Functions

The Advanced Programming ROM extends the capabilities of the mainframe uppercase (UPC\$) and lowercase (LWC\$) functions to correctly handle the upper and lowercase transformations of the various local language characters.

The results of the UPC\$ and LWC\$ functions are determined by the entries in the upper and lowercase sections of the lexical order array for Roman Extension characters (refer to Appendix B). The upper and lower cases of characters in the main ASCII set are fixed.

Descriptions of the French, German, Spanish and Swedish upper and lowercase transformations which are provided in the respective files on the cartridge are explained in Appendix B.

The UPC\$ function can be used to obtain a proper LEX comparison if it is not known if the two strings are in the same case (i.e., if "ABC" should equal "abc"). Example:

```
X = LEX (UPC$ (A$), UPC$ (B$) )
```

---

#### NOTE

If a program containing UPC\$ or LWC\$ statements is STORE'd in a System 45B which has an Advanced Programming ROM, and is then LOAD'ed into a System 45B which does not have an Advanced Programming ROM, a MISSING ROM error message occurs.

If a program is STORE'd in a System 45B which has no Advanced Programming ROM, and is then LOAD'ed into a System 45B which has an Advanced Programming ROM, no error message results. However, if any lexical statements are added, the current UPC\$/LWC\$ lines must be re-STORE'd in order for the Advanced Programming ROM to affect them.

It is suggested that you SAVE rather than STORE programs which will be switched between those System 45B Desktop Computers which have an Advanced Programming ROM and those which do not.

---



# Statement Summary

**MAT SORT** source array (key specifier) [ [substring specifier] ] [DES] [, (secondary key specifier) [ [substring specifier] ] [DES] ] [TO pointer array]

Arranges specified numeric or string data in ascending (lexical) or descending (reverse lexical) order. Optional pointer array stores sorting order.

**MAT REORDER** object array BY pointer array [, dimension specifier]

Reorders the object array according to the order specified by the pointer array. Optional dimension specifier selects the object array dimension along which reordering is performed.

**MAT SEARCH** source array (location specifier), condition ; variable [, starting address]

Searches the specified condition of the source array for the specified condition. The satisfactory condition is returned to the variable. The optional starting address specifier selects the location at which the search begins.

**LEXICAL ORDER IS** lexical order designator

Establishes the particular collating table by which sorting and searching is performed.

**LEX**(string expression 1, string expression 2)

Compares two strings and determines which precedes the other according to the current lexical order.

**UPC#** and **LWC#**

Extends the capabilities of the uppercase and lowercase transformations to the various local language characters.

# Appendix A





## User-Defined Lexical Order

It is recommended that you use the existing collating tables as they appear on the tape cartridge. However, an experienced programmer can create a lexical order table to accommodate his particular application. The following information describes the procedures involved. A knowledge of binary and octal numbers is necessary to utilize these procedures.

To modify a local language collating table, it is first input to an integer array. Once in the array, the specified table can be accessed for modifications.

Along with the local language tables, a special ASCII table ("ASCII") is included on the cartridge for the purpose of modifying the ASCII collating sequence. The LEXICAL ORDER IS STANDARD statement automatically uses the standard ASCII sequence. But in order to modify it, the "ASCII" file on the cartridge may be input to an integer array in the same manner as is a local language table. Example:

```
10 | *****
20 | *** This program segment illustrates the ***
30 | *** LEXICAL ORDER IS statement. ***
40 | *****
50 |
60 | OPTION BASE 1 | Select OPTION BASE.
70 | INTEGER B(400) | Dimension integer array.
80 | ASSIGN #1 TO "ASCII" | Select ASCII file.
90 | MAT READ #1;B | Read in the ASCII table.
100 |
110 | | User modifications.
120 |
130 | LEXICAL ORDER IS B(*) | Establish contents of
140 | | array B as lexical order.
150 | END
```

Line 40 inputs the "ASCII" file contents to the integer array B. Lines 50 through 70 represent the user modifications to the ASCII table, and line 80 establishes the modified table as the lexical order. Notice that in order for the table in the given array to dictate lexical comparison results, a LEXICAL ORDER IS statement must be included after any modifications to the array have been made.

# Appendix B

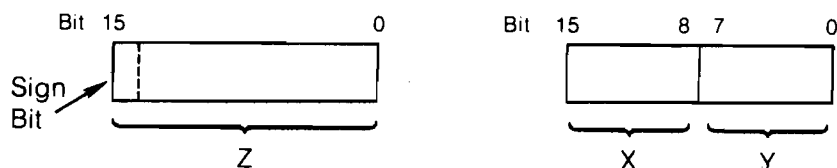
Once a LEXICAL ORDER IS statement has been executed, the integer array can be used for other purposes. Be sure, however, to save all necessary modifications on the tape cartridge before using the array for other purposes. The program lines shown next can be used to maintain a copy of the lexical table modifications for the previous example.

```
100 ASSIGN #3 TO "NEW"
110 MAT PRINT #3;B
```

The lexical order tables give you the capability to define the collating sequence (collating section), define the uppercase/lowercase transformation of Roman Extension characters (UPC/LWC section), and prescribe "special handling" of certain characters (mode section). The lexical table with its individual sections is an integer array organized as shown in the following illustration.

1	Total Length
2	Mode Section Length
3 . . . . 258	Collating Section (Length = 256)
259 . . 354	UPC/LWC Section (Length = 96)
355 . . . . 354 + N	Mode Section (Length = N)

Each element of the lexical table consists of 16 bits. This arrangement can accommodate one number or an integer combination of two (unsigned) numbers as shown next.



If one number is stored, a binary two's complement storage format is used. In this manner, each element has the ability to contain two distinct values.

The first element of the array contains the complete length of the lexical table. This length is 354 plus the number of elements (N) needed in the user-defined mode section (see explanation of mode section).

The second element of the array contains N, the length of the mode section of the table. Once the length of the mode section is defined, the first two elements of the array table can be filled with the proper information.

Elements 3 – 258 of the table contain the actual collating sequence and the appropriate pointers into the mode section, if one exists.

Elements 259 – 354 define the uppercase/lowercase transformations for characters in the Roman Extension Set (refer to the ASCII and Roman Extension chart at the end of this appendix).

The mode section consists of elements 355 to (354 + N) and provides facilities for handling certain special case characters.

## Collating Sequences

When the lexical order is STANDARD, the ASCII codes define the collating sequence. Thus, A lexically precedes B (i.e.,  $\text{LEX}(\text{"A"}, \text{"B"})$  is  $-1$ ) because the ASCII code for A is 65, the ASCII code for B is 66, and 65 is less than 66.

For computer processing, each character has been assigned an ASCII code. The ASCII codes are fixed, however, and cannot be changed to reflect a new user-defined lexical order. The lexical order tables give you the capability of redefining the lexical order by assigning each character a sequence number. Characters are still represented internally by their ASCII codes, but a lexical comparison compares sequence numbers rather than ASCII codes. Now,  $\text{LEX}(\text{"A"}, \text{"B"})$  is  $-1$  if the sequence number of A is less than the sequence number of B. For example, the strings "ABC" and "XYZ" are stored in the computer as strings of ASCII codes:

ASCII code for A	ASCII code for B	ASCII code for C
---------------------	---------------------	---------------------

ASCII code for X	ASCII code for Y	ASCII code for Z
---------------------	---------------------	---------------------

The calculation of LEX ("ABC", "XYZ") involves a comparison of sequence numbers:

Sequence # for A	Sequence # for X
Sequence # for B	Sequence # for Y
Sequence # for C	Sequence # for Z

By properly assigning each character a sequence number, you can define any lexical ordering of the characters.

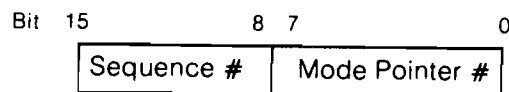
### Collating Section

The collating section of the lexical order table (elements 3-258) defines the actual collating sequence by assigning a sequence number to each character. There are 256 possible 8-bit character codes (refer to the ASCII and Roman Extension chart). The user-defined sequence number for the character with ASCII code 0 goes in element 3 of the lexical table (element 1 of the collating section), sequence number for character with ASCII code 1 goes in element 4 of the lexical table (element 2 of the collating section), and so on.

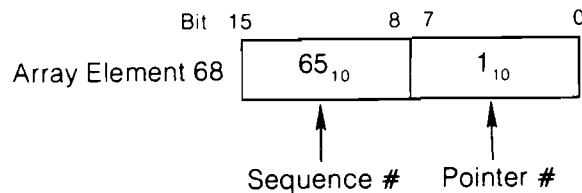
For example, the letter A has ASCII code 65<sub>10</sub>. When an A is encountered, the manipulation routine goes to the 66th position of the collating section to fetch a sequence number. Note that the sequence number is fetched from the collating section element number equal to the ASCII value of the character + 1. This is because the first character in the collating section is null, and its ASCII value is zero. Usually, the sequence number equals the ASCII value. However, if a different collating order is selected, the sequence number may differ from the ASCII value.

In addition to the sequence number, each entry of the collating section may contain a pointer into the mode section of the lexical table. If no pointer is specified, 0 is entered. The mode section contains instructions governing special case letters. The pointer number indicates in which position of the mode section a special case is handled.

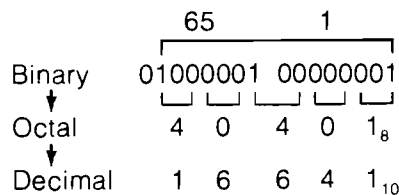
The sequence number and mode pointer number for each character are combined to form a single integer entry in the lexical table:



Assume that A is a special case letter with the sequence number 65. Also, assume that the special case is handled in position 1 of the mode section. The collating section entry containing this information is represented by:



The actual content of this entry is 16 641<sub>10</sub>. The number is obtained as shown next.

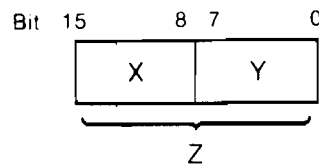


Both the sequence number (65) and the pointer number (1) are combined to form one integer value within the table element, in this case, 16 641<sub>10</sub>.

An alternate method of converting the binary and octal elements of the lexical table entries to integer form is given below. For example, if the sequence number you are using is greater than 127, you may find it more convenient to use the formulas shown next.

X and Y are each 8-bit unsigned integers.

Z is a 16-bit signed integer formed by concatenating the bit patterns of X and Y.



1. Given X and Y, calculate Z:  

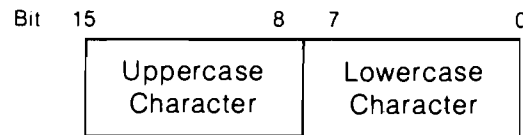
$$Z = X \cdot 256 + Y - (X > 127) \cdot 65536$$
2. Given Z, calculate X and Y:  

$$X = \text{INT}(Z / 256) + (Z < 0) \cdot 256$$

$$Y = Z \text{ MOD } 256$$

## Uppercase/Lowercase Section

Elements 259-354 of the lexical table are used for uppercase/lowercase information for characters residing in the Roman Extension Set. Data contained in these elements is in the form:



The UPC/LWC section contains an element for each character with an 8-bit code of 160 through 255. The first element of the UPC/LWC section (element 259 of the lexical table) contains the upper/lowercase for the character with 8-bit code 160, the next element contains the transformations for the character with 8-bit code 161, and so on.

The purpose of this section of the lexical table is to allow you to define the upper and lowercase versions of each letter in the Roman Extension Set for use by the UPC\$ and LWC\$ functions. All characters in the main ASCII table have fixed upper and lowercases, and as such, are not changed. If both the uppercase and lowercase 8-bit codes in the table are 0 for a given character, the character is unchanged by the UPC\$ and LWC\$ functions.

For example, if a word containing a lowercase *é* (code 197<sub>10</sub>) is to be converted by French conventions, the accent is omitted when the letter is converted to uppercase. That is, *é* → E. But under German conventions, an umlaute in the letter remains so that *ä* → Ä. Therefore, for French, the 38th entry of the uppercase/lowercase section (296th entry of the table) is

	E	é
Array Element 296	69 <sub>10</sub>	197 <sub>10</sub>

while the 45th entry in the German uppercase/lowercase section is

	Ä	ä
Array Element 303	216 <sub>10</sub>	204 <sub>10</sub>

## Mode Section

The mode section (starting at element 355) of the lexical table handles special case characters. The three special cases shown next are possible.

- Accent Priority
- 1 For 2 Character Replacement
- 2 For 1 Character Replacement

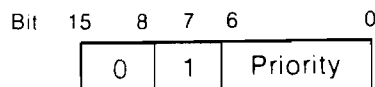
### Accent Priority

You may specify the priority of letters with accent marks. Example:

$e < \hat{e} < \ddot{e} < \acute{e} < \grave{e}$

The symbol (<) means “precedes alphabetically”. In this case, these letters have the same sequence number in the collating section, but different pointer numbers. That is, each has an individual entry in the mode section.

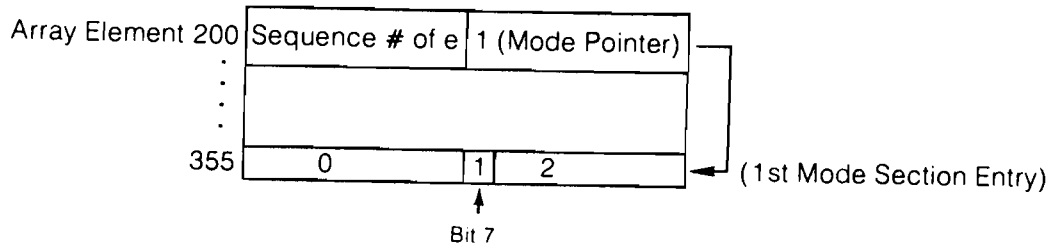
An accent priority requires only one entry in the mode section (some special cases require more than one). The format of the mode section entry for this special case is:



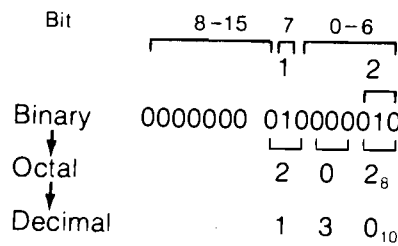
The zero indicates that the eight bits of this special case mode entry are not used, and must contain zero (0) so as not to be confused with one of the other special cases. A 1 must be present in bit seven to indicate that this special case is an accent priority. The remaining bits describe the user-defined priority number.



For example, assume that the letter  $\acute{e}$  in the collating section has a pointer number which specifies the first entry of the mode section (or element 355 of the table). Also, assume that this mode entry assigns a priority of 2 to the letter  $\acute{e}$ . The collating and mode sections contain the information shown next.



The following scheme is used to find the correct mode section entry.



130<sub>10</sub> is the proper entry for the mode section element. Bits 6 through 0 are used to describe the priority number. Bit 7 is always 1 for an accent priority special case. Bits 15 through 8 are not used, but must contain zeros to insure proper results.

Note that the value of an accent priority entry in the mode section is always 128 plus the assigned priority. This means that the maximum accent priority is 127.

Accent priority mode section entries assign a relative priority to characters having the same sequence number. The accent priority is used only to distinguish two otherwise identical strings which differ only in their accent marks.

For example, assume  $\acute{e}$  and  $\grave{e}$  have the same sequence number and each has been assigned an accent priority;  $\acute{e}$  has priority 5, and  $\grave{e}$  has priority 2. The routine to calculate LEX (" $\acute{e}$ ", " $\grave{e}$ ") first compares the sequence numbers of  $\acute{e}$  and  $\grave{e}$ .

Since the sequence numbers are the same (and each character has been assigned an accent priority), the accent priorities are next compared. Five is greater than 2, so LEX (" $\acute{e}$ ", " $\grave{e}$ ") is 1 (i.e.,  $\acute{e} > \grave{e}$ , or  $\acute{e}$  comes after  $\grave{e}$  alphabetically).

A character which is not assigned an accent priority in the mode section lexically precedes any other character with the same sequence number and which has an accent priority. For example, assume *e* and *ê* have the same sequence number, but that *ê* has been assigned an accent priority in the mode section while the mode pointer of *e* is 0. Since *e* has not been assigned an accent priority,  $\text{LEX} ("e", "ê") = -1$ , or *e* precedes *ê* alphabetically.

In the following example, the two strings have the letter *e* (but with different accent marks) as the first character. Since the last characters of each string differ, a mismatch is encountered without using any accent priority information from the mode section.

$$\text{LEX} ("êz", "ëa") = -1$$

In the next example, both letters in one string match those in the other; the only difference being the accent marks on the first letters. In this case, the strings match and accent priority information is used to make a distinction. If the accent priority for *ë* has been specified to be less than the accent priority for *ê*, the results would be:

$$\text{LEX} ("êb", "ëb") = 1$$

### 1 For 2 Character Replacement

Another special case that you can specify is a 1 for 2 character replacement. This specifies that for collating purposes, a given combination of two characters is to be treated as a single character. That is, two characters are assigned a single sequence number.

An example of this are the letters "CH" together in the Spanish standard collating sequence. The correct alphabetical sequence in Spanish is A, B, C, CH, D... All words beginning with C followed by any letter other than H come before words beginning with C followed by H. Therefore, the two letters CH can be taken as a single letter coming between C and D in the alphabetical sequence.

The string "CHA" is to be stored as

ASCII Code for C	ASCII Code for H	ASCII Code for A
---------------------	---------------------	---------------------

but for collating purposes, the sequence numbers fetched for this string would be:

Sequence # for CH
Sequence # for A

The mode section format is:

Seq. # of 2-char. combination	ppercase of d character
Seq. # of 2-char. combination	vercase of d character
0	128 <sub>10</sub>

The front, or upper part of the first (and possibly second) entry, contains the sequence number of the 2-character combination. The last, or lower part of the entry, contains the ASCII value of the 2nd character.

Usually an entry for both the uppercase and lowercase versions of the 2nd character are specified. This is to catch both occurrences of the given character combination; for example, CH and Ch. However, this is not required. If the special 1 for 2 character replacement is to specify CH only (and not Ch), the lowercase entry would be omitted.

In the case that both an uppercase and lowercase entry appears, the sequence numbers in the upper part can be the same in both entries. They are the same in the Spanish table. However, different sequence numbers for each combination are acceptable also.

No matter how many entries are given, the last entry for this special case is equal to an octal 200, or a decimal 128. It is a terminator that indicates there are no further character combinations.

When a lexical calculation on a string encounters a character with a mode pointer to a mode entry of this type, the next character in the string is compared one-by-one with the possible second characters given in the mode entries. If a match is found, the two characters together are represented by the single sequence number given in the corresponding mode table entry. If no match is found, the first character is assigned the sequence number given in its collating section entry and processing continues with the next character in the string. Note that "don't care" characters in the string are not regarded as such for this type of replacement.

For example, the sequence numbers fetched for the string "CAH" (even if "A" is a "don't-care" character) would be:

Sequence # for C
Sequence # for A
Sequence # for H

Note that this special case does not cause any actual substitutions to be made in a string containing CH; CHA is still stored as:

ASCII Code for C	ASCII Code for H	ASCII Code for A
---------------------	---------------------	---------------------

This means that for collating purposes, the character pair CH causes a single sequence number to be fetched.

Consider again the example of CH for Spanish. Assume that the letter C has sequence number 67 and D has sequence number 68. To include the combination CH, a new sequence number must be created between 67 and 68. To do this, the collating sequence is rearranged so that C has sequence number 67 and D has 69. Sequence number 68 is reserved for CH and the remaining sequence numbers are adjusted accordingly.

Assume that the 4th, 5th, and 6th elements in the mode sequence contain the information governing the CH special case. The mode sequence entries for this example are shown next.

Array Element 70	67	4	← (4th Mode Section Entry)
71	69		
⋮			
⋮			
358	Sequence # for CH	H	
359	Sequence # for CH	h	
360	0	128 <sub>10</sub>	

Numerically, the mode section entries for this example would contain:

68	72
68	104
0	128

Note that 72 is the ASCII value for uppercase H and 104 is the ASCII value for lowercase h.

Whenever a C is encountered in collating and H or h immediately follows in the string, the pair of characters (CH or Ch) is assigned a single sequence number. Processing the string "CHA" would fetch the sequence numbers:

Sequence # for CH

Sequence # for A

In the original example, the actual contents of the mode sequence entries are:

4th Mode Section Entry:

	68	72		
Binary	01000100	01001000		
↓	┌───┐ ┌───┐ ┌───┐ ┌───┐			
Octal	4	2	1	1 0 <sub>8</sub>
↓				
Decimal	1	7	4	8 0 <sub>10</sub>

5th Mode Section Entry:

	68	104		
Binary	01000100	01101000		
↓	┌───┐ ┌───┐ ┌───┐ ┌───┐			
Octal	4	2	1	5 0 <sub>8</sub>
↓				
Decimal	1	7	5	1 2 <sub>10</sub>

Therefore, lexical table element 358 (1+1+256+96+4) contains 17 480<sub>10</sub>, element 359 contains 17 512<sub>10</sub>, and element 360 contains the terminator value 128<sub>10</sub>.

It is possible to specify additional character replacements for the same character. For example, if the combination CZ were to follow CH in the previous example, the mode section would be expanded to include the additional characters as shown below.

Sequence # for CH	H
Sequence # for Ch	h
Sequence # for CZ	Z
Sequence # for Cz	z
0	128

The terminator does not occur until all additional characters have been included.

### 2 For 1 Character Replacement

The third special case is a 2 for 1 character replacement. This type of mode entry specifies that a single character is to be assigned two sequence numbers for lexical comparisons. For example, the letter  $\ddot{A}$  in German is alphabetically equivalent to AE (or Ae). In a lexical comparison,  $\ddot{A}$  is assigned two sequence numbers; the sequence number for A and the sequence number for E (or e).

In general, if the mode entry specifies that a character is to be assigned two sequence numbers, the first sequence number is given in the normal collating section of the table. The second sequence number is given in the mode table entries. The format of the mode table entries is shown next.

Bit	15		8	7	0
Sequence # of 2nd character (upper)			0		
Sequence # of 2nd character (lower)			0		

Since there is only one number required for this case, the second portion of the element (bits 7-0) is not used, and must be set to zero to avoid confusion.

Note that mode entries of this special case require no terminator element. For example, assume  $\ddot{A}$  is a special case character that is alphabetically equivalent to AE (or Ae), and that the special case for  $\ddot{A}$  is handled in the first entry of the mode section. The collating section entry for  $\ddot{A}$  (which is character code 216) and the mode section entry are shown next.

Array Element 219	Sequence # for A	1 (Mode Pointer)	<div style="border-left: 1px solid black; border-right: 1px solid black; height: 15px; width: 100%;"></div> <div style="border-left: 1px solid black; border-right: 1px solid black; height: 15px; width: 100%;"></div> <div style="border-left: 1px solid black; border-right: 1px solid black; height: 15px; width: 100%;"></div>
...			
355	Sequence # for E	0	
356	Sequence # for e	0	<div style="border-left: 1px solid black; border-right: 1px solid black; height: 15px; width: 100%;"></div>

(1st Mode Section Entry)

Assuming that the sequence numbers of A, E, and e are the same as their ASCII values, the following illustration represents the numerical contents of the table entries.

219	65 <sub>10</sub>	1	16641 <sub>10</sub>
...			
355	69 <sub>10</sub>	0	17664 <sub>10</sub>
356	101 <sub>10</sub>	0	25856 <sub>10</sub>

Note that the sequence number for Ä in the collating section is equal to the sequence number for A. Differentiation occurs in the mode section information. That is, A probably will have no mode section pointer, while Ä will. The sequence difference between A and Ä, then, is handled by Ä's entry in the mode section.

The choice between lower and uppercase for the second sequence number is determined by the case of the character immediately following the special case letter. For ÄN, A is equivalent to AE. For Än, A is equivalent to Ae. The character Ä alone is equivalent to AE.

This special case does not cause any actual substitutions to be made in a string. In the previous example, the string "Än" is stored as:

ASCII CODE for Ä	ASCII Code for n
---------------------	---------------------

For collating purposes, the sequence numbers fetched for the string "Än" would be:

Sequence # for A
Sequence # for e
Sequence # for n

### "Don't Care" Characters

There is a fourth special case which doesn't require a mode section entry. This is the "don't care" special case where the specified character is to be ignored for collating purposes. An example of such a case is the hyphen in hyphenated words.

A character is ignored alphabetically if it has a sequence number of 255 in the collating section. That is, wherever a "don't care" condition exists in the collating section, simply enter 255 as the sequence number in that element.

Note that once a character is specified to be a "don't care" character, it is always treated as if it did not appear for collating purposes. For example, if the hyphen is designated as a "don't care" character, then "user-defined" would be collated as "userdefined", and, consequently, "-2" would be collated as "2". Therefore, consideration should be given to the desired results before a character is designated as a "don't care" condition.

It should also be noted that the null string precedes a string containing only a "don't care" character. That is, if A\$ = "" and B\$ = "-", then A\$ precedes B\$ even if the hyphen is a "don't care" character.

# ASCII Character Codes

ASCII Char.	EQUIVALENT FORMS			
	Binary	Oct	Hex	Dec
NULL	00000000	000	00	0
SOH	00000001	001	01	1
STX	00000010	002	02	2
ETX	00000011	003	03	3
EOT	00000100	004	04	4
ENQ	00000101	005	05	5
ACK	00000110	006	06	6
BELL	00000111	007	07	7
BS	00001000	010	08	8
HT	00001001	011	09	9
LF	00001010	012	0A	10
VT	00001011	013	0B	11
FF	00001100	014	0C	12
CR	00001101	015	0D	13
SO	00001110	016	0E	14
SI	00001111	017	0F	15
DLE	00010000	020	10	16
DC1	00010001	021	11	17
DC2	00010010	022	12	18
DC3	00010011	023	13	19
DC4	00010100	024	14	20
NAK	00010101	025	15	21
SYNC	00010110	026	16	22
ETB	00010111	027	17	23
CAN	00011000	030	18	24
EM	00011001	031	19	25
SUB	00011010	032	1A	26
ESC	00011011	033	1B	27
FS	00011100	034	1C	28
GS	00011101	035	1D	29
RS	00011110	036	1E	30
US	00011111	037	1F	31

ASCII Char.	EQUIVALENT FORMS			
	Binary	Oct	Hex	Dec
space	00100000	040	20	32
!	00100001	041	21	33
"	00100010	042	22	34
#	00100011	043	23	35
\$	00100100	044	24	36
%	00100101	045	25	37
&	00100110	046	26	38
'	00100111	047	27	39
(	00101000	050	28	40
)	00101001	051	29	41
*	00101010	052	2A	42
+	00101011	053	2B	43
,	00101100	054	2C	44
-	00101101	055	2D	45
.	00101110	056	2E	46
/	00101111	057	2F	47
0	00110000	060	30	48
1	00110001	061	31	49
2	00110010	062	32	50
3	00110011	063	33	51
4	00110100	064	34	52
5	00110101	065	35	53
6	00110110	066	36	54
7	00110111	067	37	55
8	00111000	070	38	56
9	00111001	071	39	57
:	00111010	072	3A	58
;	00111011	073	3B	59
<	00111100	074	3C	60
=	00111101	075	3D	61
>	00111110	076	3E	62
?	00111111	077	3F	63

ASCII Char.	EQUIVALENT FORMS			
	Binary	Oct	Hex	Dec
@	01000000	100	40	64
A	01000001	101	41	65
B	01000010	102	42	66
C	01000011	103	43	67
D	01000100	104	44	68
E	01000101	105	45	69
F	01000110	106	46	70
G	01000111	107	47	71
H	01001000	110	48	72
I	01001001	111	49	73
J	01001010	112	4A	74
K	01001011	113	4B	75
L	01001100	114	4C	76
M	01001101	115	4D	77
N	01001110	116	4E	78
O	01001111	117	4F	79
P	01010000	120	50	80
Q	01010001	121	51	81
R	01010010	122	52	82
S	01010011	123	53	83
T	01010100	124	54	84
U	01010101	125	55	85
V	01010110	126	56	86
W	01010111	127	57	87
X	01011000	130	58	88
Y	01011001	131	59	89
Z	01011010	132	5A	90
[	01011011	133	5B	91
\	01011100	134	5C	92
]	01011101	135	5D	93
^	01011110	136	5E	94
_	01011111	137	5F	95

ASCII Char.	EQUIVALENT FORMS			
	Binary	Oct	Hex	Dec
`	01100000	140	60	96
a	01100001	141	61	97
b	01100010	142	62	98
c	01100011	143	63	99
d	01100100	144	64	100
e	01100101	145	65	101
f	01100110	146	66	102
g	01100111	147	67	103
h	01101000	150	68	104
i	01101001	151	69	105
j	01101010	152	6A	106
k	01101011	153	6B	107
l	01101100	154	6C	108
m	01101101	155	6D	109
n	01101110	156	6E	110
o	01101111	157	6F	111
p	01110000	160	70	112
q	01110001	161	71	113
r	01110010	162	72	114
s	01110011	163	73	115
t	01110100	164	74	116
u	01110101	165	75	117
v	01110110	166	76	118
w	01110111	167	77	119
x	01111000	170	78	120
y	01111001	171	79	121
z	01111010	172	7A	122
{	01111011	173	7B	123
	01111100	174	7C	124
}	01111101	175	7D	125
~	01111110	176	7E	126
DEL	01111111	177	7F	127



## Roman Extension Character Codes

ASCII Char.	EQUIVALENT FORMS		
	Binary	Octal	Decimal
CLEAR	10000000	200	128
IV	10000001	201	129
BL	10000010	202	130
IV/BL	10000011	203	131
UL	10000100	204	132
IV/UL	10000101	205	133
BL/UL	10000110	206	134
IV/BL/UL	10000111	207	135
	10010111	227	151
	10101000	250	168
	10101001	251	169
	10101010	252	170
	10001000	210	136
	10001001	211	137
	10001010	212	138
	10001011	213	139
	10001100	214	140
	10001101	215	141
	10001110	216	142
	10001111	217	143
	10010000	220	144
	10010001	221	145
	10010010	222	146
	10010011	223	147
	10010100	224	148
	10010101	225	149
	10010110	226	150
	10010111	227	157
	10011000	230	152
	10011001	231	153
	10011010	232	154
	10011011	233	155
	10011100	234	156

ASCII Char.	EQUIVALENT FORMS		
	Binary	Octal	Decimal
	10011101	235	157
	10011110	236	158
	10011111	237	159
	10100000	240	160
À	10100001	241	161
Á	10100010	242	162
Â	10100011	243	163
Ã	10100100	244	164
Ä	10100101	245	165
Å	10100110	246	166
Ö	10100111	247	167
Ç	00101000	250	168
È	00101001	251	169
É	00101010	252	170
Ê	10101011	253	171
Ë	10101100	254	172
Ë	10101101	255	173
Ï	10101110	256	174
Ë	10101111	257	175
—	10110000	260	176
Ë	10110001	261	177
Ë	10110010	262	178
◊	10110011	263	179
Ç	10110100	264	180
Ç	10110101	265	181
Ë	10110110	266	182
Ë	10110111	267	183
ï	10111000	270	184
◊	10111001	271	185
◊	10111010	272	186
Ë	10111011	273	187
◊	10111100	274	188
Ë	10111101	275	189

ASCII Char.	EQUIVALENT FORMS		
	Binary	Octal	Decimal
◊	10111110	276	190
+	10111111	277	191
ä	11000000	300	192
ë	11000001	301	193
ö	11000010	302	194
ü	11000011	303	195
ä	11000100	304	196
ë	11000101	305	197
ö	11000110	306	198
ü	11000111	307	199
ä	11001000	310	200
ë	11001001	311	201
ö	11001010	312	202
ü	11001011	313	203
ä	11001100	314	204
ë	11001101	315	205
ö	11001110	316	206
ü	11001111	317	207
Ä	11010000	320	208
Ë	11010001	321	209
Ö	11010010	322	210
Ë	11010011	323	211
ä	11010100	324	212
ï	11010101	325	213
ö	11010110	326	214
ü	11010111	327	215
Ä	11011000	330	216
ï	11011001	331	217
Ö	11011010	332	218
Ü	11011011	333	219
Ë	11011100	334	220
Ë	11011101	335	221
Ë	11011110	336	222

ASCII Char.	EQUIVALENT FORMS		
	Binary	Octal	Decimal
	11011111	337	223
	11100000	340	224
	11100001	341	225
	11100010	342	226
	11100011	343	227
	11100100	344	228
	11100101	345	229
	11100110	346	230
	11100111	347	231
	11101000	350	232
	11101001	351	233
	11101010	352	234
	11101011	353	235
	11101100	354	236
	11101101	355	237
	11101110	356	238
	11101111	357	239
	11110000	360	240
	11110001	361	241
	11110010	362	242
	11110011	363	243
	11110100	364	244
	11110101	365	245
	11110110	366	246
	11110111	367	247
	11111000	370	248
	11111001	371	249
	11111010	372	250
	11111011	373	251
	11111100	374	252
	11111101	375	253
	11111110	376	254
	11111111	377	255

IV - Inverse video  
 BL - Blinking  
 UL - Underline

## Advanced Programming Lexical Tables

The following listings are the Advanced Programming lexical tables. The tables are recorded on the Lexical Tables Cartridge under the names ASCII, FRENCH, GERMAN, SWEDSH and SPANSH. For your convenience, backup tables are included on the cartridge under the names BKUPAS, BKUPFR, BKUPSW and BKUPSP.

ASCII TABLE											
-------------	--	--	--	--	--	--	--	--	--	--	--

ENTRY #1 -- LENGTH OF COLLATING TABLE = 354  
 ENTRY #2 -- LENGTH OF MODE TABLE = 0

----- COLLATING SECTION -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
3	1	0	NUL	0	0	51	49	48	0	48	0
4	2	1	SOH	1	0	52	50	49	1	49	0
5	3	2	STX	2	0	53	51	50	2	50	0
6	4	3	ETX	3	0	54	52	51	3	51	0
7	5	4	EOT	4	0	55	53	52	4	52	0
8	6	5	ENQ	5	0	56	54	53	5	53	0
9	7	6	ACK	6	0	57	55	54	6	54	0
10	8	7	BEL	7	0	58	56	55	7	55	0
11	9	8	BS	8	0	59	57	56	8	56	0
12	10	9	HT	9	0	60	58	57	9	57	0
13	11	10	LF	10	0	61	59	58	:	58	0
14	12	11	VT	11	0	62	60	59	;	59	0
15	13	12	FF	12	0	63	61	60	<	60	0
16	14	13	CR	13	0	64	62	61	=	61	0
17	15	14	SO	14	0	65	63	62	>	62	0
18	16	15	SI	15	0	66	64	63	?	63	0
19	17	16	DLE	16	0	67	65	64	@	64	0
20	18	17	DC1	17	0	68	66	65	A	65	0
21	19	18	DC2	18	0	69	67	66	B	66	0
22	20	19	DC3	19	0	70	68	67	C	67	0
23	21	20	DC4	20	0	71	69	68	D	68	0
24	22	21	NAK	21	0	72	70	69	E	69	0
25	23	22	SYN	22	0	73	71	70	F	70	0
26	24	23	ETB	23	0	74	72	71	G	71	0
27	25	24	CAN	24	0	75	73	72	H	72	0
28	26	25	EM	25	0	76	74	73	I	73	0
29	27	26	SUB	26	0	77	75	74	J	74	0
30	28	27	ESC	27	0	78	76	75	K	75	0
31	29	28	FS	28	0	79	77	76	L	76	0
32	30	29	GS	29	0	80	78	77	M	77	0
33	31	30	RS	30	0	81	79	78	N	78	0
34	32	31	US	31	0	82	80	79	O	79	0
35	33	32	SP	32	0	83	81	80	P	80	0
36	34	33	!	33	0	84	82	81	Q	81	0
37	35	34	"	34	0	85	83	82	R	82	0
38	36	35	#	35	0	86	84	83	S	83	0
39	37	36	\$	36	0	87	85	84	T	84	0
40	38	37	%	37	0	88	86	85	U	85	0
41	39	38	&	38	0	89	87	86	V	86	0
42	40	39	<	39	0	90	88	87	W	87	0
43	41	40	(	40	0	91	89	88	X	88	0
44	42	41	)	41	0	92	90	89	Y	89	0
45	43	42	*	42	0	93	91	90	Z	90	0
46	44	43	+	43	0	94	92	91	[	91	0
47	45	44	,	44	0	95	93	92	\	92	0
48	46	45	-	45	0	96	94	93	]	93	0
49	47	46	.	46	0	97	95	94	^	94	0
50	48	47	/	47	0	98	96	95	_	95	0

## ----- COLLATING SECTION - CONT -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
99	97	96	\	96	0	147	145	144		144	0
100	98	97	a	97	0	148	146	145		145	0
101	99	98	b	98	0	149	147	146		146	0
102	100	99	c	99	0	150	148	147		147	0
103	101	100	d	100	0	151	149	148		148	0
104	102	101	e	101	0	152	150	149		149	0
105	103	102	f	102	0	153	151	150		150	0
106	104	103	g	103	0	154	152	151		151	0
107	105	104	h	104	0	155	153	152		152	0
108	106	105	i	105	0	156	154	153		153	0
109	107	106	j	106	0	157	155	154		154	0
110	108	107	k	107	0	158	156	155		155	0
111	109	108	l	108	0	159	157	156		156	0
112	110	109	m	109	0	160	158	157		157	0
113	111	110	n	110	0	161	159	158		158	0
114	112	111	o	111	0	162	160	159		159	0
115	113	112	p	112	0	163	161	160		160	0
116	114	113	q	113	0	164	162	161	A	161	0
117	115	114	r	114	0	165	163	162	I	162	0
118	116	115	s	115	0	166	164	163	O	163	0
119	117	116	t	116	0	167	165	164	U	164	0
120	118	117	u	117	0	168	166	165	A	165	0
121	119	118	v	118	0	169	167	166	E	166	0
122	120	119	w	119	0	170	168	167	O	167	0
123	121	120	x	120	0	171	169	168	C	168	0
124	122	121	y	121	0	172	170	169	\	169	0
125	123	122	z	122	0	173	171	170	^	170	0
126	124	123	{	123	0	174	172	171	~	171	0
127	125	124		124	0	175	173	172	^	172	0
128	126	125	}	125	0	176	174	173	E	173	0
129	127	126	~	126	0	177	175	174	O	174	0
130	128	127	DEL	127	0	178	176	175	I	175	0
131	129	128	CLR	128	0	179	177	176	A	176	0
132	130	129	IV	129	0	180	178	177	A	177	0
133	131	130	BL	130	0	181	179	178	A	178	0
134	132	131	IV-B	131	0	182	180	179	o	179	0
135	133	132	UL	132	0	183	181	180	C	180	0
136	134	133	IV-U	133	0	184	182	181	C	181	0
137	135	134	BL-U	134	0	185	183	182	N	182	0
138	136	135	I-B-U	135	0	186	184	183	N	183	0
139	137	136		136	0	187	185	184	i	184	0
140	138	137		137	0	188	186	185	z	185	0
141	139	138		138	0	189	187	186	o	186	0
142	140	139		139	0	190	188	187	E	187	0
143	141	140		140	0	191	189	188	e	188	0
144	142	141		141	0	192	190	189	s	189	0
145	143	142		142	0	193	191	190	e	190	0
146	144	143		143	0	194	192	191	.	191	0

## ----- COLLATING SECTION - CONT -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
195	193	192	ä	192	0	243	241	240		240	0
196	194	193	ë	193	0	244	242	241		241	0
197	195	194	ö	194	0	245	243	242		242	0
198	196	195	û	195	0	246	244	243		243	0
199	197	196	ä	196	0	247	245	244		244	0
200	198	197	ë	197	0	248	246	245		245	0
201	199	198	ö	198	0	249	247	246		246	0
202	200	199	û	199	0	250	248	247		247	0
203	201	200	ä	200	0	251	249	248		248	0
204	202	201	ë	201	0	252	250	249		249	0
205	203	202	ö	202	0	253	251	250		250	0
206	204	203	û	203	0	254	252	251		251	0
207	205	204	ä	204	0	255	253	252		252	0
208	206	205	ë	205	0	256	254	253		253	0
209	207	206	ö	206	0	257	255	254		254	0
210	208	207	û	207	0	258	256	255		255	0
211	209	208	ä	208	0						
212	210	209	ë	209	0						
213	211	210	ö	210	0						
214	212	211	û	211	0						
215	213	212	ä	212	0						
216	214	213	ë	213	0						
217	215	214	ö	214	0						
218	216	215	û	215	0						
219	217	216	ä	216	0						
220	218	217	ë	217	0						
221	219	218	ö	218	0						
222	220	219	û	219	0						
223	221	220	ä	220	0						
224	222	221	ë	221	0						
225	223	222	ö	222	0						
226	224	223		223	0						
227	225	224		224	0						
228	226	225		225	0						
229	227	226		226	0						
230	228	227		227	0						
231	229	228		228	0						
232	230	229		229	0						
233	231	230		230	0						
234	232	231		231	0						
235	233	232		232	0						
236	234	233		233	0						
237	235	234		234	0						
238	236	235		235	0						
239	237	236		236	0						
240	238	237		237	0						
241	239	238		238	0						
242	240	239		239	0						

## ----- UPPERCASE/LOWERCASE SECTION -----

ENTRY #	U/L ENTRY	CHAR CODE	CHAR	UPPER CASE	LOWER CASE	ENTRY #	U/L ENTRY	CHAR CODE	CHAR	UPPER CASE	LOWER CASE
259	1	160				307	49	208	À	À	à
260	2	161	À	À	à	308	50	209	Í	Í	í
261	3	162	Ī	Ī	ī	309	51	210	Ō	Ō	ō
262	4	163	Ō	Ō	ō	310	52	211	Ē	Ē	ē
263	5	164	Ū	Ū	ū	311	53	212	Ā	Ā	ā
264	6	165	Ā	Ā	ā	312	54	213	Ī	Ī	ī
265	7	166	Ē	Ē	ē	313	55	214	Ō	Ō	ō
266	8	167	Ō	Ō	ō	314	56	215	Ē	Ē	ē
267	9	168	·	·	·	315	57	216	Ā	Ā	ā
268	10	169	·	·	·	316	58	217	Ī	Ī	ī
269	11	170	·	·	·	317	59	218	Ō	Ō	ō
270	12	171	·	·	·	318	60	219	Ū	Ū	ū
271	13	172	·	·	·	319	61	220	Ē	Ē	ē
272	14	173	Ē	Ē	ē	320	62	221	Ī	Ī	ī
273	15	174	Ō	Ō	ō	321	63	222	Ī	Ī	ī
274	16	175	Ē	Ē	ē	322	64	223			
275	17	176				323	65	224			
276	18	177	Ā	Ā	ā	324	66	225			
277	19	178	ā	ā	ā	325	67	226			
278	20	179	°	°	°	326	68	227			
279	21	180	Ç	Ç	ç	327	69	228			
280	22	181	Ç	Ç	ç	328	70	229			
281	23	182	Ñ	Ñ	ñ	329	71	230			
282	24	183	Ñ	Ñ	ñ	330	72	231			
283	25	184	ı	ı	ı	331	73	232			
284	26	185	ı	ı	ı	332	74	233			
285	27	186	ı	ı	ı	333	75	234			
286	28	187	£	£	£	334	76	235			
287	29	188	£	£	£	335	77	236			
288	30	189	£	£	£	336	78	237			
289	31	190	£	£	£	337	79	238			
290	32	191	·	·	·	338	80	239			
291	33	192	ā	ā	ā	339	81	240			
292	34	193	ē	ē	ē	340	82	241			
293	35	194	ō	ō	ō	341	83	242			
294	36	195	ū	ū	ū	342	84	243			
295	37	196	ā	ā	ā	343	85	244			
296	38	197	ē	ē	ē	344	86	245			
297	39	198	ō	ō	ō	345	87	246			
298	40	199	ū	ū	ū	346	88	247			
299	41	200	ā	ā	ā	347	89	248			
300	42	201	ē	ē	ē	348	90	249			
301	43	202	ō	ō	ō	349	91	250			
302	44	203	ū	ū	ū	350	92	251			
303	45	204	ā	ā	ā	351	93	252			
304	46	205	ē	ē	ē	352	94	253			
305	47	206	ō	ō	ō	353	95	254			
306	48	207	ū	ū	ū	354	96	255			



## ----- MODE SECTION -----

ENTRY #	MODE ENTRY	TYPE DESCRIPTION
---------	------------	------------------

||||| TYPE DETAILS |||||

-- DON'T CARES --

0 ASCII - 255 = Don't care

-- ACCENT PRIORITIES --

-- TWO FOR ONE REPLACEMENTS --

-- ONE FOR TWO REPLACEMENTS --

FRENCH TABLE											
--------------	--	--	--	--	--	--	--	--	--	--	--

ENTRY #1 -- LENGTH OF COLLATING TABLE = 374  
 ENTRY #2 -- LENGTH OF MODE TABLE = 20

----- COLLATING SECTION -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
3	1	0	NUL	0	0	51	49	48	0	48	0
4	2	1	SOH	1	0	52	50	49	1	49	0
5	3	2	STX	2	0	53	51	50	2	50	0
6	4	3	ETX	3	0	54	52	51	3	51	0
7	5	4	EOT	4	0	55	53	52	4	52	0
8	6	5	ENO	5	0	56	54	53	5	53	0
9	7	6	ACK	6	0	57	55	54	6	54	0
10	8	7	BEL	7	0	58	56	55	7	55	0
11	9	8	BS	8	0	59	57	56	8	56	0
12	10	9	HT	9	0	60	58	57	9	57	0
13	11	10	LF	10	0	61	59	58	:	58	0
14	12	11	VT	11	0	62	60	59	;	59	0
15	13	12	FF	12	0	63	61	60	<	60	0
16	14	13	CR	13	0	64	62	61	=	61	0
17	15	14	SO	14	0	65	63	62	>	62	0
18	16	15	SI	15	0	66	64	63	?	63	0
19	17	16	DLE	16	0	67	65	64	@	64	0
20	18	17	DC1	17	0	68	66	65	A	65	0
21	19	18	DC2	18	0	69	67	66	B	68	0
22	20	19	DC3	19	0	70	68	67	C	69	0
23	21	20	DC4	20	0	71	69	68	D	70	0
24	22	21	NAK	21	0	72	70	69	E	71	0
25	23	22	SYN	22	0	73	71	70	F	72	0
26	24	23	ETB	23	0	74	72	71	G	73	0
27	25	24	CAN	24	0	75	73	72	H	74	0
28	26	25	EM	25	0	76	74	73	I	75	0
29	27	26	SUB	26	0	77	75	74	J	76	0
30	28	27	ESC	27	0	78	76	75	K	77	0
31	29	28	FS	28	0	79	77	76	L	78	0
32	30	29	GS	29	0	80	78	77	M	79	0
33	31	30	RS	30	0	81	79	78	N	80	0
34	32	31	US	31	0	82	80	79	O	82	0
35	33	32	SP	32	0	83	81	80	P	84	0
36	34	33	!	33	0	84	82	81	Q	85	0
37	35	34	"	34	0	85	83	82	R	86	0
38	36	35	#	35	0	86	84	83	S	87	0
39	37	36	\$	36	0	87	85	84	T	88	0
40	38	37	%	37	0	88	86	85	U	89	0
41	39	38	&	38	0	89	87	86	V	91	0
42	40	39	<	39	0	90	88	87	W	92	0
43	41	40	(	40	0	91	89	88	X	93	0
44	42	41	)	41	0	92	90	89	Y	94	0
45	43	42	*	42	0	93	91	90	Z	95	0
46	44	43	+	43	0	94	92	91	[	96	0
47	45	44	,	44	0	95	93	92	\	97	0
48	46	45	-	255	0	96	94	93	]	98	0
49	47	46	.	46	0	97	95	94	^	99	0
50	48	47	/	47	0	98	96	95	_	100	0



## ----- COLLATING SECTION - CONT -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
99	97	96	\	101	0	147	145	144		0	0
100	98	97	a	102	0	148	146	145		0	0
101	99	98	b	103	0	149	147	146		0	0
102	100	99	c	104	0	150	148	147		0	0
103	101	100	d	106	0	151	149	148		0	0
104	102	101	e	107	0	152	150	149		0	0
105	103	102	f	108	0	153	151	150		0	0
106	104	103	g	109	0	154	152	151		0	0
107	105	104	h	110	0	155	153	152		0	0
108	106	105	i	111	0	156	154	153		0	0
109	107	106	j	112	0	157	155	154		0	0
110	108	107	k	113	0	158	156	155		0	0
111	109	108	l	114	0	159	157	156		0	0
112	110	109	m	115	0	160	158	157		0	0
113	111	110	n	116	0	161	159	158		0	0
114	112	111	o	118	0	162	160	159		0	0
115	113	112	p	119	0	163	161	160		0	0
116	114	113	q	120	0	164	162	161	A	65	0
117	115	114	r	121	0	165	163	162	i	75	0
118	116	115	s	122	0	166	164	163	o	82	0
119	117	116	t	123	0	167	165	164	U	89	0
120	118	117	u	124	0	168	166	165	A	65	0
121	119	118	v	125	0	169	167	166	E	71	0
122	120	119	w	126	0	170	168	167	o	82	0
123	121	120	x	127	0	171	169	168		137	0
124	122	121	y	128	0	172	170	169		138	0
125	123	122	z	129	0	173	171	170		139	0
126	124	123	(	133	0	174	172	171		140	0
127	125	124		134	0	175	173	172		141	0
128	126	125	)	135	0	176	174	173	E	71	0
129	127	126	^	136	0	177	175	174	o	82	0
130	128	127	DEL	0	0	178	176	175	E	142	0
131	129	128	CLR	0	0	179	177	176	T	143	0
132	130	129	IV	0	0	180	178	177	A	65	0
133	131	130	BL	0	0	181	179	178	A	102	0
134	132	131	IV-B	0	0	182	180	179	o	144	0
135	133	132	UL	0	0	183	181	180	G	69	0
136	134	133	IV-U	0	0	184	182	181	C	105	0
137	135	134	BL-U	0	0	185	183	182	N	81	0
138	136	135	I-B-U	0	0	186	184	183	A	117	0
139	137	136		0	0	187	185	184	i	145	0
140	138	137		0	0	188	186	185	o	146	0
141	139	138		0	0	189	187	186	o	147	0
142	140	139		0	0	190	188	187	E	148	0
143	141	140		0	0	191	189	188	o	149	0
144	142	141		0	0	192	190	189	S	150	0
145	143	142		0	0	193	191	190	o	151	0
146	144	143		0	0	194	192	191	.	152	0

## ----- COLLATING SECTION - CONT -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
195	193	192	ä	102	0	243	241	240		0	0
196	194	193	é	107	0	244	242	241		0	0
197	195	194	ó	118	0	245	243	242		0	0
198	196	195	ú	124	0	246	244	243		0	0
199	197	196	ä	102	0	247	245	244		0	0
200	198	197	é	107	0	248	246	245		0	0
201	199	198	ó	118	0	249	247	246		0	0
202	200	199	ú	124	0	250	248	247		0	0
203	201	200	ä	102	0	251	249	248		0	0
204	202	201	é	107	0	252	250	249		0	0
205	203	202	ó	118	0	253	251	250		0	0
206	204	203	ú	124	0	254	252	251		0	0
207	205	204	ä	102	0	255	253	252		0	0
208	206	205	é	107	0	256	254	253		0	0
209	207	206	ó	118	0	257	255	254		0	0
210	208	207	ú	124	0	258	256	255		0	0
211	209	208	ä	65	0						
212	210	209	é	111	0						
213	211	210	ó	82	0						
214	212	211	ú	65	0						
215	213	212	ä	102	0						
216	214	213	é	111	0						
217	215	214	ó	118	0						
218	216	215	ú	102	0						
219	217	216	ä	65	0						
220	218	217	é	111	0						
221	219	218	ó	82	0						
222	220	219	ú	89	0						
223	221	220	ä	71	0						
224	222	221	é	111	0						
225	223	222	ó	122	1						
226	224	223		0	0						
227	225	224		0	0						
228	226	225		0	0						
229	227	226		0	0						
230	228	227		0	0						
231	229	228		0	0						
232	230	229		0	0						
233	231	230		0	0						
234	232	231		0	0						
235	233	232		0	0						
236	234	233		0	0						
237	235	234		0	0						
238	236	235		0	0						
239	237	236		0	0						
240	238	237		0	0						
241	239	238		0	0						
242	240	239		0	0						

## ----- UPPERCASE/LOWERCASE SECTION -----

ENTRY #	U/L ENTRY	CHAR CODE	CHAR	UPPER CASE	LOWER CASE	ENTRY #	U/L ENTRY	CHAR CODE	CHAR	UPPER CASE	LOWER CASE
259	1	160				307	49	208	À	À	à
260	2	161	À	À	à	308	50	209	Í	Í	í
261	3	162	Ī	Ī	ī	309	51	210	Ō	Ō	ō
262	4	163	Ó	Ó	ó	310	52	211	Ė	Ė	ė
263	5	164	Ū	Ū	ū	311	53	212	À	À	à
264	6	165	Ā	Ā	ā	312	54	213	Ī	Ī	ī
265	7	166	Ē	Ē	ē	313	55	214	Ō	Ō	ō
266	8	167	Ō	Ō	ō	314	56	215	Ė	Ė	ė
267	9	168	·	·	·	315	57	216	À	À	à
268	10	169	·	·	·	316	58	217	Ī	Ī	ī
269	11	170	·	·	·	317	59	218	Ō	Ō	ō
270	12	171	·	·	·	318	60	219	Ū	Ū	ū
271	13	172	·	·	·	319	61	220	Ē	Ē	ē
272	14	173	Ē	Ē	ē	320	62	221	Ī	Ī	ī
273	15	174	Ō	Ō	ō	321	63	222	Ō	Ō	ō
274	16	175	Ē	Ē	ē	322	64	223			
275	17	176	Ē	Ē	ē	323	65	224			
276	18	177	Ā	Ā	ā	324	66	225			
277	19	178	ā	ā	ā	325	67	226			
278	20	179	°	°	°	326	68	227			
279	21	180	Ç	Ç	ç	327	69	228			
280	22	181	Ç	Ç	ç	328	70	229			
281	23	182	Ñ	Ñ	ñ	329	71	230			
282	24	183	Ñ	Ñ	ñ	330	72	231			
283	25	184	İ	İ	ı	331	73	232			
284	26	185	İ	İ	ı	332	74	233			
285	27	186	Œ	Œ	œ	333	75	234			
286	28	187	£	£	£	334	76	235			
287	29	188	£	£	£	335	77	236			
288	30	189	£	£	£	336	78	237			
289	31	190	£	£	£	337	79	238			
290	32	191	·	·	·	338	80	239			
291	33	192	À	À	à	339	81	240			
292	34	193	Ē	Ē	ē	340	82	241			
293	35	194	Ō	Ō	ō	341	83	242			
294	36	195	Ū	Ū	ū	342	84	243			
295	37	196	Ā	Ā	ā	343	85	244			
296	38	197	Ē	Ē	ē	344	86	245			
297	39	198	Ō	Ō	ō	345	87	246			
298	40	199	Ū	Ū	ū	346	88	247			
299	41	200	À	À	à	347	89	248			
300	42	201	Ē	Ē	ē	348	90	249			
301	43	202	Ō	Ō	ō	349	91	250			
302	44	203	Ū	Ū	ū	350	92	251			
303	45	204	Ā	Ā	ā	351	93	252			
304	46	205	Ē	Ē	ē	352	94	253			
305	47	206	Ō	Ō	ō	353	95	254			
306	48	207	Ū	Ū	ū	354	96	255			

## ----- MODE SECTION -----

ENTRY #	MODE ENTRY	TYPE DESCRIPTION
355	1	TWO FOR ONE CHARACTER REPLACEMENT
356	2	TWO FOR ONE CHARACTER REPLACEMENT

||||| TYPE DETAILS |||||

-- DON'T CARES --

0 - ASCII - 45 = Don't care

-- ACCENT PRIORITIES --

-- TWO FOR ONE REPLACEMENTS --

ENTRY #	REPLACEMENT	ASCII
1	B = ss = ss	222

-- ONE FOR TWO REPLACEMENTS --

GERMAN TABLE											
--------------	--	--	--	--	--	--	--	--	--	--	--

ENTRY #1 -- LENGTH OF COLLATING TABLE = 374

ENTRY #2 -- LENGTH OF MODE TABLE = 20

## ----- COLLATING SECTION -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
3	1	0	NUL	0	0	51	49	48	0	48	0
4	2	1	SOH	1	0	52	50	49	1	49	0
5	3	2	STX	2	0	53	51	50	2	50	0
6	4	3	ETX	3	0	54	52	51	3	51	0
7	5	4	EOT	4	0	55	53	52	4	52	0
8	6	5	ENQ	5	0	56	54	53	5	53	0
9	7	6	ACK	6	0	57	55	54	6	54	0
10	8	7	BEL	7	0	58	56	55	7	55	0
11	9	8	BS	8	0	59	57	56	8	56	0
12	10	9	HT	9	0	60	58	57	9	57	0
13	11	10	LF	10	0	61	59	58	:	58	0
14	12	11	VT	11	0	62	60	59	;	59	0
15	13	12	FF	12	0	63	61	60	<	60	0
16	14	13	CR	13	0	64	62	61	=	61	0
17	15	14	SO	14	0	65	63	62	>	62	0
18	16	15	SI	15	0	66	64	63	?	63	0
19	17	16	DLE	16	0	67	65	64	@	64	0
20	18	17	DC1	17	0	68	66	65	A	65	0
21	19	18	DC2	18	0	69	67	66	B	71	0
22	20	19	DC3	19	0	70	68	67	C	72	0
23	21	20	DC4	20	0	71	69	68	D	74	0
24	22	21	NAK	21	0	72	70	69	E	75	0
25	23	22	SYN	22	0	73	71	70	F	79	0
26	24	23	ETB	23	0	74	72	71	G	80	0
27	25	24	CAN	24	0	75	73	72	H	81	0
28	26	25	EM	25	0	76	74	73	I	82	0
29	27	26	SUB	26	0	77	75	74	J	84	0
30	28	27	ESC	27	0	78	76	75	K	85	0
31	29	28	FS	28	0	79	77	76	L	86	0
32	30	29	GS	29	0	80	78	77	M	87	0
33	31	30	RS	30	0	81	79	78	N	88	0
34	32	31	US	31	0	82	80	79	O	90	0
35	33	32	SP	32	0	83	81	80	P	95	0
36	34	33	!	33	0	84	82	81	Q	96	0
37	35	34	"	34	0	85	83	82	R	97	0
38	36	35	#	35	0	86	84	83	S	98	0
39	37	36	\$	36	0	87	85	84	T	99	0
40	38	37	%	37	0	88	86	85	U	100	0
41	39	38	&	38	0	89	87	86	V	102	0
42	40	39	/	39	0	90	88	87	W	103	0
43	41	40	(	40	0	91	89	88	X	104	0
44	42	41	)	41	0	92	90	89	Y	105	0
45	43	42	*	42	0	93	91	90	Z	106	0
46	44	43	+	43	0	94	92	91	[	107	0
47	45	44	,	44	0	95	93	92	\	108	0
48	46	45	-	45	0	96	94	93	]	109	0
49	47	46	.	46	0	97	95	94	^	110	0
50	48	47	/	47	0	98	96	95	_	111	0

## ----- COLLATING SECTION - CONT -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
99	97	96	`	112	0	147	145	144		0	0
100	98	97	a	113	0	148	146	145		0	0
101	99	98	b	120	0	149	147	146		0	0
102	100	99	c	121	0	150	148	147		0	0
103	101	100	d	123	0	151	149	148		0	0
104	102	101	e	124	0	152	150	149		0	0
105	103	102	f	129	0	153	151	150		0	0
106	104	103	g	130	0	154	152	151		0	0
107	105	104	h	131	0	155	153	152		0	0
108	106	105	i	132	0	156	154	153		0	0
109	107	106	j	137	0	157	155	154		0	0
110	108	107	k	138	0	158	156	155		0	0
111	109	108	l	139	0	159	157	156		0	0
112	110	109	m	140	0	160	158	157		0	0
113	111	110	n	141	0	161	159	158		0	0
114	112	111	o	143	0	162	160	159		0	0
115	113	112	p	148	0	163	161	160		0	0
116	114	113	q	149	0	164	162	161	A	68	0
117	115	114	r	150	0	165	163	162	I	83	0
118	116	115	s	151	0	166	164	163	O	91	0
119	117	116	t	152	0	167	165	164	U	101	0
120	118	117	u	153	0	168	166	165	A	69	0
121	119	118	v	157	0	169	167	166	E	77	0
122	120	119	w	158	0	170	168	167	O	92	0
123	121	120	x	159	0	171	169	168		166	0
124	122	121	y	160	0	172	170	169		167	0
125	123	122	z	161	0	173	171	170		168	0
126	124	123	{	162	0	174	172	171		169	0
127	125	124		163	0	175	173	172		170	0
128	126	125	}	164	0	176	174	173	E	78	0
129	127	126	~	165	0	177	175	174	O	93	0
130	128	127	DEL	0	0	178	176	175	E	171	0
131	129	128	CLR	0	0	179	177	176	I	172	0
132	130	129	IV	0	0	180	178	177	B	70	0
133	131	130	BL	0	0	181	179	178	B	119	0
134	132	131	IV-B	0	0	182	180	179	O	173	0
135	133	132	UL	0	0	183	181	180	C	73	0
136	134	133	IV-U	0	0	184	182	181	C	122	0
137	135	134	BL-U	0	0	185	183	182	M	89	0
138	136	135	I-B-U	0	0	186	184	183	N	142	0
139	137	136		0	0	187	185	184	i	174	0
140	138	137		0	0	188	186	185	L	175	0
141	139	138		0	0	189	187	186	O	176	0
142	140	139		0	0	190	188	187	E	177	0
143	141	140		0	0	191	189	188	O	178	0
144	142	141		0	0	192	190	189	S	179	0
145	143	142		0	0	193	191	190	O	180	0
146	144	143		0	0	194	192	191	.	181	0

## ----- COLLATING SECTION - CONT -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
195	193	192	ä	118	0	243	241	240		0	0
196	194	193	é	127	0	244	242	241		0	0
197	195	194	ô	146	0	245	243	242		0	0
198	196	195	û	156	0	246	244	243		0	0
199	197	196	ä	116	0	247	245	244		0	0
200	198	197	é	125	0	248	246	245		0	0
201	199	198	ô	144	0	249	247	246		0	0
202	200	199	û	154	0	250	248	247		0	0
203	201	200	ä	117	0	251	249	248		0	0
204	202	201	é	126	0	252	250	249		0	0
205	203	202	ô	145	0	253	251	250		0	0
206	204	203	û	155	0	254	252	251		0	0
207	205	204	ä	113	1	255	253	252		0	0
208	206	205	é	128	0	256	254	253		0	0
209	207	206	ô	143	3	257	255	254		0	0
210	208	207	û	153	5	258	256	255		0	0
211	209	208	Ä	67	0						
212	210	209	ï	135	0						
213	211	210	ø	94	0						
214	212	211	Æ	66	0						
215	213	212	ä	115	0						
216	214	213	ï	133	0						
217	215	214	ø	147	0						
218	216	215	æ	114	0						
219	217	216	Ä	65	7						
220	218	217	ï	134	0						
221	219	218	ö	90	9						
222	220	219	ü	100	11						
223	221	220	é	76	0						
224	222	221	ï	136	0						
225	223	222	ø	151	13						
226	224	223		0	0						
227	225	224		0	0						
228	226	225		0	0						
229	227	226		0	0						
230	228	227		0	0						
231	229	228		0	0						
232	230	229		0	0						
233	231	230		0	0						
234	232	231		0	0						
235	233	232		0	0						
236	234	233		0	0						
237	235	234		0	0						
238	236	235		0	0						
239	237	236		0	0						
240	238	237		0	0						
241	239	238		0	0						
242	240	239		0	0						

## ----- UPPERCASE/LOWERCASE SECTION -----

ENTRY #	U/L ENTRY	CHAR CODE	CHAR	UPPER CASE	LOWER CASE	ENTRY #	U/L ENTRY	CHAR CODE	CHAR	UPPER CASE	LOWER CASE
259	1	160				307	49	208	À	À	à
260	2	161	À	À	à	308	50	209	Á	Á	á
261	3	162	Ã	Ã	ã	309	51	210	Â	Â	â
262	4	163	Ô	Ô	ô	310	52	211	Ë	Ë	ë
263	5	164	Ù	Ù	ù	311	53	212	À	À	à
264	6	165	Ä	Ä	ä	312	54	213	Í	Í	í
265	7	166	È	È	è	313	55	214	Ë	Ë	ë
266	8	167	Ò	Ò	ò	314	56	215	Ë	Ë	ë
267	9	168	Û	Û	û	315	57	216	Ä	Ä	ä
268	10	169	¸	¸	¸	316	58	217	Ï	Ï	ï
269	11	170	¸	¸	¸	317	59	218	Û	Û	û
270	12	171	¸	¸	¸	318	60	219	Ü	Ü	ü
271	13	172	¸	¸	¸	319	61	220	É	É	é
272	14	173	È	È	è	320	62	221	Ï	Ï	ï
273	15	174	Ò	Ò	ò	321	63	222	Ë	Ë	ë
274	16	175	Ë	Ë	ë	322	64	223			
275	17	176				323	65	224			
276	18	177	Ä	Ä	ä	324	66	225			
277	19	178	Å	Å	å	325	67	226			
278	20	179	Ö	Ö	ö	326	68	227			
279	21	180	Ç	Ç	ç	327	69	228			
280	22	181	Ç	Ç	ç	328	70	229			
281	23	182	Ñ	Ñ	ñ	329	71	230			
282	24	183	Ñ	Ñ	ñ	330	72	231			
283	25	184	Ï	Ï	ï	331	73	232			
284	26	185	Ç	Ç	ç	332	74	233			
285	27	186	È	È	è	333	75	234			
286	28	187	È	È	è	334	76	235			
287	29	188	È	È	è	335	77	236			
288	30	189	È	È	è	336	78	237			
289	31	190	È	È	è	337	79	238			
290	32	191	¸	¸	¸	338	80	239			
291	33	192	À	À	à	339	81	240			
292	34	193	È	È	è	340	82	241			
293	35	194	Ò	Ò	ò	341	83	242			
294	36	195	Ù	Ù	ù	342	84	243			
295	37	196	À	À	à	343	85	244			
296	38	197	È	È	è	344	86	245			
297	39	198	Ò	Ò	ò	345	87	246			
298	40	199	Ù	Ù	ù	346	88	247			
299	41	200	À	À	à	347	89	248			
300	42	201	È	È	è	348	90	249			
301	43	202	Ò	Ò	ò	349	91	250			
302	44	203	Ù	Ù	ù	350	92	251			
303	45	204	À	À	à	351	93	252			
304	46	205	È	È	è	352	94	253			
305	47	206	Ò	Ò	ò	353	95	254			
306	48	207	Ù	Ù	ù	354	96	255			



## ----- MODE SECTION -----

ENTRY #	MODE ENTRY	TYPE DESCRIPTION
355	1	TWO FOR ONE CHARACTER REPLACEMENT
356	2	TWO FOR ONE CHARACTER REPLACEMENT
357	3	TWO FOR ONE CHARACTER REPLACEMENT
358	4	TWO FOR ONE CHARACTER REPLACEMENT
359	5	TWO FOR ONE CHARACTER REPLACEMENT
360	6	TWO FOR ONE CHARACTER REPLACEMENT
361	7	TWO FOR ONE CHARACTER REPLACEMENT
362	8	TWO FOR ONE CHARACTER REPLACEMENT
363	9	TWO FOR ONE CHARACTER REPLACEMENT
364	10	TWO FOR ONE CHARACTER REPLACEMENT
365	11	TWO FOR ONE CHARACTER REPLACEMENT
366	12	TWO FOR ONE CHARACTER REPLACEMENT
367	13	TWO FOR ONE CHARACTER REPLACEMENT
368	14	TWO FOR ONE CHARACTER REPLACEMENT

||||| TYPE DETAILS |||||

-- DON'T CARES ---- ACCENT PRIORITIES ---- TWO FOR ONE REPLACEMENTS --

ENTRY #	REPLACEMENT	ASCII
1	ä = ae = ae	204
3	ö = oe = oe	206
5	ü = ue = ue	207
7	Å = AE = Ae	216
9	Ö = OE = Oe	218
11	Ü = UE = Ue	219
13	ß = ss = ss	222

-- ONE FOR TWO REPLACEMENTS --

SPANISH TABLE											
---------------	--	--	--	--	--	--	--	--	--	--	--

ENTRY #1 -- LENGTH OF COLLATING TABLE = 374  
 ENTRY #2 -- LENGTH OF MODE TABLE = 20

## ----- COLLATING SECTION -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
3	1	0	NUL	0	0	51	49	48	0	48	0
4	2	1	SOH	1	0	52	50	49	1	49	0
5	3	2	STX	2	0	53	51	50	2	50	0
6	4	3	ETX	3	0	54	52	51	3	51	0
7	5	4	EOT	4	0	55	53	52	4	52	0
8	6	5	ENQ	5	0	56	54	53	5	53	0
9	7	6	ACK	6	0	57	55	54	6	54	0
10	8	7	BEL	7	0	58	56	55	7	55	0
11	9	8	BS	8	0	59	57	56	8	56	0
12	10	9	HT	9	0	60	58	57	9	57	0
13	11	10	LF	10	0	61	59	58	:	58	0
14	12	11	VT	11	0	62	60	59	;	59	0
15	13	12	FF	12	0	63	61	60	<	60	0
16	14	13	CR	13	0	64	62	61	=	61	0
17	15	14	SO	14	0	65	63	62	>	62	0
18	16	15	SI	15	0	66	64	63	?	63	0
19	17	16	DLE	16	0	67	65	64	@	64	0
20	18	17	DC1	17	0	68	66	65	A	65	0
21	19	18	DC2	18	0	69	67	66	B	66	0
22	20	19	DC3	19	0	70	68	67	C	67	1
23	21	20	DC4	20	0	71	69	68	D	69	0
24	22	21	NAK	21	0	72	70	69	E	70	0
25	23	22	SYN	22	0	73	71	70	F	71	0
26	24	23	ETB	23	0	74	72	71	G	72	0
27	25	24	CAN	24	0	75	73	72	H	73	0
28	26	25	EM	25	0	76	74	73	I	74	0
29	27	26	SUB	26	0	77	75	74	J	75	0
30	28	27	ESC	27	0	78	76	75	K	76	0
31	29	28	FS	28	0	79	77	76	L	77	4
32	30	29	GS	29	0	80	78	77	M	79	0
33	31	30	RS	30	0	81	79	78	N	80	0
34	32	31	US	31	0	82	80	79	O	82	0
35	33	32	SP	32	0	83	81	80	P	83	0
36	34	33	!	33	0	84	82	81	Q	84	0
37	35	34	"	34	0	85	83	82	R	85	0
38	36	35	#	35	0	86	84	83	S	86	0
39	37	36	\$	36	0	87	85	84	T	87	0
40	38	37	%	37	0	88	86	85	U	88	0
41	39	38	&	38	0	89	87	86	V	89	0
42	40	39	<	39	0	90	88	87	W	90	0
43	41	40	(	40	0	91	89	88	X	91	0
44	42	41	)	41	0	92	90	89	Y	92	0
45	43	42	*	42	0	93	91	90	Z	93	0
46	44	43	+	43	0	94	92	91	[	94	0
47	45	44	,	44	0	95	93	92	\	95	0
48	46	45	-	45	0	96	94	93	]	100	0
49	47	46	.	46	0	97	95	94	^	101	0
50	48	47	/	47	0	98	96	95	_	102	0

## ----- COLLATING SECTION - CONT -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
99	97	96	`	103	0	147	145	144		0	0
100	98	97	a	104	0	148	146	145		0	0
101	99	98	b	105	0	149	147	146		0	0
102	100	99	c	106	7	150	148	147		0	0
103	101	100	d	108	0	151	149	148		0	0
104	102	101	e	109	0	152	150	149		0	0
105	103	102	f	110	0	153	151	150		0	0
106	104	103	g	111	0	154	152	151		0	0
107	105	104	h	112	0	155	153	152		0	0
108	106	105	i	113	0	156	154	153		0	0
109	107	106	j	114	0	157	155	154		0	0
110	108	107	k	115	0	158	156	155		0	0
111	109	108	l	116	10	159	157	156		0	0
112	110	109	m	118	0	160	158	157		0	0
113	111	110	n	119	0	161	159	158		0	0
114	112	111	o	121	0	162	160	159		0	0
115	113	112	p	122	0	163	161	160		0	0
116	114	113	q	123	0	164	162	161	A	65	0
117	115	114	r	124	0	165	163	162	I	74	0
118	116	115	s	125	0	166	164	163	O	82	0
119	117	116	t	126	0	167	165	164	U	88	0
120	118	117	u	127	0	168	166	165	A	65	0
121	119	118	v	128	0	169	167	166	E	70	0
122	120	119	w	129	0	170	168	167	O	82	0
123	121	120	x	130	0	171	169	168	`	137	0
124	122	121	y	131	0	172	170	169	`	138	0
125	123	122	z	132	0	173	171	170	`	139	0
126	124	123	(	133	0	174	172	171	`	140	0
127	125	124		134	0	175	173	172	`	141	0
128	126	125	)	135	0	176	174	173	E	70	0
129	127	126	~	136	0	177	175	174	O	82	0
130	128	127	DEL	0	0	178	176	175	E	142	0
131	129	128	CLR	0	0	179	177	176	I	143	0
132	130	129	IV	0	0	180	178	177	A	65	0
133	131	130	BL	0	0	181	179	178	A	104	0
134	132	131	IV-B	0	0	182	180	179	O	144	0
135	133	132	UL	0	0	183	181	180	C	67	0
136	134	133	IV-U	0	0	184	182	181	C	106	0
137	135	134	BL-U	0	0	185	183	182	N	81	0
138	136	135	I-B-U	0	0	186	184	183	N	120	0
139	137	136		0	0	187	185	184	i	145	0
140	138	137		0	0	188	186	185	L	146	0
141	139	138		0	0	189	187	186	O	147	0
142	140	139		0	0	190	188	187	E	148	0
143	141	140		0	0	191	189	188	E	149	0
144	142	141		0	0	192	190	189	S	150	0
145	143	142		0	0	193	191	190	2	151	0
146	144	143		0	0	194	192	191	.	152	0

## ----- COLLATING SECTION - CONT -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
195	193	192	ä	104	0	243	241	240		0	0
196	194	193	ë	109	0	244	242	241		0	0
197	195	194	ö	121	0	245	243	242		0	0
198	196	195	û	127	0	246	244	243		0	0
199	197	196	ä	104	0	247	245	244		0	0
200	198	197	ë	109	0	248	246	245		0	0
201	199	198	ö	121	0	249	247	246		0	0
202	200	199	û	127	0	250	248	247		0	0
203	201	200	ä	104	0	251	249	248		0	0
204	202	201	ë	109	0	252	250	249		0	0
205	203	202	ö	121	0	253	251	250		0	0
206	204	203	û	127	0	254	252	251		0	0
207	205	204	ä	104	0	255	253	252		0	0
208	206	205	ë	109	0	256	254	253		0	0
209	207	206	ö	121	0	257	255	254		0	0
210	208	207	û	127	0	258	256	255		0	0
211	209	208	Ä	65	0						
212	210	209	Ï	113	0						
213	211	210	Ø	82	0						
214	212	211	Æ	65	0						
215	213	212	ä	104	0						
216	214	213	ï	113	0						
217	215	214	ø	121	0						
218	216	215	æ	104	0						
219	217	216	Ä	65	0						
220	218	217	Ï	113	0						
221	219	218	Ø	82	0						
222	220	219	Ù	88	0						
223	221	220	é	70	0						
224	222	221	ï	113	0						
225	223	222	Ø	125	13						
226	224	223		0	0						
227	225	224		0	0						
228	226	225		0	0						
229	227	226		0	0						
230	228	227		0	0						
231	229	228		0	0						
232	230	229		0	0						
233	231	230		0	0						
234	232	231		0	0						
235	233	232		0	0						
236	234	233		0	0						
237	235	234		0	0						
238	236	235		0	0						
239	237	236		0	0						
240	238	237		0	0						
241	239	238		0	0						
242	240	239		0	0						

## ----- UPPER CASE      LOWER CASE SECTION -----

ENTRY #	U/L ENTRY	CHAR CODE	CHAR	UPPER CASE	LOWER CASE	ENTRY #	U/L ENTRY	CHAR CODE	CHAR	UPPER CASE	LOWER CASE
259	1	160				307	49	208	À	À	à
260	2	161	Á	Á	á	308	50	209	Í	Í	í
261	3	162	Ī	Ī	ī	309	51	210	Ō	Ō	ō
262	4	163	Ō	Ō	ō	310	52	211	Ĳ	Ĳ	ĳ
263	5	164	Ū	Ū	ū	311	53	212	À	À	à
264	6	165	Ā	Ā	ā	312	54	213	Ī	Ī	ī
265	7	166	Ē	Ē	ē	313	55	214	Ō	Ō	ō
266	8	167	Ō	Ō	ō	314	56	215	Ĳ	Ĳ	ĳ
267	9	168	·	·	·	315	57	216	Ā	Ā	ā
268	10	169	·	·	·	316	58	217	Ī	Ī	ī
269	11	170	·	·	·	317	59	218	Ō	Ō	ō
270	12	171	·	·	·	318	60	219	Ū	Ū	ū
271	13	172	·	·	·	319	61	220	Ē	Ē	ē
272	14	173	Ē	Ē	ē	320	62	221	Ī	Ī	ī
273	15	174	Ō	Ō	ō	321	63	222	Ĳ	Ĳ	ĳ
274	16	175	Ĳ	Ĳ	ĳ	322	64	223			
275	17	176	Ĳ	Ĳ	ĳ	323	65	224			
276	18	177	Ā	Ā	ā	324	66	225			
277	19	178	Ī	Ī	ī	325	67	226			
278	20	179	Ō	Ō	ō	326	68	227			
279	21	180	Ū	Ū	ū	327	69	228			
280	22	181	Ā	Ā	ā	328	70	229			
281	23	182	Ī	Ī	ī	329	71	230			
282	24	183	Ō	Ō	ō	330	72	231			
283	25	184	Ū	Ū	ū	331	73	232			
284	26	185	Ā	Ā	ā	332	74	233			
285	27	186	Ī	Ī	ī	333	75	234			
286	28	187	Ō	Ō	ō	334	76	235			
287	29	188	Ū	Ū	ū	335	77	236			
288	30	189	Ā	Ā	ā	336	78	237			
289	31	190	Ī	Ī	ī	337	79	238			
290	32	191	Ō	Ō	ō	338	80	239			
291	33	192	Ū	Ū	ū	339	81	240			
292	34	193	Ā	Ā	ā	340	82	241			
293	35	194	Ī	Ī	ī	341	83	242			
294	36	195	Ō	Ō	ō	342	84	243			
295	37	196	Ū	Ū	ū	343	85	244			
296	38	197	Ā	Ā	ā	344	86	245			
297	39	198	Ī	Ī	ī	345	87	246			
298	40	199	Ō	Ō	ō	346	88	247			
299	41	200	Ū	Ū	ū	347	89	248			
300	42	201	Ā	Ā	ā	348	90	249			
301	43	202	Ī	Ī	ī	349	91	250			
302	44	203	Ō	Ō	ō	350	92	251			
303	45	204	Ū	Ū	ū	351	93	252			
304	46	205	Ā	Ā	ā	352	94	253			
305	47	206	Ī	Ī	ī	353	95	254			
306	48	207	Ō	Ō	ō	354	96	255			

## ----- MODE SECTION -----

ENTRY #	MODE ENTRY	TYPE DESCRIPTION
355	1	ONE FOR TWO CHARACTER REPLACEMENT
356	2	ONE FOR TWO CHARACTER REPLACEMENT
357	3	* TERMINATOR WORD *
358	4	ONE FOR TWO CHARACTER REPLACEMENT
359	5	ONE FOR TWO CHARACTER REPLACEMENT
360	6	* TERMINATOR WORD *
361	7	ONE FOR TWO CHARACTER REPLACEMENT
362	8	ONE FOR TWO CHARACTER REPLACEMENT
363	9	* TERMINATOR WORD *
364	10	ONE FOR TWO CHARACTER REPLACEMENT
365	11	ONE FOR TWO CHARACTER REPLACEMENT
366	12	* TERMINATOR WORD *
367	13	TWO FOR ONE CHARACTER REPLACEMENT
368	14	TWO FOR ONE CHARACTER REPLACEMENT

||||| TYPE DETAILS |||||

-- DON'T CARES --

-- ACCENT PRIORITIES --

-- TWO FOR ONE REPLACEMENTS --

ENTRY #	REPLACEMENT	ASCII
13	B = ss = ss	222

-- ONE FOR TWO REPLACEMENTS --

ENTRY #	REPLACEMENT AND SEQUENCE NUMBER
1	CH = 68
2	Ch = 68
4	LL = 78
5	Ll = 78
7	cH = 107
8	ch = 107
10	lL = 117
11	ll = 117

SWEDISH TABLE											
---------------	--	--	--	--	--	--	--	--	--	--	--

ENTRY #1 -- LENGTH OF COLLATING TABLE = 374

ENTRY #2 -- LENGTH OF MODE TABLE = 20

## ----- COLLATING SECTION -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
3	1	0	NUL	0	0	51	49	48	0	48	0
4	2	1	SOH	1	0	52	50	49	1	49	0
5	3	2	STX	2	0	53	51	50	2	50	0
6	4	3	ETX	3	0	54	52	51	3	51	0
7	5	4	EOT	4	0	55	53	52	4	52	0
8	6	5	ENO	5	0	56	54	53	5	53	0
9	7	6	ACK	6	0	57	55	54	6	54	0
10	8	7	BEL	7	0	58	56	55	7	55	0
11	9	8	BS	8	0	59	57	56	8	56	0
12	10	9	HT	9	0	60	58	57	9	57	0
13	11	10	LF	10	0	61	59	58	:	58	0
14	12	11	VT	11	0	62	60	59	;	59	0
15	13	12	FF	12	0	63	61	60	<	60	0
16	14	13	CR	13	0	64	62	61	=	61	0
17	15	14	SO	14	0	65	63	62	>	62	0
18	16	15	SI	15	0	66	64	63	?	63	0
19	17	16	DLE	16	0	67	65	64	@	64	0
20	18	17	DC1	17	0	68	66	65	A	65	0
21	19	18	DC2	18	0	69	67	66	B	66	0
22	20	19	DC3	19	0	70	68	67	C	67	0
23	21	20	DC4	20	0	71	69	68	D	68	0
24	22	21	NAK	21	0	72	70	69	E	69	0
25	23	22	SYN	22	0	73	71	70	F	70	0
26	24	23	ETB	23	0	74	72	71	G	71	0
27	25	24	CAN	24	0	75	73	72	H	72	0
28	26	25	EM	25	0	76	74	73	I	73	0
29	27	26	SUB	26	0	77	75	74	J	74	0
30	28	27	ESC	27	0	78	76	75	K	75	0
31	29	28	FS	28	0	79	77	76	L	76	0
32	30	29	GS	29	0	80	78	77	M	77	0
33	31	30	RS	30	0	81	79	78	N	78	0
34	32	31	US	31	0	82	80	79	O	80	0
35	33	32	SP	32	0	83	81	80	P	81	0
36	34	33	!	33	0	84	82	81	Q	82	0
37	35	34	"	34	0	85	83	82	R	83	0
38	36	35	#	35	0	86	84	83	S	84	0
39	37	36	\$	36	0	87	85	84	T	85	0
40	38	37	%	37	0	88	86	85	U	86	0
41	39	38	&	38	0	89	87	86	V	87	0
42	40	39	^	39	0	90	88	87	W	88	0
43	41	40	(	40	0	91	89	88	X	89	0
44	42	41	)	41	0	92	90	89	Y	90	0
45	43	42	*	42	0	93	91	90	Z	91	0
46	44	43	+	43	0	94	92	91	[	111	0
47	45	44	,	44	0	95	93	92	\	112	0
48	46	45	-	45	0	96	94	93	]	113	0
49	47	46	.	46	0	97	95	94	^	114	0
50	48	47	/	47	0	98	96	95	_	115	0

## ----- COLLATING SECTION - CONT -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
99	97	96		116	0	147	145	144		0	0
100	98	97	a	117	0	148	146	145		0	0
101	99	98	b	118	0	149	147	146		0	0
102	100	99	c	119	0	150	148	147		0	0
103	101	100	d	120	0	151	149	148		0	0
104	102	101	e	122	0	152	150	149		0	0
105	103	102	f	123	0	153	151	150		0	0
106	104	103	g	124	0	154	152	151		0	0
107	105	104	h	125	0	155	153	152		0	0
108	106	105	i	126	0	156	154	153		0	0
109	107	106	j	127	0	157	155	154		0	0
110	108	107	k	128	0	158	156	155		0	0
111	109	108	l	129	0	159	157	156		0	0
112	110	109	m	130	0	160	158	157		0	0
113	111	110	n	131	0	161	159	158		0	0
114	112	111	o	132	0	162	160	159		0	0
115	113	112	p	133	0	163	161	160		0	0
116	114	113	q	134	0	164	162	161	A	94	0
117	115	114	r	135	0	165	163	162	i	102	0
118	116	115	s	136	0	166	164	163	o	103	0
119	117	116	t	137	0	167	165	164	U	109	0
120	118	117	u	138	0	168	166	165	A	95	0
121	119	118	v	139	0	169	167	166	E	100	0
122	120	119	w	140	0	170	168	167	o	104	0
123	121	120	x	141	0	171	169	168	.	173	0
124	122	121	y	142	0	172	170	169	,	174	0
125	123	122	z	143	0	173	171	170	^	175	0
126	124	123	(	169	0	174	172	171	-	176	0
127	125	124		170	0	175	173	172	~	177	0
128	126	125	)	171	0	176	174	173	E	101	0
129	127	126	^	172	0	177	175	174	O	105	0
130	128	127	DEL	0	0	178	176	175	E	178	0
131	129	128	CLR	0	0	179	177	176	-	179	0
132	130	129	IV	0	0	180	178	177	A	97	0
133	131	130	BL	0	0	181	179	178	o	150	0
134	132	131	IV-B	0	0	182	180	179	o	180	0
135	133	132	UL	0	0	183	181	180	C	98	0
136	134	133	IV-U	0	0	184	182	181	C	151	0
137	135	134	BL-U	0	0	185	183	182	N	108	0
138	136	135	I-B-U	0	0	186	184	183	R	164	0
139	137	136		0	0	187	185	184	i	181	0
140	138	137		0	0	188	186	185	o	182	0
141	139	138		0	0	189	187	186	o	183	0
142	140	139		0	0	190	188	187	E	184	0
143	141	140		0	0	191	189	188	o	185	0
144	142	141		0	0	192	190	189	o	186	0
145	143	142		0	0	193	191	190	o	187	0
146	144	143		0	0	194	192	191	.	188	0



## ----- COLLATING SECTION - CONT -----

ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR	ENTRY #	COLL. ENTRY	CHAR CODE	CHAR	SEQ #	MODE PTR
195	193	192	á	148	0	243	241	240		0	0
196	194	193	è	153	0	244	242	241		0	0
197	195	194	ò	161	0	245	243	242		0	0
198	196	195	ó	167	0	246	244	243		0	0
199	197	196	â	146	0	247	245	244		0	0
200	198	197	é	122	0	248	246	245		0	0
201	199	198	ó	159	0	249	247	246		0	0
202	200	199	ú	165	0	250	248	247		0	0
203	201	200	â	147	0	251	249	248		0	0
204	202	201	è	152	0	252	250	249		0	0
205	203	202	ò	160	0	253	251	250		0	0
206	204	203	ú	166	0	254	252	251		0	0
207	205	204	â	149	0	255	253	252		0	0
208	206	205	é	154	0	256	254	253		0	0
209	207	206	ó	162	0	257	255	254		0	0
210	208	207	ú	168	0	258	256	255		0	0
211	209	208	á	93	0						
212	210	209	í	157	0						
213	211	210	ó	107	0						
214	212	211	æ	92	0						
215	213	212	â	145	0						
216	214	213	í	155	0						
217	215	214	ø	163	0						
218	216	215	æ	144	0						
219	217	216	á	96	0						
220	218	217	í	156	0						
221	219	218	ó	106	0						
222	220	219	ú	110	0						
223	221	220	é	99	0						
224	222	221	í	158	0						
225	223	222	ó	136	1						
226	224	223		0	0						
227	225	224		0	0						
228	226	225		0	0						
229	227	226		0	0						
230	228	227		0	0						
231	229	228		0	0						
232	230	229		0	0						
233	231	230		0	0						
234	232	231		0	0						
235	233	232		0	0						
236	234	233		0	0						
237	235	234		0	0						
238	236	235		0	0						
239	237	236		0	0						
240	238	237		0	0						
241	239	238		0	0						
242	240	239		0	0						

----- UPPERCASE/LOWERCASE SECTION -----

ENTRY #	U/L ENTRY	CHAR CODE	CHAR	UPPER CASE	LOWER CASE	ENTRY #	U/L ENTRY	CHAR CODE	CHAR	UPPER CASE	LOWER CASE
259	1	160				307	49	208	À	À	à
260	2	161	Á	Á	á	308	50	209	Í	Í	í
261	3	162	Ī	Ī	ī	309	51	210	Ō	Ō	ō
262	4	163	Ō	Ō	ō	310	52	211	Ē	Ē	ē
263	5	164	Ū	Ū	ū	311	53	212	À	À	à
264	6	165	Ā	Ā	ā	312	54	213	Ī	Ī	ī
265	7	166	Ē	Ē	ē	313	55	214	Ō	Ō	ō
266	8	167	Ō	Ō	ō	314	56	215	Ē	Ē	ē
267	9	168	ˆ	ˆ	ˆ	315	57	216	Ā	Ā	ā
268	10	169	ˆ	ˆ	ˆ	316	58	217	Ī	Ī	ī
269	11	170	ˆ	ˆ	ˆ	317	59	218	Ō	Ō	ō
270	12	171	ˆ	ˆ	ˆ	318	60	219	Ū	Ū	ū
271	13	172	ˆ	ˆ	ˆ	319	61	220	É	É	é
272	14	173	Ē	Ē	ē	320	62	221	Ī	Ī	ī
273	15	174	Ō	Ō	ō	321	63	222	Ī	Ī	ī
274	16	175	Ē	Ē	ē	322	64	223			
275	17	176				323	65	224			
276	18	177	Ā	Ā	ā	324	66	225			
277	19	178	Ā	Ā	ā	325	67	226			
278	20	179	°	°	°	326	68	227			
279	21	180	Ç	Ç	ç	327	69	228			
280	22	181	Ç	Ç	ç	328	70	229			
281	23	182	Ñ	Ñ	ñ	329	71	230			
282	24	183	Ñ	Ñ	ñ	330	72	231			
283	25	184	ı	ı	ı	331	73	232			
284	26	185	ı	ı	ı	332	74	233			
285	27	186	ı	ı	ı	333	75	234			
286	28	187	ı	ı	ı	334	76	235			
287	29	188	ı	ı	ı	335	77	236			
288	30	189	ı	ı	ı	336	78	237			
289	31	190	ı	ı	ı	337	79	238			
290	32	191	ı	ı	ı	338	80	239			
291	33	192	À	À	à	339	81	240			
292	34	193	Ē	Ē	ē	340	82	241			
293	35	194	Ō	Ō	ō	341	83	242			
294	36	195	Ū	Ū	ū	342	84	243			
295	37	196	Ā	Ā	ā	343	85	244			
296	38	197	Ē	Ē	ē	344	86	245			
297	39	198	Ō	Ō	ō	345	87	246			
298	40	199	Ū	Ū	ū	346	88	247			
299	41	200	Ā	Ā	ā	347	89	248			
300	42	201	Ē	Ē	ē	348	90	249			
301	43	202	Ō	Ō	ō	349	91	250			
302	44	203	Ū	Ū	ū	350	92	251			
303	45	204	Ā	Ā	ā	351	93	252			
304	46	205	Ē	Ē	ē	352	94	253			
305	47	206	Ō	Ō	ō	353	95	254			
306	48	207	Ū	Ū	ū	354	96	255			

----- MODE SECTION -----

ENTRY #	MODE ENTRY	TYPE DESCRIPTION
355	1	TWO FOR ONE CHARACTER REPLACEMENT
356	2	TWO FOR ONE CHARACTER REPLACEMENT

||||||| TYPE DETAILS |||||

-- DON'T CARES --

-- ACCENT PRIORITIES --

-- TWO FOR ONE REPLACEMENTS --

ENTRY #	REPLACEMENT	ASCII
1	B = ss = ss	222

-- ONE FOR TWO REPLACEMENTS --

# Error Messages

## Mainframe Errors

- 1 Missing ROM or configuration error. Also, check to see if all option ROMs are installed properly.
- 2 Memory overflow; subprogram larger than block of memory. Also check to see if your arrays are too large to fit in memory.
- 3 Line not found or not in current program segment. Check the spelling of line labels and line identifiers.
- 4 Improper return. Branched into the middle of a subroutine.
- 5 Abnormal program termination; no END or STOP statement.
- 6 Improper FOR/NEXT matching.
- 7 Undefined function or subroutine. Check spellings.
- 8 Improper parameter matching. Check the parameter lists in SUB and CALL, and DEF FN and FN statements to see if they match in number and type.
- 9 Improper number of parameters. Check the number of arguments used in an FN or CALL reference.
- 10 String value required.
- 11 Numeric value required.
- 12 Attempt to redeclare variable. Once a variable name has been declared in a DIM, COM, REAL, SHORT or INTEGER statement, it can't be redeclared in that program segment.
- 13 Array dimensions not specified. You must dimension the array, either explicitly or implicitly.
- 14 Multiple OPTION BASE statements or OPTION BASE statement preceded by variable declarative statements.
- 15 Invalid bounds on array dimension or string length in DIM, COM, REAL, SHORT or INTEGER statement. Strings can't be longer than 32 767 characters. The range of array subscripts is -32 767 through 32 767.
- 16 Dimensions are improper or inconsistent; more than 32 767 elements in an array. Check for wrong number of subscripts in an array reference. Check any matrix multiplication for proper sizes.
- 17 Subscript out of range.
- 18 Substring out of range or string too long. Check substring specifiers against length of string.
- 19 Improper value. Check numbers being entered, especially their exponents.
- 20 Integer precision overflow. The range is -32 768 through 32 767.

# Appendix C

- 21 Short precision overflow. Short-precision numbers have six significant digits and an exponent in the range  $-63$  through  $63$ .
- 22 Real precision overflow. Full-precision numbers have twelve significant digits and an exponent in the range  $-99$  through  $99$ .
- 23 Intermediate result overflow.
- 24  $\text{TAN}(n \cdot \pi / 2)$ , when  $n$  is odd.
- 25 Magnitude of argument of  $\text{ASN}$  or  $\text{ACS}$  is greater than  $1$ .
- 26 Zero to negative power.
- 27 Negative base to non-integer power.
- 28  $\text{LOG}$  or  $\text{LGT}$  of negative number.
- 29  $\text{LOG}$  or  $\text{LGT}$  of zero.
- 30  $\text{SQR}$  of negative number.
- 31 Division by zero; or  $X \text{ MOD } Y$  with  $Y = 0$ .
- 32 String does not represent valid number or string response when numeric data required. Check any use of  $\text{VAL}$  function and its argument. Check for correct spelling of variable name.
- 33 Improper argument for  $\text{NUM}$ ,  $\text{CHR}\#$ , or  $\text{RPT}\#$  function.
- 34 Referenced line is not  $\text{IMAGE}$  statement. Check the line identifier in the  $\text{PRINT USING}$  statement.
- 35 Improper format string.
- 36 Out of  $\text{DATA}$ . Make sure  $\text{READ}$  and  $\text{DATA}$  statements correspond. Use  $\text{RESTORE}$  if appropriate.
- 37  $\text{EDIT}$  string longer than 160 characters. Try using a substring.
- 38 I/O function not allowed.  $\text{TYP}$  and other I/O functions aren't allowed in any I/O statement like  $\text{DISP}$  or  $\text{PRINT}$ . Place the value into a variable.
- 39 Function subprogram not allowed. An  $\text{FN}$  reference isn't allowed in any I/O statement, or in redim subscripts. Place the value into a variable.
- 40 Improper replace, delete or  $\text{REN}$  command.  $\text{SUB}$  and  $\text{DEF FN}$  can only be replaced by another  $\text{SUB}$  or  $\text{DEF FN}$ . They can only be deleted if the rest of the corresponding subprogram is deleted. A renumbering may cause out-of-range line numbers if completed, so an error occurs; check increment value.
- 41 First line number greater than second.
- 42 Attempt to replace or delete a busy line or subprogram. Typically, this is caused by trying to delete an input statement that is still requesting values.
- 43 Matrix not square. The dimensions of an identity matrix or of one used to find an inverse or determinant must be the same size.
- 44 Illegal operand in matrix transpose or matrix multiply. The result matrix can't be one of the operands.
- 45 Nested keyboard entry statements.
- 46 No binary in memory for  $\text{STORE BIN}$  or no program in memory for  $\text{SAVE}$ . Check line numbers in  $\text{SAVE}$  against program in memory.
- 47 Subprogram  $\text{COM}$  declaration is not consistent with main program. Check number, type and dimensions of variables.

- 48 Recursion in single-line DEF FN function. Only subprograms can be called recursively.
- 49 Line specified in ON declaration not found.
- 50 File number less than 1 or greater than 10.
- 51 File not currently assigned. Execute an ASSIGN statement for the file, or check the accuracy of the file number used.
- 52 Improper mass storage unit specifier. Check the values of the select code, unit code and controller address.
- 53 Improper file name. A file name can have 1-6 characters and can't contain a colon, quote mark, NULL or CHR\$(255).
- 54 Duplicate file name. Choose another name or PURGE the old one.
- 55 Directory overflow. There is a maximum number of files that a mass storage medium can hold. A file will have to be removed to add another.
- 56 File name is undefined. Check the spelling.
- 57 Mass Storage ROM is missing. Check to see that the ROM is installed properly.
- 58 Improper file type. Use LOAD for PROG files, ASSIGN and GET on DATA files and LOADKEY for KEYS files.
- 59 Physical or logical end-of-file found. Attempting to READ# or PRINT# past the end of the file. Compare the data list to the file size.
- 60 Physical or logical end-of-record found in random mode. Compare the data list to the record size.
- 61 Defined record size is too small for data item. You can either PURGE and RE-CREATE the file with longer records or regroup the data being recorded.
- 62 File is protected or wrong protect code specified. Check to see that the protect code is included and spelled properly.
- 63 The number of physical records is greater than 32 767. That's the limit; use something smaller.
- 64 Medium overflow (out of user storage space). A file can't be set up because there isn't enough space. Use another medium or purge unwanted files.
- 65 Incorrect data type. You can't use GET on a DATA file that doesn't contain a program. Use TYP to find out what kind of data the computer is trying to be read.
- 66 Excessive rejected tracks during a mass storage initialization. The medium can't be initialized. If the medium is a flexible disk, use a different one. If the medium is a hard disc, call your HP Sales and Service Office for assistance, to determine whether there has been a hardware failure.
- 67 Mass storage parameter less than or equal to 0. Check values of variables. Record numbers, record lengths and number of defined records must be positive numbers.
- 68 Invalid line number in GET or LINK operation. Check line numbers. May be trying to LINK to file that doesn't contain a program.
- 69 - 79 See Mass Storage ROM errors.
- 80 Cartridge out or door open. Also check to see if interface is connected properly.

81	Mass storage device failure. Possible power failure.
82	Mass storage device not present. Check mass storage unit specifier.
83	Write protected. Check the write-protection device on the medium or drive.
84	Record not found. There is a bad spot on the medium.
85	Mass storage medium is not initialized.
86	Not a compatible tape cartridge.
87	Record address error; information can't be read. Hardware failure. Check for a dirty read head.
88	Read data error. Hardware failure. Check for a dirty read head.
89	Check read error.
90	Mass storage system error.
91-99	See Mass Storage ROM errors.
100	Item in print using list is string but image specifier is numeric.
101	Item in print using list is numeric but image specifier is string.
102	Numeric field specifier wider than printer width.
103	Item in print using list has no corresponding image specifier.
104	ON KBD statement not allowed in a subprogram.
105-109	Unused
110-113	See Plotter ROM errors.

System Error octal number; octal number

This error indicates a malfunction in the machine's firmware system. Contact your Sales and Service Office.

## I/O Device Errors

Two error messages can occur when attempting to direct an operation to an I/O device that is not ready for use. A printer which is out of paper or no device at a specified select code are examples. The first message that appears is -

I/O ERROR ON SELECT CODE select code

If the condition is not corrected, the machine beeps intermittently and the following message replaces the first -

I/O TIMEOUT ON SELECT CODE select code

The I/O device can be made usable by correcting the error (loading paper, or changing the select code, for example), then executing the READY# command -

READY# select code

This command readies the I/O device and the operation which was attempted is attempted again. The select code must be specified by an integer.

If you get an I/O error on select code 0 and the printer is not out of paper, call your Sales and Service Office.

In some cases, such as an interface which is not connected, READY# for that select code may not solve the I/O error. In this case, **STOP** should be pressed to regain control of the computer. Be sure to turn the power off before inserting an interface. After the problem is remedied, the operation or program can be tried again.

If you get an I/O error and you have an ON KBD statement in effect, you must press **STOP** to gain control of the computer. Otherwise, the READY# command will be trapped by ON KBD.

## Mass Storage ROM Errors

69	Format switch on the disc off. Turn it on.
70	Not a disc interface. Check mass storage unit specifier.
71	Disc interface power off. Turn it on.
72	Incorrect controller address, controller power off, or disc time out. Check mass storage unit specifier; make sure controller is on.
73	Incorrect device type in mass storage unit specifier.
74	Drive missing or power off.
75	Disc system error, type I <sup>1</sup> .
76	Incorrect unit code in mass storage unit specifier.
77	Disc system error, type II <sup>1</sup> .
78-79	Unused
91-99	Unused

## Graphics ROM Errors

110	Plotter type specification not recognized. Check spelling of "GRAPHICS", "9872A" or "INCREMENTAL".
111	Plotter has not been specified. Check select codes.
112	No graphics hardware installed in the System 45B.
113	LIMIT specifications out of range.

<sup>1</sup> See the Mass Storage Techniques Manual.



## I/O ROM Errors

- 114 98036 card improperly configured.
- 115 TDISP not allowed unless peripheral keyboard active.
- 116 TOPEN is active on another select code.
- 150 Improper select code.
- 151 A negative select code was specified that does not match present bus addressing.
- 152 Parity error.
- 153 Either insufficient input data to satisfy enter list or attempt to ENTER from source into source, or enter count exhausted without linefeed.
- 154 Integer overflow, or ENTER count greater than 32 767 bytes or 16 383 words.
- 155 Invalid interface register number. (Can only specify 4-7.)
- 156 Improper expression type in READIO, WRITEIO, or STATUS list.
- 157 No linefeed was found to satisfy “/” ENTER image specifier, or no linefeed record delimiter was found in 512 characters of input.
- 158 Improper image specifier or nesting image specifiers more than 4 levels deep.
- 159 Numeric data was not received for numeric enter list item.
- 160 Repetition of input character more than 32 768 times.
- 161 Attempted to create CONVERT table or EOL sequence for source or destination variable which is locally defined in a subprogram.
- 162 Attempted to delete a nonexistent CONVERT table or EOL sequence.
- 163 I/O error, such as interface card not present, device timeout, interface or peripheral failure, or stop key hit. (Interface FLAG line=0.)
- 164 Transfer type specified is incorrect type for interface card.
- 165 A FHS or DMA transfer with no format specifies a count which exceeds size of variable, or image specifier indicates more characters than will fit in the specified variable.
- 166 A NOFORMAT FHS or DMA type transfer does not start on an odd numbered character position, such as A\$[3].
- 167 Interface status error, TRL Character or an EOI was received on an HP-IB Interface before ENTER list or image specification was satisfied.
- 184 Improper argument for OCTAL or DECIMAL Function.

## Advanced Programming ROM Error Messages

330	Lexical table size exceeds array.size.
331	Improper pointer array*.
332	Non-existent dimension specified in MAT REORDER.
333	Pointer array contains out-of-range subscript value.
334	Pointer array length does not equal number of records.
335	Pointer array is not one-dimensional
336	Number of records (plus twice the number of secondary keys plus twice the number of substrings) exceeds 16 383.
337	Subscript extends beyond dimensioned maximum length.
338	Subscript out-of-range in key specifier.
339	Starting location is an out-of-range subscript value.
340	Lexical table is too small to include all characters.
341	Main lexical table length plus mode section length does not equal specified table length.
342	Array is not one-dimensioned or is not integer.
343	Lexical mode section pointer out-of-range.
344	Lexical table length exceeds 16 383.

\* This error occurs when data is lost in the process of reordering the array. If this error does not occur, it does not necessarily imply that the pointer array contains a permutation.

## Subject Index

### a

Accent Priority .....	51
Advanced Programming ROM Error Messages .....	93
Appendices:	
A - Statement Summary .....	43
B - User-Defined Lexical Order .....	45
C - System 45B Error Messages .....	87
Array Key .....	7
Array Record .....	6
Array Structure and Terminology .....	6
Array, Pointer .....	13
Ascending Order Sorting .....	11
ASCII .....	34,38,59,62
ASCII Character Codes .....	59
ASCII Table .....	62
Asterisk .....	7

### b

Brackets .....	4
----------------	---

### c

Chapter 1 .....	2
Chapter 2 .....	6
Chapter 3 .....	38
Collating Section .....	48
Collating Sequence .....	34,47
Condition .....	30

### d

Data Manipulation .....	6
Default Sorting Order .....	11
Descending Order Sorting .....	11,18
DEScending Parameter .....	11,18
Dimension .....	6
Dimension Specifier .....	24
Dissimilar Characters .....	14
"Don't Care" Characters .....	58
Dot Matrix .....	4
Duplicate Numbers .....	26

### e

Error Messages:	
Advanced Programming .....	93
Mainframe .....	87
Extended Character Sets .....	37

### f

FRENCH .....	38
French Table .....	67

### g

General Information .....	2
GERMAN .....	38
German Table .....	72

### i

I/O Device Errors .....	90
Identical Data .....	15
Indeterminate Sorting .....	16
Information, General .....	2
INTEGER Mode Sorting .....	11
Introductions:	
Chapter 1 .....	2
Chapter 2 .....	6
Chapter 3 .....	38

### k

Key .....	7
Key Specifier .....	7
Key Specifier, Secondary .....	15
Key String .....	21
Key, Array .....	7

**l**

LEX	39
LEX Function	18
Lexical	11,15
Lexical Order	11
Lexical Order Designator	38
LEXICAL ORDER IS Statement	38
Lexical Order Tables	45,61,62,67,72,77,82
Lexical Ordering	18
Lexical Table Organization	46
Lexical Tables Cartridge	3
Lexical Table Sections:	
Collating Section	48
Mode Section	51
Uppercase/Lowercase Section	50
Local Language Collating Sequences	38
Location Specifier	30,31
Lowercase Function	40
LWC\$	40

**m**

Mainframe Errors	87
Manual Syntax	4
MAT REORDER Statement	27
MAT SEARCH Statement	30
MAT SORT Statement	11

**n**

Null String	19
Numeric Array Searching	30
Numeric Data Sorting	11

**o**

Object Array	25
One-for-Two Character Replacement	53
Option Base 1	6
Out-of-Range Numbers	26

**p**

Parallel Arrays	28
Pointer Array	13

**r**

REAL Mode Sorting	11
Record	7
Record, Array	7
Relational Operators	30
Return Variable	30
Reverse Lexical	11
ROM Installation	2
Roman Extension Character Codes	60

**s**

Search Condition	30
Search Variable	30
Searching Numeric Arrays	30
Searching String Arrays	34
Secondary Key Specifier	15
Sequences, Collating	47
Sorting	11
Sorting Numeric Data	11
Sorting String Data	18
Source Array	11,18,30
Spanish Table	77
SPANSH(Spanish)	38
Specifier, Dimension	24
Specifier, Key	7
Specifier, Substring	21
STANDARD Designator	38
STANDARD Lexical Order	38
Starting Address Specifier	30
Statement Summary	43
String Array Searching	34
String Comparisons	39
String Data Sorting	18
String Expression	39
Structure, Array	6
Substring Specifier	21
Swedish Table	82
SWEDSH(Swedish)	38
System 45B Advanced	
Programming ROM	2
System 45B Error Messages	87

**t**

The LEX Function	39
The Null String	19
Transformations, Uppercase and	
Lowercase	40
Two-for-One Character Replacement	57

## u

Unsatisfied Search .....	32
UPC\$ .....	40
Uppercase Function .....	40
Uppercase/Lowercase Section .....	50
User-Defined Lexical Order .....	45

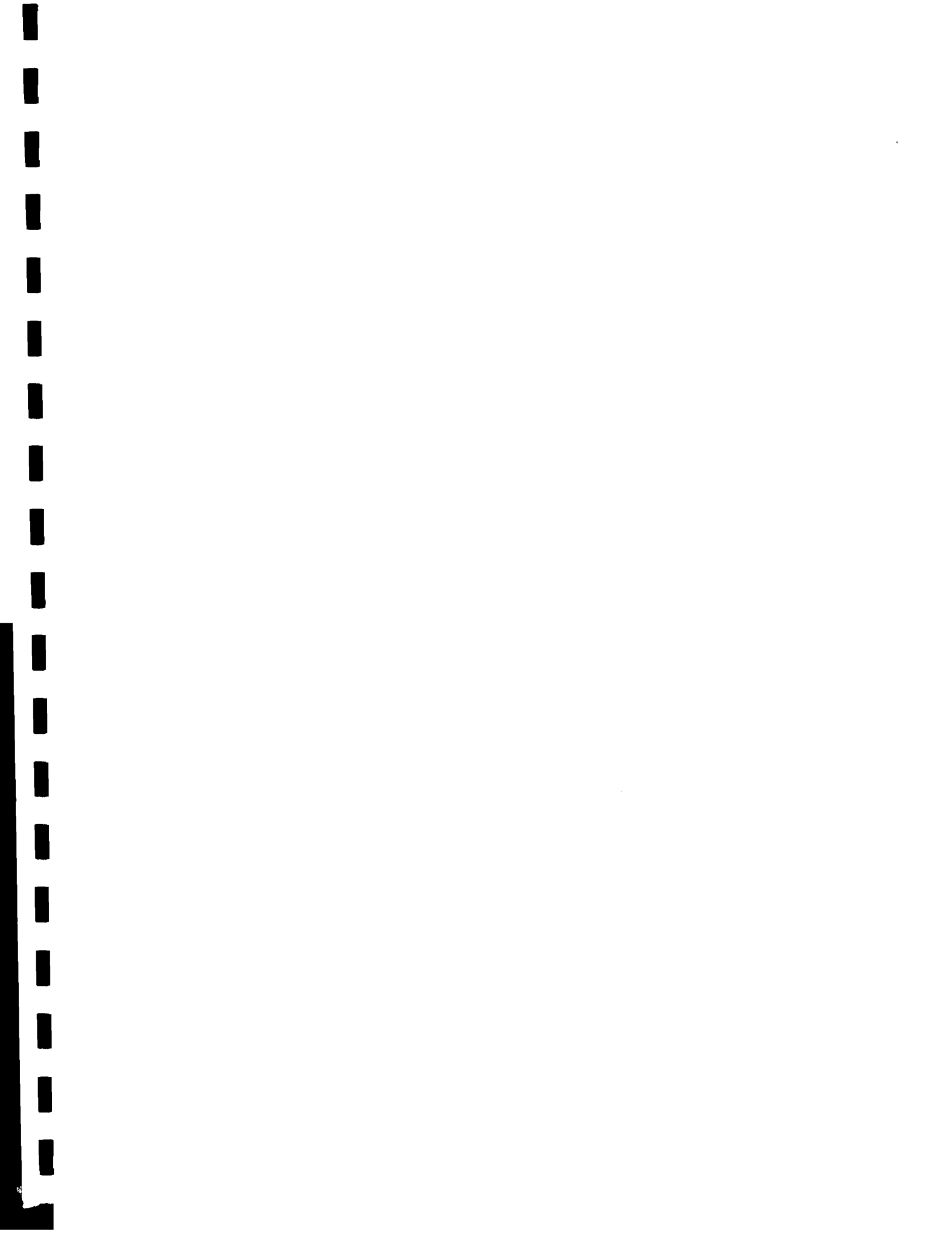
## v

Variable .....	30
----------------	----



- AUSTRIA** Hewlett-Packard Gas m B H Handelskai 52 P.O. Box 7 A-1205 Vienna Tel: 351621-27 Cable: HEWPAK Vienna Fax: 5923 Hewpak a
- BAHRAIN** Al Hamoya Trading and Contracting P.O. Box 20074 Manama Tel: 259978, 259958 Tel: 8895 KALDIA GJ
- BELGIUM** Hewlett-Packard Benlux S.A. Av. Avenue du Col-Vert, 1 (Groenkraaglaan) B-1170 Brussels Tel: (02) 660 50 50 Cable: PALOBEN Brussels Tel: 23-494 paloben bru
- CYPRUS** Kypricos 19 Gregorios Xenopoulos Street P.O. Box 1152 Nicosia Tel: 45628/29 Cable: Kypricos Pandehis Tel: 3018
- CZECHOSLOVAKIA** Vyojovka a Provozni Zakladna Vyzkumnych Ustavu v Bechovicich CSSR-25097 Bechovice u Prahy Tel: 89 93 41 25 Tel: 121333 Institute of Medical Bionics Vyskumny Ustav Lekarsky Bioniky Jedlova 5 CS-88346 Bratislava-Kramare Tel: 4251 Tel: 932229
- DDR** Entwicklungslabor der TU Dresden Forschungsinstitut Meinsberg DDR-7305 Waldheim/Meinsberg Tel: 37 867 Export Contact AG Zurich Guenther Forberg Scheffelstrasse 15 1040 Berlin Tel: 42 74 12 Tel: 111889
- DENMARK** Hewlett-Packard A/S Datavej 52 DK-3460 Birkerød Tel: (02) 81 66 40 Cable: HEWPAK AS Tel: 37499 hps dk Hewlett-Packard A/S Nørvæn 1 DK-8600 Silkeborg Tel: (06) 82 71 86 Tel: 37409 hps dk Cable: HEWPAK AS
- EGYPT** I.E.A. International Engineering Associates 24 Hussein Hegazi Street Kass-el-Aini Cairo Tel: 23 829 Tel: 93830 Cable: INTENGASSO
- SAMITRO** Sami Amr Trading Office 18 Abdel Aziz Gawah Abidine-Cairo Tel: 24932 Cable: SAMITRO CAIRO
- ALABAMA** P.O. Box 4207 8290 Whitesburg Dr. Huntsville 35802 Tel: (205) 981-4591 8533 E. Rouboux Blvd. Birmingham 35208 Tel: (205) 836-2203/2
- ARIZONA** 2336 E. Magnolia St. Phoenix 85034 Tel: (602) 244-1361 2424 East Aragon Rd. Tucson 85706 Tel: (602) 889-4861
- ARKANSAS** Medical Service Only P.O. Box 14865 Little Rock 72215 Tel: (501) 376-1844
- CALIFORNIA** 1579 W. Shaw Ave. Fresno 93711 Tel: (209) 274-0562 1430 East Orangeflorpe Ave. Fullerton 92631 Tel: (714) 870-1000 3939 Lankershim Boulevard North Hollywood 91604 Tel: (213) 877-1282 Tel: (213) 499-2871 5400 West Rosecrans Blvd. P.O. Box 91205 World Way Postal Center Los Angeles 90009 Tel: (213) 776-7500 Tel: (213) 910-325-6608 \*Los Angeles Tel: (213) 776-7500 3003 Scott Boulevard Santa Clara 95050 Tel: (408) 968-7000 \*Ridgeway Tel: (414) 448-8165 646 W North Market Blvd. Sacramento 95834 Tel: (916) 929-7222
- FLORIDA** P.O. Box 24210 2727 N.W. 62nd Street Ft. Lauderdale 33309 Tel: (305) 973-2800 4428 Emerson Street Unit 103 Jacksonville 32207 Tel: (904) 725-6333 P.O. Box 13910 6177 Lake Eleanor Dr. Orlando 32809 Tel: (305) 859-2900 P.D. Box 12826 Suite 5, 9ldg I Office Park North Pensacola 32575 Tel: (904) 476-6272
- GEORGIA** P.O. Box 105005 450 Interstate North Parkway Atlanta 30348 Tel: (404) 955-1500 Medical Service Only \*Augusta 30903 P.O. Box 1044736-0592 P.O. Box 2103 1172 N. Davis Drive Warner Robins 31098 Tel: (912) 822-0449
- HAWAII** 2875 So. King Street Honolulu 96826 Tel: (808) 955-4455
- 9606 Aero Drive P.O. Box 23333 San Diego 92123 Tel: (714) 279-3200 \*Tarzana Tel: (213) 705-3344**
- COLORADO** 5600 OTI Parkway Englewood 80110 Tel: (303) 771-3455
- CONNECTICUT** 12 Lunar Drive New Haven 06525 Tel: (203) 389-8551 Tel: (203) 389-8551 Tel: (203) 389-8551
- FLORIDA** P.O. Box 24210 2727 N.W. 62nd Street Ft. Lauderdale 33309 Tel: (305) 973-2800 4428 Emerson Street Unit 103 Jacksonville 32207 Tel: (904) 725-6333 P.O. Box 13910 6177 Lake Eleanor Dr. Orlando 32809 Tel: (305) 859-2900 P.D. Box 12826 Suite 5, 9ldg I Office Park North Pensacola 32575 Tel: (904) 476-6272
- GEORGIA** P.O. Box 105005 450 Interstate North Parkway Atlanta 30348 Tel: (404) 955-1500 Medical Service Only \*Augusta 30903 P.O. Box 1044736-0592 P.O. Box 2103 1172 N. Davis Drive Warner Robins 31098 Tel: (912) 822-0449
- HAWAII** 2875 So. King Street Honolulu 96826 Tel: (808) 955-4455
- 9606 Aero Drive P.O. Box 23333 San Diego 92123 Tel: (714) 279-3200 \*Tarzana Tel: (213) 705-3344**
- COLORADO** 5600 OTI Parkway Englewood 80110 Tel: (303) 771-3455
- CONNECTICUT** 12 Lunar Drive New Haven 06525 Tel: (203) 389-8551 Tel: (203) 389-8551 Tel: (203) 389-8551
- FLORIDA** P.O. Box 24210 2727 N.W. 62nd Street Ft. Lauderdale 33309 Tel: (305) 973-2800 4428 Emerson Street Unit 103 Jacksonville 32207 Tel: (904) 725-6333 P.O. Box 13910 6177 Lake Eleanor Dr. Orlando 32809 Tel: (305) 859-2900 P.D. Box 12826 Suite 5, 9ldg I Office Park North Pensacola 32575 Tel: (904) 476-6272
- GEORGIA** P.O. Box 105005 450 Interstate North Parkway Atlanta 30348 Tel: (404) 955-1500 Medical Service Only \*Augusta 30903 P.O. Box 1044736-0592 P.O. Box 2103 1172 N. Davis Drive Warner Robins 31098 Tel: (912) 822-0449
- HAWAII** 2875 So. King Street Honolulu 96826 Tel: (808) 955-4455
- 9606 Aero Drive P.O. Box 23333 San Diego 92123 Tel: (714) 279-3200 \*Tarzana Tel: (213) 705-3344**
- COLORADO** 5600 OTI Parkway Englewood 80110 Tel: (303) 771-3455
- CONNECTICUT** 12 Lunar Drive New Haven 06525 Tel: (203) 389-8551 Tel: (203) 389-8551 Tel: (203) 389-8551
- FLORIDA** P.O. Box 24210 2727 N.W. 62nd Street Ft. Lauderdale 33309 Tel: (305) 973-2800 4428 Emerson Street Unit 103 Jacksonville 32207 Tel: (904) 725-6333 P.O. Box 13910 6177 Lake Eleanor Dr. Orlando 32809 Tel: (305) 859-2900 P.D. Box 12826 Suite 5, 9ldg I Office Park North Pensacola 32575 Tel: (904) 476-6272
- GEORGIA** P.O. Box 105005 450 Interstate North Parkway Atlanta 30348 Tel: (404) 955-1500 Medical Service Only \*Augusta 30903 P.O. Box 1044736-0592 P.O. Box 2103 1172 N. Davis Drive Warner Robins 31098 Tel: (912) 822-0449
- HAWAII** 2875 So. King Street Honolulu 96826 Tel: (808) 955-4455
- 9606 Aero Drive P.O. Box 23333 San Diego 92123 Tel: (714) 279-3200 \*Tarzana Tel: (213) 705-3344**
- COLORADO** 5600 OTI Parkway Englewood 80110 Tel: (303) 771-3455
- CONNECTICUT** 12 Lunar Drive New Haven 06525 Tel: (203) 389-8551 Tel: (203) 389-8551 Tel: (203) 389-8551
- FLORIDA** P.O. Box 24210 2727 N.W. 62nd Street Ft. Lauderdale 33309 Tel: (305) 973-2800 4428 Emerson Street Unit 103 Jacksonville 32207 Tel: (904) 725-6333 P.O. Box 13910 6177 Lake Eleanor Dr. Orlando 32809 Tel: (305) 859-2900 P.D. Box 12826 Suite 5, 9ldg I Office Park North Pensacola 32575 Tel: (904) 476-6272
- GEORGIA** P.O. Box 105005 450 Interstate North Parkway Atlanta 30348 Tel: (404) 955-1500 Medical Service Only \*Augusta 30903 P.O. Box 1044736-0592 P.O. Box 2103 1172 N. Davis Drive Warner Robins 31098 Tel: (912) 822-0449
- HAWAII** 2875 So. King Street Honolulu 96826 Tel: (808) 955-4455
- 9606 Aero Drive P.O. Box 23333 San Diego 92123 Tel: (714) 279-3200 \*Tarzana Tel: (213) 705-3344**
- COLORADO** 5600 OTI Parkway Englewood 80110 Tel: (303) 771-3455
- CONNECTICUT** 12 Lunar Drive New Haven 06525 Tel: (203) 389-8551 Tel: (203) 389-8551 Tel: (203) 389-8551
- FLORIDA** P.O. Box 24210 2727 N.W. 62nd Street Ft. Lauderdale 33309 Tel: (305) 973-2800 4428 Emerson Street Unit 103 Jacksonville 32207 Tel: (904) 725-6333 P.O. Box 13910 6177 Lake Eleanor Dr. Orlando 32809 Tel: (305) 859-2900 P.D. Box 12826 Suite 5, 9ldg I Office Park North Pensacola 32575 Tel: (904) 476-6272
- GEORGIA** P.O. Box 105005 450 Interstate North Parkway Atlanta 30348 Tel: (404) 955-1500 Medical Service Only \*Augusta 30903 P.O. Box 1044736-0592 P.O. Box 2103 1172 N. Davis Drive Warner Robins 31098 Tel: (912) 822-0449
- HAWAII** 2875 So. King Street Honolulu 96826 Tel: (808) 955-4455

- ITALY** Hewlett-Packard Italiana S.p.A. Via C. Vittorino, 9 20083 Carraro Tel: (059) 691991 Fax: (059) 691991 Hewlett-Packard Italiana S.p.A. Via Vittorino, 9 20083 Carraro Tel: (059) 691991
- POLAND** Biuro Informacji Technicznej Hewlett-Packard Ul. Stawki 2, 6P 00-242 Warszawa Tel: 32 35 88 09 87 43 Tel: 91 24 53 hps pl
- PORTUGAL** Telecra-Empresa Técnica de Equipamentos Electricos S.A.r.l. Rua Rodrigo da Fonseca 103 P.O. Box 2531 P. Lisboa Tel: (19) 68 80 72 Cable: TELETRA Lisbon Tel: 12588 Medical only Mündster InterCambio Mundial de Comercio S.A.r.l. P.O. Box 2761 Avenida Antonio Augusto de Aguiar 138 P. Lisboa Tel: (01) 53 21 317 Cable: HPAG CH Tel: 18691 munter p Cable: INTERCAMBIO Lisbon
- QATAR** Nasser Trading & Contracting P.O. Box 1563 Doha Tel: 221720 Tel: 4439 NASSER Cable: NASSER
- RUMANIA** Hewlett-Packard Reprzentanta Bd. N. Balcescu 16 Bucurorestel Tel: 15 80 231/3 68 85 Tel: 10440 I.I.R.U.C. Intreprinderea Pentru Intretinerea Si Repararea Utilajelor de Calcul B-d. Prof. Dimitrie Pompei 6 Bucurorestel-Sectorul 6 Tel: 66-20-70, 66-24-40, 66-67-95 Tel: 118
- SAUDI ARABIA** Modern Electronic Establishment (Head Office) P.O. Box 1228, Baghdadiah Street Jeddah Tel: 27 796 Tel: 40035 Cable: ELECTA JEDDAH Modern Electronic Establishment (Branch) P.O. Box 2728 Riyadh Tel: 82596,8232 Cable: ALFOUCO Modern Electronic Establishment (Branch) P.O. Box 193 Al-Khobar Tel: 44478-44813
- SPAIN** Hewlett-Packard Española, S.A. Calle Jerez 3 E-Madrid 16 Tel: (1) 456 28 00 (10 lines) Tel: 52515 hps Hewlett-Packard Española S.A. Colonia Miraflores Edificio Juban \*C/ Costa Brava, 13 Madrid 34 Hewlett-Packard Española, S.A. Milanesado 21-23 E-Barcelona 17 Tel: (3) 203 6200 (5 lines)
- NORWAY** Hewlett-Packard Norge A/S Osterdalen 18 P.O. Box 34 1345 Osterås Tel: (02) 17111 80 Tel: (02) 17111 2081
- NETHERLANDS** Hewlett-Packard Benlux N.V. Van Heuven Goedhartlaan 121 P.O. Box 667 NL-Amstelveen 1134 Tel: (020) 47 20 21
- NORWAY** Hewlett-Packard Norge A/S Osterdalen 18 P.O. Box 34 1345 Osterås Tel: (02) 17111 80 Tel: (02) 17111 2081
- NEW YORK** 6 Automation Lane Computer Park Albany 12205 Tel: (518) 458-1550 TWX: 710-444-4961 650 Perinton Hill Office Park Fairport 14450 Tel: (518) 223-9950 TWX: 510-653-0092 1021 8th Avenue King of Prussia Industrial Park King of Prussia 19406 Tel: (610) 265-7000 TWX: 640-890-2670
- PENNSYLVANIA** 3027 Vantage Dr. Director's Plaza Chester 36311 Tel: (610) 346-8370 \*Meachville Medical Service only Tel: (610) 346-8370
- TENNESSEE** 8914 Kingston Pike Knoxville 37922 Tel: (615) 323-0522 3027 Vantage Dr. Director's Plaza Chester 36311 Tel: (610) 346-8370
- TEXAS** 4171 North Meachville Suite C110 El Paso 79902 Tel: (915) 833-3555 P.O. Box 1270 201 E. Arapahoe Rd. Richardson 75080 Tel: (214) 231-8100
- UTAH** 2160 South 3270 West Street Salt Lake City 84119 Tel: (801) 972-4711
- VIRGINIA** P.O. Box 12778 Norfolk 23502 Tel: (804) 450-2871 P.O. Box 9669 2914 Hungary Springs Road Richmond 23228 Tel: (804) 285-3431
- WASHINGTON** Bellefield Office Pk. 1203 114th Ave. S.E. Bellevue 98005 Tel: (206) 454-3971 TWX 910-443-2448 P.O. Box 4010 Spokane 99202 Tel: (509) 355-0864 \*WEST VIRGINIA Medical/Analytical Only Charleston 25311 Tel: (604) 345-1640 WISCONSIN 9004 West Lincoln Ave. West Allie 53227 Tel: (414) 541-0555
- ANAKA** Tel: 23 03 09 - 17 80 26 Annex: 42578 OZEK Tel: Cable: OZUYUREK ANAKA
- UNITED STATES (Head Office)** P.O. Box 1641 Sharjah Tel: 354121-3 Extension 4787 Hewlett-Packard Office P.O. Box 2771 Abu Dhabi Tel: 331370 1
- UNITED KINGDOM** Hewlett-Packard Ltd King Street Lane Winchester, Wokingham Berks RG1 1 SAR Tel: (0734) 784774 Tel: 8471784
- Hewlett-Packard Ltd** Lygon Court Hewlett-Packard House Navigation Road Cheshire WA14 1HU Tel: (061) 928 8422 Tel: 668066
- Hewlett-Packard Ltd** Lygon Court Hewlett-Packard House Navigation Road Cheshire WA14 1HU Tel: (061) 928 8422 Tel: 668066
- Hewlett-Packard Ltd** Lygon Court Hewlett-Packard House Navigation Road Cheshire WA14 1HU Tel: (061) 928 8422 Tel: 668066
- Hewlett-Packard Ltd** Lygon Court Hewlett-Packard House Navigation Road Cheshire WA14 1HU Tel: (061) 928 8422 Tel: 668066





Printed in U.S.A.  
May 1, 1979