## Hardware Internal Design

## Specification

For the HP-71


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Hewlett-Packard -- Portable Computer Division
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## Corvallis, Oregon

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#### Abstract

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# HP-71 Hardware IDS - Detailed Design Description 

8 SYSTEM DIAGRAMS
9 SYSTEM ELECTRICAL SPECIFICATION


The HP-71 mainframe chip set is composed of the 4 custom CMOS integrated circuits shown below:

CHIP SET:
--------
1 1LF2 CPU
3 1LF3 Display Driver (@ . 5 K bytes RAM/chip)
1 1LG7 ROM hybrid (@ 4 chips/hybrid and 16K bytes ROM/chip)
4 1LG8 RAM hybrids (@ 4 chips/hybrid and 1 K bytes RAM/chip)

The CPU and display drivers are mounted on the underside of the keyboard/LCD PC board. The ROM hybrid and 4 RAM hybrids are mounted on the I/O PC board. Connection to the 4 RAM/ROM ports in the front of the HP-71 and the HP-IL port in the back of the machine is made via the I/O PCB. The card reader port connects to the keyboard PC board.

The CPU (1LF2) is an arithmetic oriented microprocessor capable of both HEX and BCD arithmetic. The 1LF2 operates on a fully multiplexed 4 -bit bus with a 512 K byte address space and an internal word size of 64 bits. The CPU has a general purpose Input Register ( 16 lines) with interrupt capability and a general purpose Output Register (12 lines). The Input and Output Registers are used to form the HP-71 keyboard matrix. The Output Register is also used to drive the daisy-chain signal to each port and to drive the piezo-electric beeper used in the system. External circuitry provides the capability of loud or soft beep.

The three display drivers (1LF3) drive the 8 row by 136 column (including annunciators) display. One of these 3 1LF3s acts as the master. The master generates the display clock, the voltage reference, and the display-on signal used by the other two display chips (slaves). There is a 4 -bit writeable register on the master that controls the display contrast by adjusting the value of the voltage reference signal. Each display chip supplies . 5 K bytes of hard-configured RAM and a 24-bit timer (crystal-oscillator controlled) with $1 / 512$ second resolution. The timers are used by the operating system to implement the clock system. The master display chip also provides a low supply voltage sensor that is used by the operating system to indicate low-battery.

Each 1LG8 chip provides 1 K bytes of soft configured RAM. The four 4 -chip RAM hybrids in combination with the 1.5 K bytes of display driver RAM provide a system total of 17.5 K bytes. Plug-in RAM modules contain one 4 -chip hybrid totalling $4 K$ bytes of RAM. Each 1 LG7 chip provides 16 K bytes of ROM. The 64 K operating system is stored in one 4 -chip ROM hybrid that is hard addressed beginning at 00000. Plug-in ROM modules contain 1 to 4 hard or soft configured ROMs. For information on hard and soft address configuration see section 2.2.

The HP-71 system supports a low-power shutdown mode. The system can be shutdown with the display either on or off. For information on power down or wake up see section 2.4 .

The system is powered from either the unregulated battery voltage or the regulated AC adapter voltage. A 470 uF capacitor supplies "keep-alive" power to the system while batteries are being replaced with the system in shutdown mode.


This section describes the HP-71 system bus structure, protocol, and timing. Also included is a description of the bus commands, device addressing, and the power down and wake up characteristics.

### 2.1 Bus Structure

The HP-71 bus consists of 8 lines including:
4 BUS[0:3] - data lines, driven by the CPU or system devices.
$1{ }^{*}$ STROBE (*STR) - driven by the CPU.
1 *COMMAND-DATA ( ${ }^{*} \mathrm{CD}$ ) - usually driven by the CPU.
1 VDD
1 GROUND

Plug-ins may have two additional lines:
1 DAISYIN - (DIN) input to device
1 DAISYOUT - (DOUT) output from device, may be tied to next device's DIN

NOTE : A '* before a signal name denotes negative true logic.

### 2.1.1 General Protocol

*STR is driven by the CPU and serves to synchronjze all bus transfers. It is not a true system clock since it can remain inactive for several cycles while the CPU is performing internal manipulations. All address, data, and bus commands are ransferred on the BUS[0:3]. The BUS[0:3] is driven during "STR low by either the CPU or the device the CPU is accessing. This data is latched either by the CPU during ${ }^{*}$ STR low or by the receiving device on the rising edge of ${ }^{*}$ STR.

All bus operations are initiated by the CPU. The CPU starts a specific transfer on the bus by driving the *CD line low before ${ }^{*}$ STR goes low. While ${ }^{*} \mathrm{CD}$ and ${ }^{* S T R}$ are low the CPU drives a bus command on the BUS[0:3] and all devices in the system latch the command on the rising edge of ${ }^{*}$ STR. This strobe is referred to as
a command strobe. The bus command issued during a command strobe specifies the operation that is to be performed on each succeeding "STR until another bus command is issued. At all times when data or address is being transferred ${ }^{*} \mathrm{CD}$ is held high. A strobe issued while "CD is high is referred to as a a data strobe.

### 2.1.2 Bus Commands

The bus commands are:

| 0 | NOP | All devices ignore ${ }^{*}$ STR until a new command is loaded. |
| :---: | :---: | :---: |
| 1 | ID | The unconfigured device that has its DAISY- |
|  |  | IN high sends its 5-nibble ID on the |
|  |  | following data strokes, starting with the |
|  |  | low- order nibble of the ID. |
| 2 | PC READ | (PC)->BUS or read using the Program Counter |
|  |  | (PC). The device addressed by its program |
|  |  | counter sends data pointed to by its local |
|  |  | program counter on each following data |
|  |  | strobe and all devices increment their local |
|  |  | program counters once each data strobe. A |
|  |  | dummy strobe immediately follows the |
|  |  | issuance of this bus command (see subsection |
|  |  | 2.1.4). |
| 3 | DP READ | (DP)->BUS or read using the Data Pointer |
|  |  | (DP). The device addressed by its data |
|  |  | pointer sends data pointed to by its local |
|  |  | data pointer on each following data strobe |
|  |  | and all devices increment their local data |
|  |  | pointers once each data strobe. A dummy |
|  |  | strobe immediately follows the issuance of |
|  |  | this bus command (see subsection 2.1.4). |
| 4 | PC WRITE | BUS->(PC) or write using PC. The device |
|  |  | addressed by its program counter loads the |
|  |  | data on the following data strobes into the |
|  |  | location pointed to by its local program |
|  |  | counter and all devices increment their |
|  |  | local program counter once each data strobe. |
| 5 | DP WRITE | BUS->(DP) or write using DP. The device |
|  |  | addressed by its data pointer loads the data |
|  |  | on the following data strobes into the |
|  |  | location pointed to by its local data |
|  |  | pointer and all devices increment their |
|  |  | local data pointer once each data strobe. |


| 6 LOAD PC | BUS->PC or load PC. All devices load the <br> data on following 5 data strobes into their |
| :--- | :--- |
| local program counter, starting with the |  |
| low-order nibble. After all 5 nibbles are |  |
| transferred the command code is |  |
| automatically changed to a 2 , PC READ (see |  |
| subsection 2.1.3). |  |

### 2.1.3 Command Auto-Switch

There exists one special ase in which all devices change their current bus command. This is called 'auto "switch' and occurs following the load of either the PC or the DP. On the rising edge of *STR after the 5th nibble of address has been loaded all devices clear bit 2 of their command latch changing the bus command from either a LOAD PC to PC READ or LOAD DP to DP READ.

### 2.1.4 Dummy Strobe

Immediately following a PC READ bus command, a DP READ bus command, and a command auto-switch the CPU issues a 'dummy strobe'. This dummy strobe appears as a data strobe except that no data is transferred during this period and devices do not insrement their local address registers. The dummy strobe provides memory devices a full strobe cycle for the first access and therefore allows data pipelining.

### 2.2 Addressing

Each device on the HP-71 bus has two 20-bit address registers; a local program counter (PC) and a local data pointer (DP). Each device is also either hard addressed at a specific address (hard configured) or capable of being dynamically located within the address space (soft configured). A device only responds to data reads and writes if its local address register (PC or DP depending on the read or write command) is within its configured address space.

The HP-71 operating system allows soft configured devices to have address spaces ranging in size from 8 bytes to 128 K bytes. All devices are configured such that the upper-order bits of the local address register can be compared with the upper-order bits of the device's configuration register (hard or soft). If these bits are j.dentical, the device has an address match and will respond to read and write commands (and the unconfigure and BUSCC commands if applicable). Each device with a given address space size compares a given number of the upper-order bits of address. For example, a device with an address space size of 1 K bytes or 2 K nibbles requires 11 bits of address leaving the upper 9 bits for its configuration register.

### 2.2.1 Soft Configuration

A soft configured device powers up unconfigured. When unconfigured a device responds only to the ID and CONFIGURE commands and drives its DAISYOUT low. A device's ID code is used to identify the device before it is configured. If a soft configured device is unconfigured and has its DAISYIN line high, it sources its 5 -nibble ID code starting with the low-order nibble on the 5 data strobes immediately following the issuance of an ID command (no dummy strobe is issued).

The 5-nibble ID code contains information on the device type and the address space required by the device as defined below:

| NIBBLE $0:$ | Determines the size of the address space and is |
| ---: | :--- |
| interpreted differently for memory and memory |  |
| mapped $\mathrm{I} / 0$. |  |

Nib 0 Memory Size MM I/O space

| F | 1 K nibble |  | 16 nibbles |  |
| :---: | :---: | :---: | :---: | :---: |
| E | 2 |  | 32 |  |
| D | 4 |  | 64 |  |
| C | 8 |  | 128 |  |
| B | 16 |  | 256 |  |
| A | 32 |  | 512 |  |
| 9 | 64 | (max RAM) | 1 K |  |
| 8 | 128 |  | 2 K |  |
| 7 | 256 | $(\max )$ | 4 K |  |
| 6 | - |  | 8 K |  |
| 7 | - |  | 16 K | (max) |

NIBBLE 1 : Reserved for future use.
NIBBLE 2 : Device type--
0 : RAM
1 : ROM
2-E : Assorted memory types
F : Memory-mapped I/O
NIBBLE 3 : Device class--
Memory : unassigned
Memory-mapped I/O - 0 : HP-IL mailbox 1-F : unassigned

NIBBLE 4 : bits 0-1 : unassigned
bit 2 : Last chip in sequence. Always assumed high for MM I/O.
bit 3 : Last chip in module.

Since the BUS [0:3] is precharged low before each strobe, the CPU will read an TD of all zeros if ajJ devices are configured (or are unconfigured but have DAISYIN low). For more information on how the operating system handles configuration see the HP-71 Software IDS Volume 1.

A soft configured device is assigned its address configuration by the CONFIGURE command. If an unconfigured device has its

DAISYIN line high, it loads the configuration address that is issued on the 5 data strobes immediately following the CONFIGURE command (low-order nibble first) into its configuration register. A device may actually latch only the number of high-order bits it requires as determined by its address space size.

After being configured a device no longer responds to either an ID or CONFIGURE command. A configured device drives DAISYOUT to the same logic level as DAISYIN. The DAISYOUT of one device may be tied to the DAISYIN of second device. In this way many devices may be daisy-chained together in a way that they can be configured one at a time to different addresses. After being configured a device waits until the next command strobe to set its configuration flag in order to delay DAISYOUT so that the next device on the daisy-chain will not be configured simultaneously.

A device may be unconfigured by either a RESET or UNCONFIGURE bus command. The bus RESET command unconfigures all soft configured devices in the system. A device responds to an UNCONFIGURE command by clearing its configuration flag if the DP is within its address configuration.

### 2.2.2 Hard Configuration

A hard configured device powers-up configured to a specific address. It will not respond to an ID, CONFIGURE, or UNCONFIGURE command and a bus RESET will not affect its configuration. If the device has a DAISYOUT, it is always driven to the same logic level as its DAISYIN.

### 2.3 Data Transfer

All information that is transferred from the CPU to other devices in the system (commands, addresses, and data) is latched from the BUS[0:3] on the rising edge of "STR. Data that is transferred from system devices to the CPU is latched off the BUS [0:3] after the falling edge of "STR (timed internally on the CPU).

The CPU loads all devices' local address registers by issuing a LOAD PC or LOAD DP bus command followed by 5 data strobes of address, least significant nibble first. After the last nibble of address has been loaded all devices auto-switch to a PC READ or DP READ bus command. The CPU may then read the contents of that address location by issuing one dummy strobe followed by 1 to 16 data strobes during which the CPU latches the BUS[0:3] data. The

CPU may read without first loading the local address registers by issuing a PC READ or DP READ, followed by a dummy strobe, followed by 1 to 16 data strobes. The CPU precharges the BUS[0:3] low each cycle before "STR goes low. Therefore if no device responds the CPU reads zeros.

The CPU writes the contents of a specific addressed location similarly. It $i s$ not required to load the local address registers immediately before issuing a PC WRITE or DP WRITE command. The write command is followed by 1 to 16 data strobes during which the addressed device latches the BUS[0:3] data.

All devices increment their local address registers once each data strobe during read and write operations. It is possible for a read or write operation to begin in one device and cross the address boundary into another device. Future controllers on the HP-71 bus may read and write more than 16 nibbles at a time.

Two other types of data transfers are ID, which is simply a $5-n i b b l e$ read with no address load or dummy strobe, and CONFIGURE, which is a 5-nibble write with no address load. Both these data transfers require that a device be unconfigured and that the DAISYIN line be high. POLL is a unique read and is discussed in section 2.5.

### 2.4 Power down; Wake up

The HP-71 system can be shutdown under software control. The CPU executes a SHUTDN i:sstruction by issuing a SHUTDOWN bus command and on the ensuing cycle stopping the system clock ("STR) and its own oscillator. While in shutdown mode all data stored in RAM and CPU resident memory is preserved. The CPU is brought out of shutdown mode by either pulling an Input Register line high, or by driving *CD low.
*CD is driven low to wake up the CPU primarily by a device in the system that needs service while the system is in shutdown mode. If a device wakes up the CPU and the CPU shuts down without satisfying its service request the device will not wake up the CPU again until its service request has been satisfied and it needs service again. This avoids a situation where the operating system does not know how to handle a device's service request and cannot shutdown. For more information on power down and wake up see section 3.5 .

### 2.5 Service Poll

If a device needs service while the CPU is operating it must wait until the CPU executes a service request instruction (SREQ), or, if it has the capability, interrupt the CPU using IR14 (available at all ports). The SREQ instruction causes the CPU to issue a POLL bus command followed by one data strobe during which the CPU latches the BUS [0:3] data in the manner of a usual read. A device may respond to the service POLL by pulling one of the BUS[0:3] lines high. Since the CPU precharges the BUS[0:3] low every cycle before *STR goes low the data read by the CPU is a binary $O R$ of all devices' responses.

The following HP-71 chips can wake-up the CPU and can respond to a service POLL on the BUS [0:3] line shown:

| Device | Bus line | Reason for Service Request |
| :---: | :---: | :---: |
| Display Driver | BUS[0] | Timer underflow. |
| HP-IL chip | BUS [1] | Data Avail; Interrupt; |
|  |  | Pwr-on Reset; Loop Service Request. |
| Card Reader | BUS[2] | FIFO servicing; Error condition. |



### 3.1 CPU Overview

The 1LF2 CPU is a proprietary CPU optimized for high-accuracy BCD math and low power consumption. The CPU's principle function is to fetch and execute micro-instructions in proper sequence. Memory is accessed in 4 -bit quantities called "nibbles". Addresses are 20 bits, providing a physical address space of 512 K bytes. The CPU operates internally on four 4 -bit busses with one or two busses acting as the source and one or two busses acting as the destination. Data for either source bus can originate from the system Bus, CPU resident memory, or other CPU register. Operations performed by the ALU include "add", "subtract", "and", "or", "1's complement", "2's complement", and in combination with the shifter "bit-shift", "digit-shift", and "rotate", on up to 64-bit operands.

For information regarding the full CPU instruction set see Chapter 4, entitled "HP-71 ASSEMBLER INSTRUCTION SET".

### 3.2 Pin Designations

The 1LF2 CPU's external pins are as follows:

| PIN | FUNCTION |
| :--- | :--- |
| VDD | Power supply. |
| GND | System ground potential. |
| BUS [0:3] | System bus. |
| "STR | *STROBE |
| "CD | *COMMAND-DATA |


| $\operatorname{IR}[0: 15]$ | INPUT REGISTER - For keyboard input and general input. |
| :---: | :---: |
| OR [0:11] | OUTPUT REGISTER - For keyboard output and general output. |
| HALT | HALT - When asserted high, the CPU completes the current instruction, tristates the Bus, *STR, and "CD lines and waits in a loop until HALT goes low again. The Bus is held passively low, "STR and *CD passively high. |
| NIE | NEXT INSTRUCTION FETCH - Goes high to indicate occurrence of the next instruction following the current "STR. |
| *INT | INTERRUPT - When pulled low the CPU completes the current instruction, pushes the PC onto the top of the subroutine return stack, and initiates the interrupt routine. |
| OSC1, OSC2 | Oscillator input/output used for LC connection. |
| ECE | EXTERNAL CLOCK ENABLE - tied low on PC board. |
| CIO | CLOCK I/O - The primary oscillator frequency is driven out on this pin. |
| DRI | DRIVE - High when the CPU is driving the Bus, low if the CPU is sensing or tristated. |
| ENP | ENABLE POWER SUPPLY - Tied low on PC board. |
| VCO | VOLTAGE CONTROL OUT - Not used. |
| VCI | VOLTAGE CONTROL IN - Not used. |
| ote that "INT is available at al.l ports except the Card Reader ; IR14 is available at all ports; and HALT is available at all s except PORT1. The Output Register is not only used to form keyboard matrix, but also drives the DAISYIN (DIN) signal to port and the piezo-electric beeper. |  |
|  |  |
|  |  |

### 3.3 Registers

There are two types of registers on the CPU; those used for data transfers and arithmetic operations, and those used for program and system control.

Arithmetic Registers:

| A | 64-bits | Working register - I/O register |
| :---: | :---: | :---: |
| B | 64-bits | Working register |
| C | 64-bits | Working register - I/O register |
| D | 64-bits | Working register |
| R0 | 64-bits | Scratch register |
| R1 | 64-bits | Scratch register |
| R2 | $64-\mathrm{bits}$ | Scratch register |
| R3 | 64-bits | Scratch register |
| R4 | 64-bits | Scratch register |
| CARRY | 1-bit | Flag set by arithmetic operations and tests |

## Control Registers:

| POINTER | 4-bits | Pointer register |
| :---: | :---: | :---: |
| DPOINTO | 20-bits | Address pointer register |
| DPOINT1 | 20-bits | Address pointer register |
| PC | 20-bits | Program Counter |
| RETURN STACK | 20-bits | 8-level subroutine stack |
| status | 16-bits | Program status flags |
| HW. STATUS | 4-bits | CPU/system status flags |
| OUTPUT | 12-bits | Keyscan/write only Output |
| INPUT | 16-bits | Keyscan/read only Input R |

### 3.3.1 Working and Scratch Registers

All arithmetic operations are performed using the 4 working registers: $A, B, C$, and $D$. Data transfers are performed principally with the A and C registers.

The scratch registers are used to temporarily hold the contents of the working registers. The lower 20 bits of scratch register R 4 are reserved by the operating system for interrupt processing, and therefore are not normally available for data storage.

### 3.3.1.1 Field Selection

Subfields of the working registers may be manipulated using field selection. The possible field selections range from the entire register to any single nibble of the register. Certain subfields are designed for use in BCD calculations and others are designed for use in general data manipulation and data access.

FIELD SELECTION FIELDS

| $P$ | Digit pointed to by P register |
| ---: | :--- |
| WP | Digit 0 through digit pointed at by $P$ |
| XS | Digit 2 $\quad$ - Exponent sign |
| X | Digits $0-2$ - Exponent and exponent sign |
| S | Digit $15 \quad$ - Mantissa sign |
| M | Digits $3-14$ - Mantissa |
| B | Digits $0-1$ - Exponent or byte field |
| W | Digits $0-15$ - Whole word |
| A | Digits $0-4$ - Address field |

Nibbles of Register

15:14:13:12:11:10: 9: 8: 7: 6: 5: 4: 3: 2: 1: 0
$|S| \quad|X S|<-B->\mid$



### 3.3.2 Carry Bit

The Carry bit is adjusted when a arithmetic operation or test is performed. During a calculation, such as incrementing or decrementing a register, it is set if the calculation overflows or borrows and cleared if it does not. During a arithmetic test, such as comparing two registers for equality, it is set if the test is true and cleared if it is not.

### 3.3.3 Pointer Registers

The Data Pointer registers, D0 and D1, provide the source addresses for all external data transfers.

The $P$ Pointer register is used in Field Selection operations with the working registers.

### 3.3.4 Program Counter and Return Stack

The program counter points to the next instruction to be executed by the CPU. It can be accessed only using jump, gosub, and return instructions.

The current value of the program counter is automatically pushed onto the 8 -level subroutine return stack when a gosub instruction is executed or an interrupt occurs. A 20-bit value is automatically popped off the return stack into the program counter when a return instruction is executed. The return stack can also be manipulated through use of the push (RSTK=C) and pop (C=RSTK) instructions.

### 3.3.5 Status Bits

Additional program control is provided by the 16-bit Program Status register and the 4 -bit Hardware Status register. Each Status bit can be individually set, reset, and tested.

The operating system uses the upper 4 Program Status bits to indicate the state of the operating system. The remaining 12 Program Status bits are generally available to applications software, and may be manipulated collectively as the ST register.

The four Hardware Status bits are set (but not cleared) by hardware-related events, and must therefore be cleared beforehand in order to detect a particular occurrence. They are individually accessible by name. The Module Pulled bit (MP) is set when a module is pulled from or added to the machine. The Sticky Bit (SB) is set when a non-zero bit or digit shifts off the right end of a working register as the result of a shift right instruction, or the least significant nibble of a working register is non-zero prior to a shift right circular instruction. The Service Request (SR) bit is set as a result of a response to the SREQ? instruction (see section 2.5). The external Module Missing bit is set by execution
of a "OO" opcode (RTNSXM instruction). Since the BUS[0:3] is precharged low, the CPU will receive a RTNSXM instruction if no device responds to a PC READ bus command.

PROGRAM STATUS: 16 bits
Bits Usage
15 thru 12 Indicate state of operating system
11 thru 0 Available to programs, may be manipulated as the ST register

HARDWARE STATUS: 4 bits

| Bit | Symbol | Name |
| :---: | :---: | :--- |
| 3 | MP | Module Pulled |
| 2 | SR | Service Request |
| 1 | SB | Sticky Bit |
| 0 | XM | External Module Missing |

### 3.3.6 Input/Output Registers

The Input and Output Registers provide the CPU with general purpose I/O. This consists of a 16-bit Input Register with interrupt capability (see section 3.5) and a 12-bit Output Register.

Data read from the Input Register corresponds to the logic levels sensed on the 16 input lines, IR[0:15]. The Input Register lines are passively held low so that if an input line is not driven a zero level will be sensed. Data written to the Output Register is driven onto the 12 output lines OR[0:11].

The Input and Output Registers are used to form the keyboard matrix. The Output Register lines OR[0:3] drive the 4 key-rows. The Input Register lines IR[0:13] sense the 14 key-columns. The most significant bit of the Input Register (IR15) is dedicated to the On-key and has additional interrupt capability (see section 4.4.2). When a key is pushed its key-row is shorted to its key-column, and the logic level driven on the Output Register line can be read from the Input Register bit. If the Output Register line is high and interrupts are enabled, pushing the key will result in an interrupt.

When one of the 8 lower order bits of the Output Register is low, the Output Register line will be actively driven low briefly
each instruction cycle, then passively held low. This limits the current and provides a deterministic state when Output Register lines are driving different logic levels and are shorted. Output Register lines are shorted when 2 or more keys in the same column are depressed at the same time. This will result in a logic one level on the corresponding bit of the Input Register if one of the Output Register lines is high. The remaining 4 Output Register lines, $O R[8: 11]$ are actively driven at all times.

The 5 Output Register lines $O R[0: 4]$ also drive the DAISYIN lines of each I/O port. During configuration the operating system individually selects each I/O port by driving its DAISYIN line high (see section 8.3).

The 2 Output Register lines $O R[10: 11]$ drive the piezo-electric beeper. Two beeper loudness settings are obtained by using one of two Output Register lines.

### 3.3.7 Loading Data from Memory

When data is read from an external device into a register, the CPU places the lowest addressed nibble in the least significant nibble of the register. For example, if the data shown below in memory is read into the $C$ register using the C=DAT1 4 instruction, the data in the register will be arranged as shown above.

## Memory

| Location | Value |  | C | Register |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1000 | 0 | I |  | 3 | 2 | 1 | 01 |
| 1, 001 | 1 | +- |  |  |  |  | + |
| 1002 | 2 | 15 |  | 3 | 2 | 1 | 0 |
| 1003 | 3 |  |  |  |  |  |  |

This principle also applies to loading constants into a CPU register such as C, DO, or D1, since the CPU must read the constant from the instruction opcode in memory. For example, the instruction LCHEX 3210 produces the opcode 330123 and the $C$ register is loaded as shown above.

### 3.3.8 Storing Data in Memory

When data is written from a register to an external device, the CPU places the least significant nibble of the register in the first addressed nibble of the memory location. For example, if the data shown above in the $C$ register is written to memory using the DAT1=C 4 instruction, the data will be written to memory as shown.

## 3. 4 HEX/DEC Modes

All arithmetic operations, except for those listed below, are performed according to the HEX or DEC mode setting. The mode is set using the SETHEX or SETDEC instruction. The following operations are performed in HEX regardless of the mode setting.

```
C+P+1
DO=DO+n DO=DO-n
D1=D1+ n D
P=P+1 P=P-1
```


### 3.5 Interrupt System

I'he 1LF2 CPU can be interrupted in the following manners:

1. Pulling the *IN' line low (module interrupt).
2. Shorting the least significant bit of the Output Register to the most significant bit of the Input Register (hardware "ON"-key).
3. An Input Register line (bits 0-15) being pulled high.

If interrupted, the CPU completes the current instruction, pushes the PC onto the return stack and sets the PC to the interrupt routine starting address of 0000 F (where $F$ is in the least significant nibble of the PC). Intexrupts are disabled upon entry into the interrupt routine and are not enabled until a Return From Interrupt (RTI) instruction is executed. Each interrupt type operates in a somewhat different manner and is described separately.

### 3.5.1 \#NT Interrupt

Pulling the *INT line low causes a non-maskable interrupt. When *INT is asserted the CPU enters the interrupt routine and disables all interrupts. When a RTI instruction is issued all interrupts may be processed immediately, including the *INT line even if it remains low from the previous interrupt.
3.5.2 Hardware "On-key" Interrupt

During the last cycle of every instruction the least significant bit of the Output Register (ORO) is driven to a one. This does not modify the contents of the Output Register, and the previous value is driven after this cycle. If the MSB of the Input Register (IR15) is driven high (On-key down) during this period, an interrupt is executed. This interrupt cannot be disabled outside of the interrupt routine.

Once a RTI instruction has been issued the CPU cannot be interrupted again by this Input Register line or any others until all Input Register lines are low. The CPU can be interrupted by the "INT line regardless of the state of the Input Register lines as long as a RTI instruction has been issued.

### 3.5.3 Input Register Interrupt

At any time except the time period allocated to the hardware "On-key", the CPU may be interrupted by any Input Register line pulled high. This interrupt may be disabled outside the interrupt routine by using the Interrupt off (INTOFF) instruction, and re-enabled using the Interrupt on (INTON) instruction.

The CPU cannot be interrupted again by any Input Register line until a Rl'I instruction is issued and all Input Register lines are again low.

### 3.6 CPU Power-down; Wake-up

The CPU supports a low-power standby mode. To place the system in standby, a SHUTDN instruction is issued. The CPU then issues the SHUTDOWN bus command and on the ensuing cycle, stops "STR and its own clock. All CPU resident memory is preserved during standby. In standby mode the CPU holds the Bus lines in the conditions as follows:

| BUS[0:3] | Low - through a high impedance |
| :--- | :--- |
| "STR | High - through a high impedance |
| "CD | High - through a high impedance |

The CPU is brought out of standby mode when either an Input Register line goes high or the ${ }^{*} C D$ line is driven low. On wake-up the CPU starts it's clock and immediately drives *CD low with a NOP bus command. If the CPU detects that a severe low voltage condition has occurred while it was in standby, the PC is set to zero and hex mode arithmetic is asserted. A LOAD PC bus command is issued after the NOP and the current CPU PC is loaded into all devices' local PCs. At this point the standard instruction fetch sequence in initiated. If the CPU was awakened by the Input Register and interrupts are enabled a normal interrupt will occur prior to the first instruction fetch. If a SHUTDN instruction is executed with the least significant bit of the Output Register (ORO) set at a "O" the CPU will wake-up immediately, set its PC to zero, assert hex mode arithmetic, and send out the PC. This avoids a situation where the CPU is shutdown and the keyboard cannot wake it up.


This chapter describes the HP-71 assembler instruction set. The instruction mnemonics shown are those provided by the assembler used by the KP-71 software development team (which is available by special arrangement with HP). Almost all the mnemonics shown are a.lso supported by the HP-71 FORTH/Assembler ROM.

### 4.1 Instruction Syntax

### 4.1.1 Labels and Symbols

A label is a symbolic name for a numeric value. A label acquires its value by appearing in the label field of certain statements. The word "symbol" is a general term for a label, and the two are used interchangeably.

Labels are one to six alphanumberic characters with the following restrictions: the characters comma (,), space () and right parenthesis are prohibited and the first character cannot be equal sign (=), sharp (\#), single quote ('), left parenthesis, or the digits 0 through 9.

A label may be immediately preceded by an equal sign which declares the label to be an external symbol. An external symbol defined in one module may be referenced as an external symbol by another module. Such references are resolved when the modules are linked together. Certain HP-71 assemblers, such as the FORTH/ASSEMBLER ROM, have no associated linker and therefore do not support external symbols. In this case, any leading equal sign is ignored.

When a label is used as part of an expression, parentheses are required to delineate it. That is, $A D 1-10$ is a label but (AD1)-10 is a computed expression.

### 4.1.2 Comments

A comment line begins with an asterisk (*) in column one, and may occur anywhere. An in-line comment may begin with any non-blank character and must follow the modifier field of an instruction (or the opcode if no modifier is required).

### 4.1.3 Expressions

Wherever an expression may appear in the modifier field of an instruction, it is represented by the symbol "expr" in the instruction descriptions below. Expressions consist of:

## EXPRESSION COMPONENTS

Item
decimal constants
hexadecimal constants ascii constants

## operators

* 

label Symbol defined in the label field of an instruction

Parenthesized expression

Two classes of instructions require a modifier field which contains a constant of a specific type that does not conform to the above rules. These are:
a) String constant which can exceed 3 characters

| LCASC | 'ascii' | or |
| :--- | :---: | :---: |
| LCASC | \ascii\ |  |
| NIBASC | 'ascii' | or |
| NIBASC | \ascii\ |  |

b) Unconditional Hex constant

LCHEX 4 FFFFF
NIBHEX $4 F F F F F$

### 4.1.4 Sample Line Image

The format below is the recommended column alignment; however, the assembler is "free format" and only a space is required to delimit the different fields. A label, if present, must start by column 2.

| 1 | 8 | 15 | 31 | 80 |
| ---: | ---: | ---: | ---: | ---: |
| v | v | v | v | v |

label opcode modifier comments

### 4.2 Explanation of Symbols

In the following descriptions of the HP-71 assembler mnemonics, these symbols have the following meanings unless specified otherwise. In particular, note the symbols used to indicate the various values encoded within the assembled opcodes.
a The hex digit used to encode the field selection in the assembled opcode of an instruction. See the Field Select Table in the next section for details.
b The hex digit used to encode the field selection in the assembled opcode of an instruction. See the Field Select Table in the next section for details.
d The number of digits represented by a field selection field. Used in calculating the execution cycle time

|  | of some instructions. See the Field Select Table in the next section for details. When used in an extended field selection fsd, represents an expression which indicates the number of nibbles of the register that will be affected by the instruction, proceeding from the low-order nibble to higher-order nibbles. |
| :---: | :---: |
| expr | An expression that evaluates to an absolute or relocatable value, usually less than or equal to 5 nibbles in length. |
| f's | Field selection symbol. See the Field Select Table in the next section for details. |
| fsd | Extended field selection symbol. Represents either a normal field selection symbol fs, or an expression that gives the number of digits $d$ of the register that will be affected by the instruction, proceeding from the low-order nibble to higher-order nibbles. |
| hh | Two-digit hex constant, such as 08 or F2. Within an opcode represents the hex digits used to store the value of the expression in the opcode in reverse order (see "Loading Data From Memory"). |
| hhhh | Four-digit hex constant, such as 38 FE . Within an opcode, represents the hex digits used to store the value of the expression in the opcode in reverse order (see "Loading Data From Memory"). |
| hhhhh | Five-digit hex constant, such as 308 FE . Within an opcode, represents the hex digits used to store the value of the expression in the opcode in reverse order (see "Loading Data From Memory"). |
| label | A symbol defined in the label field of an instruction. |
| m | A one-digit decimal integer constant. |
| n | Represents an expression that evaluates to a 1-nibble value, unless specified otherwise. Within an opcode, represents the hex digit used to store the assembled value of the expression in the opcode. |
| n | Represents an expression that evaluates to a 2-nibble value, unless specified otherwise. Within an opcode, represents the hex digits used to store the assembled value of the expression in the opcode. |
| nnnn | Represents an expression that evaluates to a 4-nibble value, unless specified otherwise. Within an opcode, |

represents the hex digits used to store the assembled value of the expression in the opcode.
nnnnn Represents an expression that evaluates to a 5-nibble value, unless specified otherwise. Within an opcode, represents the hex digits used to store the assembled value of the expression in the opcode.

### 4.2.1 Field Select Table

The following symbols are $u$. d in the instruction descriptions to denote the various possible field selections.

There are two ways in which field selection is encoded in the opcode of an instruction. These two patterns are shown in the table below, and are designated by the letter ' $a$ ' or ' $b$ ' in the opcode value given in the mnemonic descriptions below.

FIELD SELECT TABLE

| Field | Name and Description | Op Repres <br> (a) | de <br> tation <br> (b) | Number of Digits <br> (d) |
| :---: | :---: | :---: | :---: | :---: |
| P | Pointer Field. Digit specified by $P$ pointer register. | 0 | 8 | 1 |
| WP | Word-through-Pointer Field. Digits 0 through (P). | 1 | 9 | (P+1) |
| XS | Exponent Sign Field. Digit 2. | 2 | A | 1 |
| X | Exponent Field. Digits 0-2. | 3 | B | 3 |
| S | Sign Field. Digit 15. | 4 | C | 1 |
| M | Mantissa Field. Digits 3-14. | 5 | D | 12 |
| B | Byte Field. Digits 0-1. | 6 | E | 2 |
| W | Word Field. All digits. | 7 | F | 16 |

The above field selects generally share the same opcode, with one

```
nibble of the opcode containing the 'a' or ' b' value as specified
in the table. The 'A' field select, however, generally is
specified by a different opcode altogether. This has the effect of
shortening and speeding up execution of 'A' field select
manipulations, optimizing the computer for address and 5-nibble
calculations.
```


### 4.3 Instruction Set Overview

The following pages briefly summarize the HP-71 instruction set. For further details please refer to the Mnemonic Dictionary which follows this summary.

### 4.3.1 GOTO Instructions

```
    l = Statement Label
```

    ---
    GOTO label Short unconditional branch
    GOC label Short branch if Carry set
    GONC label Short branch if no Carry set
    GOLONG label Long GOTO
    GOVJ,NG label Very long GOTO
    GOYES label Short branch if test true (must
                        follow a Test Instruction)
    
### 4.3.2 GOSUB Instructions

| GOSUB label | Short transfer to subroutine |
| :--- | :--- |
| GOSUBL label. | Long GOSUB |
| GOSBVL label | Very long GOSUB |

### 4.3.3 Subroutine Returns

| RTN | Unconditional return |
| :--- | :--- |
| RTNSC | Return and set Carry |
| RTNCC | Return and clear Carry |
| RTNSXM | Return and set XM bit (Module Missing) |
| RTI | Return and enable interrupts |
| RTNC | Return if Carry set |


| RTNNC | Return if no Carry set |
| :--- | :--- |
| RTNYES | Return if test true (must follow a |
|  | Test Instruction) |

### 4.3.4 Test Instructions

All test instructions must be followed with a GOYES or a RTNYES instruction. Although they appear to be two statements, in fact they combine to be one. Each test adjusts the Carry bit when performed: Carry is set if the test is true, and cleared if false.
4.3.4.1 Register Tests
$r, s=A, B, C$ or $(r, s)=(C, D),(D, C)$ fs = Field Select
----
?r=s fs Equal
?r\#s fs Not equal
?r=0 fs Equal to zero
?r\#0 fs Not equal to zero
?r>s fs Greater than
?r<s fs Less than
?r>=s fs Greater than or equal
? $r<=s$ fs Less than or equal
4.3.4.2 P Pointer Tests
----
$0<n<=15$
--
$? \mathrm{P}=\mathrm{n} \quad$ Is P Pointer equal to n ?
?P\# $n \quad P$ Pointer not equal to $n$ ?

### 4.3.4.3 Hardware Status Bit Tests

| $? \mathrm{XM}=0$ | Module Missing bit equal to zero? |
| :--- | :--- |
| $? \mathrm{SB}=0$ | Sticky Bit equal to zero? |
| $? \mathrm{SR}=0$ | Service Request bit equal to zero? |
| $? \mathrm{MP}=0$ | Module Pulled bit equal to zero? |

4.3.4.4 Program Status Bit Tests
--. -
$0<=\mathrm{n}<=15$
----

| $? S T=1$ | $n$ | Status $n$ equal to $1 ?$ |
| :--- | :--- | :--- |
| ?ST=0 | $n$ | Status $n$ equal to $0 ?$ |
| ?ST\#1 | $n$ | Status not equal to $1 ?$ |
| ?ST\#0 | $n$ | Status not equal to $0 ?$ |

4.3.5 P Pointer Instructions

---
$0<=n<=15$
----
$\mathrm{P}=\mathrm{n}$ Set P Pointer to n
$\mathrm{P}=\mathrm{P}+1 \quad$ Increment P Pointer, adjust Carry
$\mathrm{P}=\mathrm{P}-1 \quad$ Decrement P Pointer, adjust Carry
C+P+1 Add P Pointer plus one to A-field of C
CPEX $n \quad$ Exchange $P$ Pointer with nibble $n$ of $C$
$\mathrm{P}=\mathrm{C} \quad \mathrm{n} \quad$ Copy nibble n of C into P Pointer
$\mathrm{C}=\mathrm{P} \quad \mathrm{n} \quad$ Copy P Pointer into nibble n of C
4.3.6 Status Instructions
4.3.6.1 Program Status

```
    0<= n <= 15
```

    ----
    \(S_{r}=1 \quad n \quad\) Set Status \(n\) to 1
    \(\mathrm{ST}=0 \mathrm{n} \quad\) Set Status n to 0
    CSTEX Exchange \(X\) field of \(C\) with Status 0-11
    \(C=S T \quad\) Copy Status 0-11 into \(X\) field of \(C\)
    ST=C Copy X field of \(C\) into Status 0-11
    CLRST Clear Status 0-11
    4.3.6.2 Hardware Status

| SB=0 | Clear Sticky Bit |
| :--- | :--- |
| SR=0 | Clear Service Request bit (see SREQ?) |
| $M P=0$ | Clear Module-Pulled bit |
| $X M=0$ | Clear External Module Missing bit |
| CLRHST | Clear all 4 Hardware Status bits |

4.3.7 System Control
SETHEX Set arithmetic mode to hexadecimal
SETDEC Set arithmetic mode to decimal
SREQ? Sets Service Request bit if service hashas been requested. $C(0)$ shows whatbit(s) are pulled high (if any)
C=RSTK Pop return stack into A-field of CRSTK=CCONFIGPush A-field of $C$ onto return stackConfigure
UNCNFG UnconfigureRESETSend Reset command to system bus
BUSCCSend Bus command $C$ onto system bus
SHUTDNStop CPU here (sleeps until wake-up)$C=I D$INTOFF
INTON
Request chip ID into $A$-field of $C$
Disable interrupts (doesn't affect ON-keyor module-pulled interrupts)
Enable interrupts
4.3.8 Keyscan Instructions

| OUT $=C$ | Copy X field of $C$ to OUTput register |
| :--- | :--- |
| OUT $=C S$ | Copy nibble of $C$ to OUTput register |
| $A=I N$ | Copy INput register to lower 4 nibbles of $A$ |
| $C=I N$ | Copy INput register to lower 4 nibbles of $C$ |

4.3.9 Register Swaps$s=R 0, R 1, R 2, R 3, R 4$
---
AsEX Exchange register $A$ with s CsEX Exchange register $C$ with s

    \(A=s \quad\) Copy \(s\) to register \(A\)
    
    \(\mathrm{C}=\mathrm{s} \quad\) Copy s to register C
    
    \(s=A \quad\) Copy register A to s
    
    \(s=C \quad\) Copy register \(C\) to \(s\)
    4.3.10 Data Pointer Manipulation

$$
d=D 0, D 1
$$

```
    1 <= n <= 16
    expr <= 5 nibbles
----
AdEX Exchange Data ptr d with A-field of A
CdEX Exchange Data ptr d with A-field of C
AdXS Exchange lower 4 nibs of Data ptr d with
    lower 4 nibs of A
CdXS Exchange lower 4 nibs of Data ptr d with
    lower 4 nibs of C
d=A Copy A-field of A to Data pointer d
d=C Copy A-field of C to Data pointer d
d=AS
d=CS
Copy lower 4 nibs of C to lower }4\mathrm{ nibs
    of Data pointer d
d=d+ n Increment Data pointer d by n
d=d- n Decrement Data pointer d by n
d=HEX hh Load hh into lower 2 nibs of Data ptr d
d=HEX hhhh Load hhhh into lower 4 nibs of Data ptr d
d=HEX hhhhh Load hhhhh into Data ptr d
d=(2) nn Load nn into lower 2 nibs of Data ptr d
    (any overflow is ignored)
d=(4) nnnn Load nnnn into lower 4 nibs of Data ptr d
    (any overflow is ignored)
d=(5) nnnnn Load nnnnn into Data ptr d (any overflow
    is ignored)
```


### 4.3.11 Data Transfer

    ---
    fsd \(=\) Field select fs, or \(d\) (\# of digits). If \(d\), then
                the copy starts at nibble 0 of the working register.
    \(1<=d<=16\)
    ----

| $A=D A T O$ fsd | Copy data from memory addressed by DO into A, field selected |
| :---: | :---: |
| C=DATO fsd | Copy data from memory addressed by DO into C, field selected |
| $\mathrm{A}=\mathrm{DAT} 1$ fsd | Copy data from memory addressed by D1 into A, field selected |
| $\mathrm{C}=$ DAT1 fsd | Copy data from memory addressed by D1 into C, field selected |
| DATO $=$ A fsd | Copy data from $A$ into memory addressed by DO, field selected |
| DATO $=$ C fsd | Copy data from $C$ into memory addressed by DO, field selected |
| DAT1=A fsd | Copy data from $A$ into memory addressed by |

```
                                    D1, field selected
DAT1=C fsd Copy data from C into memory addressed by
                                    D1, field selected
```


### 4.3.12 Load Constants

LCHEX hhhhhhhh Load hex constant into C (1 to 16 digits)
$L C(m)$ expr Load the m-nibble constant into $C$
LCASC 'ascii' Load up to 8 ASCII characters into C
LCASC \ascii\ Load up to 8 ASCII characters into C

### 4.3.13 Shift Instructions

Note that right shifts (circular or non-circular) will set the Sticky Bit ( $\mathrm{SB}=1$ ) if a nonzero bit is shifted off to the right. Otherwise, $S B$ is unchanged.

```
----
    r = A,B,C,D
    fs = Field Select
    ----
```

| rSL | fs | Shift register $r$ fs field Left 1 nibble |  |
| :--- | :--- | :--- | :--- |
| rSR | fs | Shift register r field Right 1 nibble |  |
| rSLC |  | Shift register r Left Circular 1 nibble |  |
| rSRC |  | Shift register r Right Circular 1 nibble |  |
| rSRB |  | Shift register | Right |

### 4.3.14 Logical Operations

Logical operations are bit-wise.

```
----
    r,s=A,B,C or (r,s)=(C,D),(D,C)
    fs = Field Select
    ----
    r=r&s fs r AND s into r, fiedd selected
    r=r!s fs r OR s into r, fjeld selected
```


### 4.3.15 Arithmetics

The two groups of arithmetics differ in the range of registers available. In the first group (General usage) almost all combinations of the four working registers are possible; however,

```
in the second group (Restricted usage) only a few select
combinations are possible.
4.3.15.1 General Usage
    ----
        r,s=A,B,C or (r,s)=(C,D),(D,C)
        fs = Field Select
    ----
    r=0 fs Set r to zero
    r=r+r fs Double r, adjust Carry
    r=r+1 fs Increment r by 1, adjust Carry
    r=r-1 fs Decrement r by 1, adjust Carry
    r=-r fs 10'S complement or 2'S complement, Carry
    set if r#O, else clear
r=-r-1 fs 9'S complement or 1'S complement
    Carry always cleared
r=r+s fs Sum r and s into r, adjust Carry
s=r+s fs Sum r and s into s, adjust Carry
r=s fs Copy s into r
s=r fs Copy r into s
rsEX fs Exchange r and s
4.3.15.2 Restricted Usage
----
    (r,s) = (A,B),(B,C),(C,A),(D,C)
----
r=r-s fs Difference of r and s into r, adjust Carry
r=s-r fs Difference of s and r into r, adjust Carry
s=s-r fs Difference of }s\mathrm{ and }r\mathrm{ into }s\mathrm{ , adjust Carry
```

4.3.16 No-Op Instructions

Execution of a No-Op affects no CPU registers except for the PC. NOP3 Three nibble No-Op NOP4 Four nibble No-Op NOP5 Five nibble No-Op

### 4.3.17 Pseudo-Ops

4.3.17.1 Data Storage Allocation

    ---
    \(1<=n<=8\)
    ----
        BSS nnnnn Allocate nnnnn number of zero nibs
        \(\operatorname{CON}(m)\) expr Generate \(m-n i b b l e ~ c o n s t a n t ~(d i g i t s ~ a r e ~\)
        reversed in the opcode)
    REL ( \(m\) ) expr Generate m-nibble relative constant (digits
        reversed in the opcode)
    NIBASC 'ascii' Generate ascii characters, byte reversed
    NIBASC \ascii\ Generate ascii characters, byte reversed
    NIBHEX hhhh Generate hexadecimal digits hhhh (digits
    are not reversed in the opcode)
    4.3.17.2 Conditional Assembly
name IF expr Start conditional assembly until ELSE or
ENDIF if flag expr was set on invocation
of assembler (optional use of name allows
nesting of IF's)
name ELSE Conditional assembly if IF test was false
name ENDIF Ends conditional assembly started by IF

### 4.3.17.3 Listing Formatting

EJECT Force new page in the assembly listing STITLE text Force new page, set subtitle value to text TITLE text Set sitle value to text

### 4.3.17.4 Symbol Definition

label EQU nnnnn Defines label to have the value expr
4.3.17.5 Assembly Mode
ABS nnnnn Specify absolute assembly at adress given END Marks end of the assembly source
4.4 Mnemonic Dictionary

This section contains a description of each HP-71 assembler instruction or pseudo-op. The description shows the binary opcode generated by the mnemonic, if any, as well as the execution cycle time required if the mnemonic is an executable instruction.

The symbols used in these descriptions are explained in the "Explanation of Symbols" section earlier in this chapter.

##  <br> * MNEMONICS <br> 

?A\#O fs - Test for $A$ not equal to 0
---------
fs $=\mathrm{A} \quad$ opcode: 8ACyy
cycles: $13+d$ (GO/RTNYES)
$6+d$ (NO)
$f s=(P, W P, X S, X, S, M, B, W) \quad o p c o d e: ~ 9 a С y y ~$
cycles: $13+d$ (GO/RTNYES) $6+d$ (NO)

Test whether the fis field of $A$ is not equal to 0 . Must ise followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?A#B fs - Test for A not equal to B
-
fs = A opcode: 8Alyy
    cycles: 13 + d (GO/RTNYES)
        6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9a4yy
    cycles: 13 + d (GO/RTNYES)
```

$$
6+d(\mathrm{NO})
$$

Test whether the fs field of $A$ is not equal to the fs field of $B$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?A#C fs - Test for A not equal to C
```

$f s=A$
opcode: 8A6yy
cycles: $13+d$ (GO/RTNYES) $6+d$ (NO)
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: 9aбyy
cycles: $13+d$ (GO/RTNYES)
$6+d$ (NO)

Test whether the fs field of $A$ is not equal to the fs field of $C$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?A<=B fs - Test for A less than or equal to B
fs = A opcode: 8BCyy
    cycles: }13+d\mathrm{ (GO/RTNYES)
                                6 + d (NO)
is = (P,WP,XS,X,S,M,B,W) opcode: 9bCyy
    cycles: }13+d\mathrm{ (GO/RTNYES)
                                6+d(NO)
```

Test whether the fs field of $A$ is less than or equal to the fs field of $B$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?A<B fs - Test for A less than B
----
fs = A opcode: 8B4yy
cycles: 13 + d (GO/RTNYES)
                                6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9b4yy
cycles: }13+d\mathrm{ (GO/RTNYES)
    6 + d (NO)
```

Test whether the fs field of $A$ is less than the fs field of $B$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?A=0 fs - Test for A equal to 0
fs = A opcode: 8A8yy
    cycles: 13 + d (GO/RTNYES)
    6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9a.8yy
    cycles: }13+d\mathrm{ (GO/RTNYES)
    6 + d (NO)
```

Test whether the fs field of $A$ is equal to 0 . Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?A=B fs - Test for A equal to B
fs =A opcode: 8AOyy
    cycles: }13+d\mathrm{ (GO/RTNYES)
        6 +d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9aOyy
cycles: 13 + d (GO/RTNYES)
    6 + d (NO)
```

Test whether the fs field of $A$ is equal to the fs field of B. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?A=C fs - Test for A equal to C
fs =A opcode: 8A2yy
cycles: }13+d\mathrm{ (GO/RTNYES)
    6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: Sa2yy
cycles: 13 + d (GO/RTNYES)
                                6 +d (NO)
```

Test whether the fs field of $A$ is equal to the fs field of $C$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?A>=B fs - Test for A greater than or equal to B
fs = A opcode: 8B8yy
    cycles: }13+d\mathrm{ (GO/RTNYES)
        6 +d (NO)
fs}=(P,WP,XS,X,S,M,B,W) opcode: 9b8yy
    cycles: 13 + d (GO/RINYES)
                                6 + d (NO)
```

Test whether the fs field of $A$ is greater than or equal. to the fs field of $B$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?A>B fs - Test for A greater than B
fs =A opcode: 8BOyy
    cycles: 13 + d (GO/RTNYES)
    6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9bOyy
    cycles: }13+d\mathrm{ (GO/RTNYES)
    4-17
```

```
6 + d (NO)
```

Test whether the fs field of $A$ is greater than the fs field of $B$. Must be followed by a.GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?B#O fs - Test for B not equal to 0
fs =A opcode: 8ADyy
    cycles: }13+d\mathrm{ (GO/RTNYES)
                                6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9aDyy
    cycles: }13+d\mathrm{ (GO/RTNYES)
    6 + d (NO)
```

Test whether the fs field of $B$ is not equal to 0 . Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.
?B\#A fs - Test for B not equal to A
$f s=A \quad$ opcode: 8A4yy
cycles: $13+d$ (GO/RINYES) $6+d$ (NO)
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: 9a4yy
cycles: $13+d$ (GO/RTNYES)

$$
6+d(\mathrm{NO})
$$

Test whether the fs field of $B$ is not equal to the fs field of $A$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RINYES or GOYES. Adjusts Carry.

```
?B#C fs - Test for B not equal to C
fs = A opcode: 8A5yy
cycles: 13 + d (GO/RTNYES)
    6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9a5yy
cycles: 13 + d (GO/RTNYES)
    6 + d (NO)
```

Test whether the fs field of $B$ is not equal to the $f$ field of $C$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.
? $B<=C$ fs - Test for $B$ less than or equal to $C$

fs = A opcode: 8BDyy
cycles: $13+d$ (GO/RTNYES) $6+d$ (NO)
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: 9bDyy
cycles: $13+d$ (GO/RTNYES)
$6+d$ (NO)
Test whether the fs field of $B$ is less than or equal to the fs field of C. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?B<C fs - Test for B less than C
fs = A opcode: 8B5yy
    cycles: 13 + d (GO/RTNYES)
    6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9b5yy
    cycles: 13 + d (GO/RTNYES)
    6 + d (NO)
```

Test whether the fs field of $B$ is less than the fs field of $C$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?B=0 fs - Test for B equal to 0
fs=A opcode: 8A9yy
cycles: 13 + d (GO/RTNYES)
                                6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9a9yy
cycles: }13+d\mathrm{ (GO/RINYES)
                                6 + d (NO)
Test whether the fs field of \(B\) is equal to 0 . Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.
```

```
?B=A fs - Test for B equal to A
```

?B=A fs - Test for B equal to A
fs = A opcode: 8AOyy
fs = A opcode: 8AOyy
cycles: }13+d\mathrm{ (GO/RTNYES)
cycles: }13+d\mathrm{ (GO/RTNYES)
6 + d (NO)
6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9aOyy
fs = (P,WP,XS,X,S,M,B,W) opcode: 9aOyy
cycles: }13+d\mathrm{ (GO/RINYES)
cycles: }13+d\mathrm{ (GO/RINYES)
6 + d (NO)

```
        6 + d (NO)
```

Test whether the fs field of $B$ is equal to the fs field of $A$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.
? $B=C$ fs - Test for $B$ equal to $C$
$f s=A$

```
opcode: 8A1yy
    cycles: }13+d\mathrm{ (GO/RTNYES)
                        6 + d (NO)
cycles: }13+d\mathrm{ (GO/RTNYES)
```

$f_{s}=(P, W P, X S, X, S, M, B, W) \quad$ opcode: 9alyy

$$
6+d \text { (NO) }
$$

Test whether the fs field of $B$ is equal to the fs field of $C$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?B>=C fs - Test for B greater than or equal to C
```



```
fs = A opcode: 8B9yy
    cycles: 13 + d (GO/RTNYES)
        6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9b9yy
    cycles: }13+d\mathrm{ (GO/RTNYES)
        6 + d (NO)
```

Test whether the fs field of $B$ is greater than or equal to the fs field of $C$. Must be followed by $a$ GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.
?B>C fs - Test for B greater than C
------

$$
f s=A
$$

opcode: 8B1yy
cycles: $13+d$ (GO/RTNYES) $6+d$ (NO)
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: 9b1yy cycles: $13+d$ (GO/RTNYES) $6+d$ (NO)

Test whether the fs field of $B$ is greater than the fis field of $C$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?C#O fs - Test for C not equal to 0
fs = A opcode: 8AEyy
    cycles: }13+d\mathrm{ (GO/RTNYES)
                                6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9aEyy
cycles: }13+d\mathrm{ (GO/RTNYES)
    6 +d (NO)
```

Test whether the fs field of $C$ is not equal to 0 . Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?C#A fs - Test for C not equal to A
fs =A opcode: 8abyy
    cycles: }13+d\mathrm{ (GO/RTNYES)
        6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9a6yy
    cycles: }13+d\mathrm{ (GO/RTNYES)
                        6 + d (NO)
```

Test whether the fs field of $C$ is not equal to the fs field of $A$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?C#B fs - Test for C not equal to B
fs =A opcode: 8A5yy
    cycles: 13 + d (GO/RTNYES)
    6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9a5yy
    cycles: }13+d\mathrm{ (GO/RTNYES)
    6 + d (NO)
```

Test whether the fs field of $C$ is not equal to the fs field of $B$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?C#D fs - Test for C not equal to D
-----
fs =A opcode: 8A7yy
cycles: }13+d\mathrm{ (GO/RTNYES)
    6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9a7yy
cycles: }13+d\mathrm{ (GO/RTNYES)
    6 + d (NO)
```

Test whether the fs field of $C$ is not equal to the fs field of $D$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?C<=A fs - Test for C less than or equal to A
-------
fs =A opcode: 8BEyy
    cycles: }13+d\mathrm{ (GO/RTNYES)
        6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9bEyy
    cycles: }13+d\mathrm{ (GO/RTNYES)
                        6 + d (NO)
Test whether the fs field of \(C\) is less than or equal to the \(f\) field of A. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.
```

```
?C<A fs - Test for C less than A
fs =A opcode: 8B6yy
    cycles: 1 13 + d (GO/RTNYES)
    6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 906yy
    cycles: }13+d\mathrm{ (GO/RTNYES)
```

$$
6+d \text { (NO) }
$$

Test whether the fs field of $C$ is less than the fs field of $A$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.
? $C=0$ fs - Test for $C$ equal to 0
$f s=A \quad$ opcode: 8AAyy
cycles: $13+d$ (GO/RTNYES)
$6+d$ (NO)
$f s=(P, W P, X S, X, S, M, B, W)$ opcode: 9aAyy cycles: $13+d$ (GO/RTNYES) $6+d$ (NO)

Test whether the fs field of $C$ is equal to 0 . Must be followed by a GOYES or RTNYES memonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?C=A fs - Test for C equal to A
fs =A opcode: 8A2yy
    cycles: }13+d\mathrm{ (GO/RTNYES)
        6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9a2yy
    cycles: }13+d\mathrm{ (GO/RTNYES)
        6 + d (NO)
```

Test whether the fs field of $\mathbb{C}$ is equal to the $f s f i e l d$ of $A$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?C=B fs - Test for C equal to B
fs =A opcode: 8A1yy
cycles: 13 + d (GO/RTNYES)
    6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9a1yy
cycles: }13+d\mathrm{ (GO/RTNYES)
    6 + d (NO)
```

Test whether the fs field of $C$ is equal to the fs field of $B$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?C=D fs - Test for C equal to D
```



```
fs =A opcode: 8A3yy
cycles: }13+d\mathrm{ (GO/RTNYES)
        6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9a3yy
cycles: }13+d\mathrm{ (GO/RTNYES)
                                6 + d (NO)
```

Test whether the fs field of $C$ is equal to the fs field of $D$. Must be followed by a GOYES or RTNYES mnemonic. yy j.s determined by the following RTNYES or GOYES. Adjusts Carry.

```
?C>=A fs - Test for C greater than or equal to A
fs = A opcode: 8BAyy
    cycles: }13+d(GO/RTNYES
    6+\alpha(NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9bAyy
    cycles: }13+d\mathrm{ (GO/RTNYES)
        6 + d (NO)
```

Test whether the $f s$ field of $C$ is greater than or equal to the fs field of A. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?C>A fs - Test for C greater than A
-------
fs = A opcode: 8B2yy
cycles: 13 + d (GO/RTNYES)
                                6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9b2yy
cycles: 13 + d (GO/RTNYES)
    6 + d (NO)
```

Test whether the fs field of $C$ is greater than the fs field of $A$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?D#O fs - Test for D not equal to 0
--------
fs = A opcode: 8AFyy
    cycles: 13 + d (GO/RTNYES)
    6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9aFyy
    cycles: 13 + d (GO/RTNYES)
    6 + d (NO)
```

Test whether the fs field of $D$ is not equal to 0 . Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?D#C Is - Test for D not equal to C
fs = A opcode: 8A7yy
    cycles: 13+d (GO/RTNYES)
        6 +d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9a'Tyy
    cycles: 13 + d (GO/RTNYES)
```

$$
6+d \text { (NO) }
$$

Test whether the fs field of $D$ is not equal to the fs field of $C$. Must be followed by a GOYES or RINYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?D<=C fs - Test for D less than or equal to C
---
fs=A opcode: 8BFyy
    cycles: }13+d\mathrm{ (GO/RINYES)
    6 + d (NO)
fs}=(P,WP,XS,X,S,M,B,W) opcode: 9bFyy
    cycles: }13+d\mathrm{ (GO/RTNYES)
                        6 + d (NO)
```

Test whether the fs field of $D$ is less than or equal to the fis field of $C$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.
? $D<C$ fs - Test for $D$ less than to $C$
$f s=A \quad$ opcode: 8B7yy
cycles: $13+d$ (GO/RTNYES) $6+d$ (NO)
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: 9b7yy
cycles: $13+d$ (GO/RTNYES)
$6+d(\mathrm{NO})$

Test whether the fs field of $D$ is less than the fs field of $C$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?D=0 fs - Test for D equal to 0
fs = A opcode: 8AByy
    cycles: }13+d\mathrm{ (GO/RTNYES)
        6 + d (NO)
fs}=(P,WP,XS,X,S,M,B,W) opcode: 9aByy
    cycles: 13 + d (GO/RTNYES)
        6 + d (NO)
```

Test whether the fs field of $D$ is equal to 0 . Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?D=C fs - Test for D equal to C
fs =A opcode: 8A3yy
    cycles: }13+d\mathrm{ (GO/RTNYES)
        6 + d (NO)
fs}=(P,WP,XS,X,S,M,B,W) opcode: 9a3yy
    cycles: }13+d\mathrm{ (GO/RTNYES)
        6 + d (NO)
```

Test whether the fs field of $D$ is equal to the fs field of $C$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?D>=C fs - Test for D greater than or equal to C
---------
fs =A opcode: 8BByy
    cycles: }13+d\mathrm{ (GO/RTNYES)
    6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opcode: 9bByy
    cycles: }13+d\mathrm{ (GO/RTNYES)
                        6+d (NO)
```

Test whether the fs field of $D$ is greater than or equal to the fs field of C. Must be followed by a GOYES or RINYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?D>C fs - Test for D greater than C
---------
fs=A opcode: 8B3yy
    cycles: }13+d\mathrm{ (GO/RTNYES)
    6 + d (NO)
fs = (P,WP,XS,X,S,M,B,W) opeode: 9b3yy
    cycles: }13+d\mathrm{ (GO/RTNYES)
    6 +d (NO)
```

Test whether the $f s$ field of $D$ is greater than the fs field of $C$. Must be followed by a GOYES or RTNYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.
?MP=0 - Test Module Pulled bit (MP)
$\qquad$
opcode: 838yy
cycles: 13 (GO/RTNYES)
6 (NO)
Test whether the Module Pulled bit (MP) is zero. This hardware status bit is set whenever a module-pulled interrupt occurs (the *INT line of the CPU is pulled high), and must be explictly cleared by the MP=0 mnemonic. See the "HP-71 Hardware Specification" for more information. Must be followed by a RTNYES or GOYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.
?P\# $n$ - Test if $P$ pointer not equal to $n$
opcode: 88nyy
cycles: $\begin{aligned} 13 & \text { (GO/RTNYES }) \\ 6 & \text { (NO) }\end{aligned}$

Test whether the $P$ pointer is not equal to $n$. Must be followed by a RTNYES or GOYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
PP= n - Test if P pointer is equal to n
-
    opcode: 89nyy
    cycles: 13 (GO/RTNYES)
        6 (NO)
```

Test whether the $P$ pointer is equal to $n$. Must be followed by a RTNYES or GOYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?SB=0 - Test Sticky Bit (SB)
```

$\qquad$
opcode: 832yy
cycles: 13 (GO/RTNYES)
6 (NO)

Test whether the Sticky Bit (SB) is zero. This hardware status bit is set on right shifts (circular or non-circular) when a non-zero nibble or bit is shifted off the end of the field. The Sticky Bit must be cleared explicitly by the $S B=0$ memonic. Must be followed by a RTNYES or GOYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.
?SR=0 - Test Service Request bit (SR) for zero
---------
opcode: 834yy
cycles: 13 (GO/RTNYES)
6 (NO)
Test whether the Service Request bit (SR) is zero. This hardware status bit is set by the SREQ? mnemonic, and must be cleared explicitly by the $S R=0$ mnemonic. Must be followed by a RTNYES or GOYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.
?ST\#0 n - Test status bit n not equal to 0

| opcode: 87 nyy |  |
| :---: | :---: |
| cycles: | 14 (GO/RTNYES) |
|  | $7($ NO $)$ |

Test whether Program Status bit $n$ is set. Must be followed by a RTNYES or GOYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.
?ST\#1 n - Test status bit $n$ not equal to 1
opcode: 86nyy
cycles: 14 (GO/RTNYES) 7 (NO)

Test whether Program Status bit $n$ is clear. Must be followed by a RTNYES or GOYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?ST=0 n - Test status bit n equal to 0
    opcode: 86nyy
    cycles: 14 (GO/RTNYES)
    7 (NO)
```

Test whether Program Status bit $n$ is clear. Must be followed by a RTNYES or GOYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
?ST=1 n - Test status bit n equal to 1
--------
    opcode: 87nyy
    cycles: 14 (GO/RTNYES)
        7(NO)
```

Test whether Program Status bit $n$ is set. Must be followed by a
RTNYES or GOYES mnemonic. yy is determined by the following RTNYES
or GOYES. Adjusts Carry.
? $\mathrm{XM}=0$ - Test External Module Missing bit (XM)
---------
opcode: 831yy
cycles: 13 (GO/RTNYES)
6 (NO)

Test the whether the External Module Missing bit (XM) is zero. This hardware status bit is set by the RTNSXM mnemonic, and must be explicitly cleared by the $X M=0$ mnemonic. Must be followed by a RTNYES or GOYES mnemonic. yy is determined by the following RTNYES or GOYES. Adjusts Carry.

```
A=-A fs - Two's complement of A into A
fs = A opcode: F8
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: Bb8
    cycles: }\quad3+
Complement the specified fs field of A. Complement is two's
complement if in HEX mode, ten's complement if in DEC mode. Carry
is set if the field is not zero, else Carry is cleared.
```

```
A=-A-1 fs - One's complement of A into A
---------
fs = A
    opcode: FC
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: BbC
    cycles: 3 + d
```

Perform a one's complement on the specified fs field of A. Carry
is always cleared.
$A=0$ is - Set $A$ equal to 0
fs $=A \quad$ opcode: DO
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: AbO
cycles: $3+d$
Set the specified fs field of $A$ to zero. Carry is not affected.
$A=A!B$ fs $-A O R B$ into $A$
---------
$\mathrm{fs}=\mathrm{A} \quad$ opcode: OEF8
cycles: $4+d$
$f_{s}=(P, W P, X S, X, S, M, B, W) \quad$ opcode: OEa 8
cycles: $4+d$

Set the fs field of register $A$ to its logical $O R$ with the corresponding field of register B. Carry is not affected.

```
A=A!C fs - AOR C into A
fs = A opcode: OEFE
    cycles: 4 + d
fs = (P,WP,XS,X,S,M,B,W) opcode: OEaE
    cycles: 4 + d
```

Set the fsfield of register $A$ to its logical $O R$ with the
corresponding field of register $C$. Carry is not affected.


Set the fs field of register $A$ to its logical AND with the corresponding field of register B. Carry is not affected.

```
A=A&C Ss - A AND C into A
fs =A opcode: OEF6
    cycles: 4 +d
fs}=(P,WP,XS,X,S,M,B,W) opcode: OEa.6
    cycles: }4+
```

Set the fs field of register $A$ to its logical AND with the corresponding field of register C. Carry is not affected.

```
A=A+1 fs - Increment A
fs = A opcode: E4
cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: Ba.4
cycles: }3+
Increment the specified fs field of register A by one. Adjusts Carry.
```

$A=A+A$ fs - Sum of $A$ and $A$ into $A$
fs $=\mathbf{A}$ opcode: C4 cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: Aa4 cycles: $\quad 3+d$

Double the specified fs field of register A. Adjusts Carry.
$A=A+B$ fs - Sum of $A$ and $B$ into $A$
$f s=A \quad$ opcode: CO
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: AaO
cycles: $\quad 3+d$
Set the specified fs field of register $A$ to the sum of itself and the corresponding field of register B. Adjusts Carry.

```
A=A+C fs - Sum of A and C into A
--------
fs =A opcode: CA
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W\rangle opcode: AaA
    cycles: }3+
```

Set the specified fs field of register $A$ to the sum of itself and the corresponding field of register C. Adjusts Carry.

```
A=A-1 fs - Decrement A
--------
fs = A opcode: CC
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: AaC
    cycles: }3+
```

Decrement the specified fs field of register A by one. Adjusts
Carry .
$A=A-B$ fs $-A$ minus $B$ into $A$
---------
fs =A opcode: EO
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W\rangle \quad$ opcode: BaO
cycles: $\quad 3+d$

Set the specified fs field of register $A$ to the difference between itself and the corresponding field of register B. Adjusts Carry.

```
A=A-C fs - A minus C into A
fs=A opcode: EA
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: BaA
    cycles: }3+
```

Set the specified fs field of register $A$ to the difference between itself and the corresponding field of register C. Adjusts Carry.

```
A=B fs - Copy B to A
fs = A opcode: D4
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: Ab4
    cycles: }3+
Copy the fs field of register B into the corresponding field of
register A. Carry is not affected.
```

$A=B-A$ fs $\quad B$ minus $A$ into $A$
---------
$f s=A$
opcode: EC
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: BaC
cycles: $\quad 3+d$

Set the specified fs field of register $A$ to the inverse difference between itself and the corresponding field of register B. Adjusts Carry.

```
A=C fs - Copy C to A
fs =A opcode: DA
cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: AbA
cycles: }3+
Copy the fs field of register C into the corresponding field of
register A. Carry is not affected.
A=DATO fsd - Load A from memory
---------
fs=A opcode: 142
    cycles: 18
fs = B opcode: 14A
    cycles: 15
fs = (P,WP,XS,X,S,M,W) opcode: 252a
    cycles: 17 + d
fs = d opcode: 15Ax (x=d-1)
    cycles: }16+
The amount of data (d nibbles) specified by fsd will be transferred from the memory address pointed to by DO into the specified field of register \(A\). The lowest-addressed nibble will be transferred into the lowest-order nibble of the register field, proceeding toward the higher-order nibbles. If fs \(=d\), d nibbles are transferred into the register starting at nibble 0. See the section on "Loading Data From Memory" earlier in this chapter.
```

| A=DAT1 fisd - Load A from memory |  |  |
| :---: | :---: | :---: |
| $f s=A$ | opcode: | 143 |
|  | cycles: | 18 |
| $f s=B$ | opcode: | 14B |
|  | cycles: | 15 |

```
fs = (P,WP,XS,X,S,M,W) opcode: 153a
cycles: 17 + d
fs = d opcode: 15Bx (x=d-1)
cycles: 16 + d
```

The amount of data (d nibbles) specified by fsd will be transferred from the memory address pointed to by D1 into the specified field of register $A$. The lowest-addressed nibble will be transferred into the lowest-order nibble of the register field, proceeding toward the higher-order nibbles. If $f s=d$, $d$ nibbles are transferred into the register starting at nibble 0 . See the section on "Loading Data From Memory" earlier in this chapter.
$A=I N \quad-L o a d A$ with IN
$\qquad$
opcode: 802
cycles: $\quad 7$
Load the low-order 4 nibbles of the $A$ register with the contents of the Input register.
$A=R 0 \quad$ - Copy RO to $A$
$\qquad$
opcode: 110
cycles: 19
The contents of the scratch register $R O$ is copied to the working register A.

```
A=R1 - Copy R1 to A
```

```
---------
```

    opcode: 111
    cycles: 19
    The contents of the scratch register $R 1$ is copied to the working register A．
$A=R 2$－Copy R2 to A

The contents of the scratch register R 2 is copied to the working register A．
$A=$ R3－Copy R3 to A
$\qquad$
opcode： 113
cycles： 19
The contents of the scratch register $R 3$ is copied to the working register A．
$A=R 4 \quad$－Copy R4 to $A$
$\qquad$
opcode： 114
cycles： 19
The contents of the scratch register $R 4$ is copied to the working register A．

```
ABEX fs - Exchange Registers A and B
---------
fs = A opcode: DC
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: AbC
    cycles: }3+
Exchange the fs fields of registers of A and B. Carry is not
affected.
ACEX fs - Exchange Registers A and C
---------
fs \(=\mathrm{A} \quad\) opcode: DE
cycles: 7
\(\mathrm{fs}=(\mathrm{P}, \mathrm{WP}, \mathrm{XS}, \mathrm{X}, \mathrm{S}, \mathrm{M}, \mathrm{B}, \mathrm{W}) \quad \begin{aligned} & \text { opcode: } \mathrm{AbE} \\ & \text { cycles: }\end{aligned} \quad 3+\mathrm{d}\)
Exchange the fs fields of registers of A and C. Carry is not affected.
```

ADOEX - Exciange $A$ and DO (nibs 0-4)
---------
opcode: 132
cycles: 8
Exchange the A field of register A with Data pointer DO. Carry is not affected.

```
ADOXS - Exchange A and DO short (nibs 0-3)
---------
    opcode: 13A
    cycles: 7
```

Exchange the lower 4 nibbles of $A$ with the lower 4 nibbles of Data pointer DO. Carry is not affected.

```
AD1EX - Exchange A and D1 (nibs 0-4)
```

---------
opcode: 133
cycles: 8

Exchange the A field of register A with Data pointer D1. Carry is not affected.

AD1XS - Exchange $A$ and D1 short (nibs 0-3)
---------
opcode: 13B cycles: 7

Exchange the lower 4 nibbles of $A$ with the lower 4 nibbles of Data pointer D1. Carry is not affected.

AROEX - Exchange A and RO
---------
opcode: 120
cycles: 19
Exchange the contents of the working register $A$ and the scratch register RO.
AR1EX - Exchange A and R1
-
opcode: 121
cycles: 19
Exchange the contents of the working register $A$ and the scratch register R1.
AR2EX - Exchange $A$ and R2
--------
opcode: 122
cycles: 19
Exchange the contents of the working register $A$ and the scratchregister R2.
AR3EX - Exchange A and R3
opcode ..... 123
cycles: ..... 19-
Exchange the contents of the working register $A$ and the scratchregister R3.
ARLEX - Exchange $A$ and $R 4$
opcode ..... 124
cycles: ..... 19Exchange the contents of the working register $A$ and the scratchregister R4.

ASL fs - A Shift Left

```
opcode: FO
cycles: 7
cycles: }3+
```

$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: BbO

Shift the contents of the specified fs field of register A left one nibble, without affecting the rest of the register. The nibble shifted off the left end of the field is lost. The new low-order nibble of the field is zero. The Sticky Bit (SB) is not affected.

ASLC - A Shift Left Circular
---------
opcode: 810
cycles: 21
Circular shift register $A$ left one $n i b b l e$. Operates on all 16 digits. The Sticky Bit (SB) is not affected.

ASR fs - A Shift Right
fs $=\mathrm{A} \quad$ opcode: F4
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: Bb4
cycles: $\quad 3+d$
Shift the contents of the specified fs field of register $A$ right one nibble, without affecting the rest of the register. The nibble shifted off the right end of the field is lost, but the Sticky Bit (SB) is set if the nibble was non-zero. The new high-order nibble of the field is zero.

## ASRB - A Shift Right Bit

$\begin{array}{lr}\text { opcode: } & 81 \mathrm{C} \\ \text { cycles: } & 20\end{array}$

Shift register A right one bit. Operates on all 16 digits. The bit shifted off the end is lost, but the Sticky Bit (SB) is set if it was non-zero. The new high-order bit of the register is zero.

## ASRC - A Shift Right Circular

$\qquad$
opcode: 814
cycles: 21
Circular shift register A right one nibble. Operates on all 16 digits. The Sticky Bit (SB) is set if the nibble shifted from low-order around to high-order position was non-zero.
$B=-B$ fs - Two's complement of $B$ into $B$
---------
$\mathrm{f} s=\mathrm{A} \quad$ opcode: F9
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: Bb 9
cycles: $\quad 3+d$
Complement the specified fs field of $B$, Complement is two's complement if in HEX mode, ten's complement if in DEC mode. Carry is set if the field is not zero, else Carry is cleared.

```
B=-B-1 fs - One's complement of B into B
--------
fs = A opcode: FD
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: BbD
    cycles: 3 + d
```

Perform a one's complement on the specified fs field of B. Carry is always cleared.

```
B=0 fs - Set B equal to 0
fs = A opcode: D1
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: Ab1
    cycles: }3+
```

Set the specified fs field of $B$ to zero. Carry is not affected.
$B=A \quad f s \quad$ - Copy $A$ to $B$
$\mathrm{fs}=\mathrm{A} \quad$ opcode: D8
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: Ab8
cycles: 3 + d

Copy the fs field of register $A$ into the corresponding field of register B. Carry is not affected.

```
B=B!A fs - B OR A into B
fs = A opcode: OEFC
    cycles: 4 + d
fs = (P,WP,XS,X,S,M,B,W) opcode: OEaC
    cycles: 4 + d
```

Set the fs field of register $B$ to its logical $O R$ with the
corresponding field of register A. Carry is not affected.

```
B=B!C fs - B OR C into B
fs = A opcode: OEF9
    cycles: 4 + d
fs = (P,WP,XS,X,S,M,B,W) opcode: OEa9
cycles: 4 + d
```

Set the fs field of register $B$ to its logical OR with the corresponding field of register C. Carry is not affected.
$B=B \& A$ fs - $B$ AND A into $B$

| $\mathrm{fs}=\mathrm{A}$ | opcode: cycles: | ${ }^{\text {OEF4 }} 4+\mathrm{d}$ |
| :---: | :---: | :---: |
| $f s=(P, W P, X S, X, S, M, B, W)$ | opcode: <br> cycles: | $\begin{aligned} & 0 \mathrm{Ea} 2 \\ & 4+\mathrm{d} \end{aligned}$ |

Set the fs field of register $B$ to its logical AND with the corresponding field of register A. Carry is not affected.

```
B=B&C fs - B AND C into B
-
fs = A opcode: OEF1
    cycles: 4 + d
fs = (P,WP,XS,X,S,M,B,W) opcode: OEal
    cycles: 4 + d
```

Set the fs field of register $B$ to its logical AND with the
corresponding field of register C. Carry is not affected.
$B=B+1$ fs - Increment $B$
$\mathrm{fs}=\mathrm{A}$
opcode: E5
cycles: 7
fs $=(P, W P, X S, X, S, M, B, W) \quad$ opcode: Ba5
cycles: $3+d$
Increment the specified fs field of register $B$ by one. Adjusts
Carry.

```
B=B+A fs - Sum of B and A into B
fs = A opcode: C8
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: Aa8
    cycles: 3 + d
```

Set the specified fs field of register $B$ to the sum of itself and the corresponding field of register A. Adjusts Carry.

```
B=B+B fs - Sum of B and B into B
fs = A opcode: C5
    cycles: 7
fs}=(P,WP,XS,X,S,M,B,W) opcode: Aa.
    cycles: }3+
```

Double the specified fs field of register B. Adjusts Carry.

```
B=B+C fs - Sum of B and C into B
--------
fs = A opcode: C1
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: Aal
    cycles: }3+
```

Set the specified is field of register $B$ to the sum of itself and
the corresponding field of register $C$. Adjusts Carry.
$B=B-1$ fs - Decrement B
$f s=A \quad$ opcode: CD
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: AaD
cycles: $\quad 3+d$
Decrement the specified is field of register $B$ by one. Adjusts
Carry.

```
B=B-A fs - B minus A into B
fs = A opcode: E8
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: Ba8
    cycles: }3+
```

Set the specified fs field of register $B$ to the difference between itself and the corresponding field of register A. Adjusts Carry.

```
B=B-C fs - B minus C into B
fs=A opcode: E1
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: Bal
    cycles: }3+
```

Set the specified fs field of register $B$ to the difference between
j.tself and the corresponding field of register C. Adjusts Carry.
$B=C$ fs - Copy $C$ to $B$
fs $=\mathrm{A} \quad$ opcode: D5
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: Ab5
cycles: $3+d$

Copy the fs field of register $C$ into the corresponding field of register B. Carry is not affected.

```
B=C-B fs - C minus B into B
fs = A opcode: ED
cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: BaD
cycles: 3 + d
```

Set the specified fs field of register $B$ to the inverse difference between itself and the corresponding field of register C. Adjusts Carry.

BAEX fs - Exchange Registers $B$ and $A$
------
$f s=A \quad$ opcode: DC cycles: $\quad 7$
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: AbC cycles: $\quad 3+d$

Exchange the fis fields of registers of $B$ and $A$. Carry is not affected.

BCEX fs - Exchange Registers B and C
fs $=A \quad$ opcode: DD
cycles: $\quad 7$
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: AbD
cycles: $\quad 3+d$
Exchange the fs fields of registers of $B$ and $C$. Carry is not affected.

```
BSL fs - B Shift Left
---------
fs = A opcode: F1
cycles: 7
fs}=(P,WP,XS,X,S,M,B,W) opcode: Bb
cycles: 3 + d
```

Shift the contents of the specified fs field of register $B$ left one nibble, without affecting the rest of the register. The nibble shifted off the left end of the field is lost. The new low-order nibble of the field is zero. The Sticky Bit (SB) is not affected.

BSLC - B Shift Left Circular

```
    opcode: }81
    cycles: 21
Circular shift register B left one nibble. Operates on all 16
digits. The Sticky Bit (SB) is not affected.
```

BSR fs - B Shift Right
$\mathrm{f}_{\mathrm{s}}=\mathrm{A} \quad$ opcode: F5
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: Bb5
cycles: $\quad 3+d$
Shift the contents of the specified fs field of register B right
one nibble, without affecting the rest of the register. The nibble
shifted off the right end of the field is lost, but the Sticky Bit
(SB) is set if the nibble was non-zero. The new high-order nibble
of the field is zero.

Shift register B right one bit. Operates on all 16 digits. The bit shifted off the end is lost, but the Sticky Bit (SB) is set if it was non-zero. The new high-order bit of the register is zero.

BSRC - B Shift Right Circular
$\qquad$
opcode: 815
cycles: 21
Circular shift register $B$ right one nibble. Operates on all 16 digits. The Sticky Bit (SB) is set if the nibble shifted from low-order around to high-order position was non-zero.

BUSCC - Bus Command " C "
$\qquad$
$\begin{array}{lr}\text { opcode: } & 80 \mathrm{~B} \\ \text { cycles: } & 6\end{array}$
Enters the HP-71 bus command "C" onto the system bus (this command is reserved for later use). No other operation is performed. See the 'HP-71 Hardware Specification' for more information.
$C+P+1$ - Increment $C$ by One Plus $P$ Pointer
---------
opcode: 809
cycles: 8

The A field of the $C$ register is incremented by one plus the value of the $P$ pointer. This instruction is always executed in HEX mode. Adjusts Carry.

```
C=-C fs - Two's complement of C into C
fs = A opcode: FA
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: BbA
    cycles: }3+
Complement the specified fs field of C. Complement is two's complement if in HEX mode, ten's complement if in DEC mode. Carry is set if the field is not zero, else Carry is cleared.
```

```
C=-C-1 fs - One's complement of C into C
fs=A opcode: FE
    cycles: 7
fs}=(P,WP,XS,X,S,M,B,W) opcode: BbE
    cycles: }3+
```

Perform a one's complement on the specified fs field of $C$. Carry is always cleared.

```
C=0 fs - Set C equal to 0
fs = A opcode: D2
    cycles: 7
fs}=(P,WP,XS,X,S,M,B,W) opcode: Ab
    cycles: }3+
```

Set the spocified fs field of $C$ to zero. Carry is not affected.
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```
C=A fs - Copy A to C
fs =A opcode: D6
cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: Ab6
    cycles: 3 + d
Copy the fs field of register A into the corresponding field of
register C. Carry is not affected.
C=A-C fs - A minus C into C
fs = A opcode: EE
cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: BaE
cycles: 3 + d
Set the specified fs field of register \(C\) to the inverse difference between itself and the corresponding field of register A. Adjusts Carry.
\(\mathrm{C}=\mathrm{B} \quad \mathrm{fs}\) - Copy B to C
\(\mathrm{fs}=\mathrm{A} \quad\) opcode: D9
cycles: 7
\(\mathrm{fs}=(\mathrm{P}, \mathrm{WP}, \mathrm{XS}, \mathrm{X}, \mathrm{S}, \mathrm{M}, \mathrm{B}, \mathrm{W}) \quad\) opcode: Ab9
cycles: \(3+d\)
Copy the fs field of register \(B\) into the corresponding field of register C. Carry is not affected.
```

```
C=C!A fs - C OR A into C
--------
fs = A opcode: OEFA
    cycles: 4 + d
fs}=(P,WP,XS,X,S,M,B,W) opcode: OEaA
    cycles: }4+
Set the fs field of register C to its logical OR with the
corresponding field of register A. Carry is not affected.
C=C!B fs - C OR B into C
fs =A opcode: OEFD
    cycles: }4+
fs = (P,WP,XS,X,S,M,B,W) opcode: OEaD
    cycles: }4+
Set the fs field of register C to its logical OR with the
corresponding field of register B. Carry is not affected.
C=C!D fs - C OR D into C
fs}=\textrm{A
opcode: OEFF
cycles: }4+
fs}=(P,WP,XS,X,S,M,B,W) opcode: OEa
    cycles: 4 + d
Set the fs field of register \(C\) to its logical \(O R\) with the corresponding field of register D. Carry is not affected.
```

```
C=C&A fs - C AND A into C
fs A A opcode: UEF2
    cycles: 4 + d
fs = (P,WP,XS,X,S,M,B,W) opcode: OEa2
    cycles: 4 + d
```

Set the fs field of register $C$ to its logical AND with the
corresponding field of register $A$. Carry is not affected.
$C=C \& B$ fs $\quad C$ AND $B$ into $C$
---------
$\mathrm{fs}=\mathrm{A} \quad$ opcode: OEF5
cycles: $\quad 4+d$
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: OEa5
cycles: $\quad 4+d$
Set the fs field of register $C$ to its logical AND with the
corresponding field of register $B$. Carry is not affected.
$C=C \& D$ fs $-C$ AND D into $C$
---------
Is $=A \quad$ opcode: OEF7
cycles: $\quad 4+d$
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: OEa7
cycles: $\quad 4+d$

Set the fs field of register $C$ to its logical AND with the corresponding field of register D. Carry is not affected.

```
C=C+1 fs - Increment C
----
fs = A
fs = (P,WP,XS,X,S,M,B,W) opcode: Ba\sigma
opcode: E6
cycles: 7
cycles: }3+
Increment the specified fs field of register C by one. Adjusts
Carry.
```

$C=C+A$ fs $\quad-$ Sum of $C$ and $A$ into $C$
$f_{s}=A \quad$ opcode: C2
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: Aa2
cycles: $\quad 3+d$

Set the specified fs field of register $C$ to the sum of itself and the corresponding field of register A. Adjusts Carry.

```
C=C+B fs - Sum of C and B into C
fs = A opcode: C9
cycles: 7
f's = (P,WP,XS,X,S,M,B,W) opcode: Aa9
    cycles: 3 + d
```

Set the specified fs field of register $C$ to the sum of itself and the corresponding field of register B. Adjusts Carry.

```
C=C+C fs - Sum of C and C into C
-------
fs = A opcode: C6
    cycles: 7
fs}=(P,WP,XS,X,S,M,B,W) opcode: Aa
    cycles: }3+
Double the specified fs field of register C. Adjusts Carry.
C=C+D fs - Sum of C and D into C
---------
fs=A opcode: CB
    cycles: 7
fs}=(P,WP,XS,X,S,M,B,W) opcode: Aa
    cycles: }3+
Set the specified fs field of register C to the sum of itself and
the corresponding field of register D. Adjusts Carry.
```

$C=C-1$ fs - Decrement $C$
$\mathrm{fs}=\mathrm{A} \quad$ opcode: CE
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: AaE
cycles: $\quad 3+d$
Decrement the specified fs field of register $C$ by one. Adjusts
Carry.

```
C=C-A fs - C minus A into C
--------
fs = A opcode: E2
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: Ba2
    cycles: }3+
Set the specified fs field of register C to the difference between
itself and the corresponding field of register A. Adjusts Carry.
```

```
C=C-B fs - C minus B into C
--------
fs =A opcode: E9
    cycles: 7
fs}=(P,WP,XS,X,S,M,B,W) opcode: Ba9,
    cycles: }3+
```

Set the specified fs field of register $C$ to the difference between
itself and the corresponding field of register B. Adjusts Carry.
$C=C-D$ fs - $C$ minus $D$ into $C$
-------
$f s=A \quad$ opcode: EB
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: BaB
cycles: $3+d$

Set the specified fs field of register $C$ to the difference between itself and the corresponding field of register D. Adjusts Carry.

```
C=D fs - Copy D to C
---------
fs}=\mathbf{A}\mathrm{ opcode: DB
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: AbB
    cycles: }3+
Copy the fs field of register D into the corresponding field of
register C. Carry is not affected.
C=DATO fsd - Load C from memory
---------
fs=A opcode: 146
    cycles: 18
fs = B opcode: 14E
    cycles: 15
fs = (P,WP,XS,X,S,M,W) opcode: 156a
    cycles: 17 + d
fs = d opcode: 15Ex (x=d-1)
    cycles: 16 + d
The amount of data (d nibbles) specified by fsd will be transferred from the memory address pointed to by DO into the specified field of register \(C\). The lowest-addressed nibble will be transferred into the lowest-order nibble of the register field, proceeding toward the higher-order nibbles. If fs \(=d\), \(d\) nibbles are transferred into the register starting at nibble 0 . See the section on "Loading Data From Memory" earlier in this chapter.
```

```
C=DAT1 fsd - Load C from memory
---------
fs = A opcode: 147
cycles: 18
fs = B opcode: 14F
cycles: 15
fs = (P,WP,XS,X,S,M,W) opcode: 157a
cycles: }17+
fs = d opcode: 15Fx (x=d-1)
cycles: }16+
```

The amount of data (d nibbles) specified by fsd will be transferred from the memory address pointed to by D1 into the specified field of register $C$. The lowest-addressed nibble will be transferred into the lowest-order nibble of the register field, proceeding toward the higher-order nibbles. If fs $=d$, $d$ nibbles are transferred into the register starting at nibble 0. See the section on "Loading Data From Memory" earlier in this chapter.

```
C=ID - Request chip ID
```

---------

$$
\text { opcode: } 806
$$

cycles: 11
The chip which has its DAISX-IN line high and its configuration flag low will send its 5 nibble ID register to the system bus which will be loaded into the low-order 5 nibbles (A field) of the $C$ register. See the "HP-71 Hardware Specification" for more information.
$C=I N \quad-$ Load $C$ with IN
opcode: 803
cycles: 7
Load the low-order 4 nibbles of the $C$ register with the contents of the Input register. See the "HP-71 Hardware Specification" for more information.

```
C=P n - Copy P Pointer into Nibble n of C
--------
opcode: 80Cn
cycles: 6
```

Copy $P$ pointer into $C$ register at digit position specified by $n$.
C=RO - Copy RO to C
opcode: 118
cycles: 19
The contents of the scratch register $R 0$ is copied to the working
register C.
$\mathrm{C}=\mathrm{R} 1 \quad$ - Copy R1 to C
opcode: 119
cycles: 19
I'he contents of the scratch register $R 1$ is copied to the working
register C.
$C=R 2 \quad$ - Copy R2 to $C$
--------
opcode: 11A
cycles: 19
The contents of the scratch register $R 2$ is copied to the working
register C.

```
C=R3 - Copy R3 to C
```

---------
opcode: 11B
cycles: 19
The contents of the scratch register $R 3$ is copied to the working
register C.
$C=R 4 \quad-\quad$ Copy 44 to $C$
opcode: 11C
cycles: 19
The contents of the scratch register $R 4$ is copied to the working
register C.
$\mathrm{C}=\mathrm{RSTK} \quad-$ Pop stack to C
opcode: 07
cycles: 8

Pop the top-most address off of the hardware return stack, placing the address in the lower 5 nibbles (A fijeld) of register $C$. The high-order nibbles of $C$ are unchanged. As the address is popped from the return stack, a zero address is inserted ai the bottom of the stack. Compare with the RTN mnemonic.

```
C=ST - Status to C
---------
    opcode: 09
    cycles: 6
Copy the low-order 12 bits of the status register into the
low-order 12 bits (X field) of the C register.
```

CAEX fs - Exchange Registers $C$ and $A$
---------
fs $=A \quad$ opcode: DE
cycles: 7
fs $=(P, W P, X S, X, S, M, B, W) \quad$ opcode: AbE
cycles: $3+\mathrm{d}$
Exchange the fs fields of registers of $C$ and $A$. Carry is not
affected.
CBEX fs - Exchange Registers C and B
--------
fs $=A \quad$ opcode: DD
cycles: 7
$f_{s}=(P, W P, X S, X, S, M, B, W) \quad$ opcode: AbD
cycles: $\quad 3+\mathrm{d}$
Exchange the fs fields of registers of $C$ and $B$. Carry is not
affected.

```
CDOEX - Exchange C and DO (nibs 0-4)
---------
    opcode: 136
    cycles: 8
```

Exchange the A field of register $C$ with Data pointer DO. Carry is
not affected.

```
CDOXS - Exchange C and DO short (nibs 0-3)
---------
    opcode: 13E
    cycles: 7
Exchange tive lower 4 nibbles of C with the lower 4 nibbles of Data
pointer DO. Carry is not affected.
```

CDIEX - Exchange $C$ and D1 (nibs 0-4)
---------
opcode: 137
cycles: 8

Exchange the A field of register C with Data pointer D1. Carry is not affected.

CD1XS - Exchange C and D1 short (nibs 0-3)
opcode: 13F
cycles: 7
Exchange the lower 4 nibbles of $C$ with the lower 4 nibbles of Data pointer D1. Carry is not affected.

```
CDEX fs - Exchange Registers C and D
--
fs = A opcode: DF
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: AbF
    cycles: }3+
Exchange the fs fields of registers of C and D. Carry is not
affected.
```

CLRHST - Clear Hardware Status bits
---------
opcode: 82F
cycles: 3

Clears the 4 Hardware Status bits $X M, S B, S R$ and MP. Note that the opcode is actually $82 x$, where $x$ is merely a mask for which Hardware Status bits to clear, as follows:
bit 0 - External Module Missing bit (see $X M=0$ mnemonic)
bit 1 - Sticky Bit (see $\mathrm{SB}=0$ mnemonic)
bit 2 - Service Request bit (see $\mathrm{SR}=0$ mnemonic)
bit 3 - Module Pulled bit (see MP=0 mnemonic)
For example opcode 829 clears $X M$ and MP. Although there is no mnemonic for this, the opcode can be inserted into the code by using, for example, NIBHEX 829.

CLRST - Clear Program Status
$\qquad$
opcode: 08
cycles: $\quad 6$
Clear the low-order 12 bits (S0 through S11) of the Program Status register ST .

```
CONFIG - Configure
```

opcode: 805
cycles: 11

Copy the low-order 5 nibbles (A field) of the $C$ register into the Configuration register of the chip which has its DAISY-IN line high and its configuration flag low. See the "HP-71 Hardware Specification" for information.

CPEX n - Exchange Nibble n of C With P Pointer
$\qquad$ opcode: 80Fn cycles: 6

Exchange the $P$ pointer with digit $n$ of the $C$ register.

CROEX - Exchange $C$ and RO
opcode: 128
cycles: 19
Exchange the contents of the working register $C$ and the scratch register RO.

CR1EX - Exchange $C$ and R1
$\qquad$
opcode: 129
cycles: 19
Exchange the contents of the working register $C$ and the scratch register R1.

```
CR2EX - Exchange C and R2
---------
```

opcode: 12A
cycles: 19

Exchange the contents of the working register $C$ and the scratch register R2.

## CR3EX <br> - Exchange C and R3

$\qquad$
opcode: 12B
cycles: 19
Exchange the contents of the working register $C$ and the scratch register R3.

CR4EX - Exchange $C$ and R4
----------
opcode: 12C
cycles: 19

Exchange the contents of the working register $C$ and the scratch register R4.

CSL fs - C Shift Left
$f_{s}=A \quad$ opcode: F2
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: Bb2
cycles: $\quad 3+d$
Shift the contents of the specified fs field of register $C$ left one nibble, without affecting the rest of the register. The nibble shifted off the left end of the field is lost. The new low-order nibble of the field is zero. The Sticky Bit (SB) is not affected.

```
CSLC - C Shift Left Circular
```

opcode: 812
cycles: 21
Circular shift register $C$ left one nibble. Operates on all 16
digits. The Sticky Bit (SB) is not affected.
CSR fs - C Shift Right
$f s=A \quad$ opcode: F6
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: Bb6
cycles: $\quad 3+d$

Shift the contents of the specified fs field of register $C$ right one nibble, without affecting the rest of the register. The nibble shifted off the right end of the field is lost, but the Sticky Bit (SB) is set if the nibble was non-zero. The new high-order nibble of the field is zero.

CSRB - C Shift Right Bit

```
    opcode: 81E
    cycles: 20
Shift register C right one bit. Operates on all 16 digits. The
bit shifted off the end is lost, but the Sticky Bit (SB) is set if
it was non-zero. The new high-order bit of the register is zero.
```

```
CSRC - C Shift Right Circular
```

-------- opcode: 816
cycles: 21

Circular shift register C right one nibble. Operates on all 16 digits. The Sticky Bit (SB) is set if the nibble shifted from low-order around to high-order position was non-zero.

CSTEX - Exchange Status with C register
$\qquad$
opcode: OB
cycles: 6
Exchange the low-order 12 bits (SO through S11) of the Program Status register ST with the low-order 12 bits ( X field) of the C register.

DO=(2) nn - Load 2 Nibbles Into Do
-----------
opcode: $\underset{4}{19 n n}$
cycles: 4
Load the low-order two nibbles of DO with nn. The upper nibbles of DO remain unchanged. Any overflow is ignored by the assembler. The assembled digits of nn are stored in the opcode in reverse order so that when the instruction is executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.

D0=(4) nnnn - Load 4 Nibbles Into DO
------------
opcode: 1Annnn cycles: 6

Load the low-order four nibbles of DO with nnnn. The upper nibble of DO remains unchanged. Any overflow is ignored by the assembler. The assembled digits of nnnn are stored in the opcode in reverse order so that when the instruction is executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.
$\mathrm{DO}=(5)$ nnnnn - Load 5 Nibbles Into DO
------------
opcode: 1Bnnnnn
cycles:
7

Load all five nibbles of DO with nnnnn. Any overflow is ignored by the assembler. The assembled digits of nnnnn are stored in the opcode in reverse order so that when the instruction is executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.

DO=A - Copy A to DO (nibs 0-4)
-----.-.
opcode: 130
cycles: 8

The A field of register A is copied into Data pointer register DO. Carry is not affected.

```
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```

DO=AS - Copy A to DO short (nibs 0-3)
--------
opcode: 138
cycles: 7

The lower 4 nibbles of $A$ are copied into the lower 4 nibbles of Data pointer register DO. Carry is not affected.

DO=C $\quad$ - Copy $C$ to DO (nibs 0-4)
opcode: 134
cycles: 8
The A field of register $C$ is copied into Data pointer register DO. Carry is not affected.

DO=CS - Copy $C$ to DO short (nibs 0-3)
$\begin{array}{rrr}\text { opcode: } & 13 C \\ \text { cycles: } & 7\end{array}$

The lower 4 nibbles of $C$ are copied into the lower 4 nibbles of Data pointer register DO. Carry is not affected.
$D O=D O+n \quad-A d d n$ to $D O \quad(1<=n<=16)$
---------
opcode: 16x ( $x=n-1$ )
cycles: 7
Increment DO by n. This instruction is always executed in HEX mode. Adjusts Carry.

```
DO=DO- n - Subtract n from DO (1<=n<=16)
```

---------
opcode: $18 x(x=n-1)$
cycles: 7

Decrement DO by n. This instruction is always executed in HEX mode. Adjusts Carry.

DO=HEX hh - Load DO with hex constant hh
$\qquad$
opcode: 19hh
cycles: 4
Load the low-order two nibbles of $D O$ with the hex constant hh. The upper nibbles of DO remain unchanged. The digits of hh are stored in the opcode in reverse order so that when the instruction is executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.

DO =HEX hhhh - Load DO with hex constant hhhh
$\qquad$
opcode: J.Ahhhh
cycles: 6

Load the low-order four nibbles of DO with the hex constant hhhh. The upper nibble of DO remains unchanged. The digits of hhhh are stored in the opcode in reverse order so that when the instruction is executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.

DO=HEX hhhhh - Load DO with hex constant hhhhh
$\qquad$
opcode: 1Bhhhhh
cycles: 7
Load all five nibbles of DO with the hex constant hhhhh. The digits of hhhhh are stored in the opcode in reverse order so that when the instruction is executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.

D1=(2) nn - Load 2 Nibbles Into D1
--------

| opcode: | 1 Dnn |
| :--- | :---: |
| cycles: | 4 |

Load the low-order two nibbles of D1 with nn. The upper nibbles of D1 remain unchanged. Any overflow is ignored by the assembler. The assembled digits of $n n$ are stored in the opcode in reverse order so that when the instruction is executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.

D1=(4) nnnn - Load 4 Nibbles Into D1
$\qquad$
opcode: 1Ennnn
cycles: 6
Load the low-order four nibbles of $D 1$ with nnnn. The upper nibble of D1 remains unchanged. Any overflow is ignored by the assembler. The assembled digits of nnnn are stored in the opcode in reverse order so that when the instruction is executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.

D1=(5) nnnnn - Load 5 Nibbles Into D1
$\qquad$

| opcode: | 1Fnnnnn |
| :--- | :---: |
| cycles: | 7 |

Load all five nibbles of D1 with nnnnn. Any overflow is ignored by the assembler. The assembled digits of nnnnn are stored in the opcode in reverse order so that when the instruction is executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.

```
D1=A - Copy A to D1 (nibs 0-4)
```

opcode: 131
cycles: 8

The A field of register A is copied into Data pointer register D1. Carry is not affected.

D1=AS - Copy A to D1 short (nibs 0-3)
opcode: 139
cycles: 7

The lower 4 nibbles of A are copied into the lower 4 nibbles of Data pointer register D1. Carry is not affected.

D1=C - Copy C to D1 (nibs 0-4)
----.-----
opcode: 135
cycles: 8

The A field of register $C$ is copied into Data pointer register D1. Carry is not affected.

```
D1=CS - Copy C to D1 short (nibs 0-3)
---------
```

    opcode: 13D
    cycles: 7
    The lower 4 nibbles of $C$ are copied into the lower 4 nibbles of Data pointer register D1. Carry is not affected.

D1=D1+ $n$ - Add $n$ to D1 ( $1<=n<=16$ )

```
-----.-.-
    opcode: 17x (x=n-1)
    cycles: 7
Increment D1 by n. This instruction is always executed in HEX
mode. Adjusts Carry.
```

D1=D1- n - Subtract n from D1 ( $1<=\mathrm{n}<=16$ )
$\qquad$
opcode: 1Cx ( $\mathrm{x}=\mathrm{n}-1$ )
cycles: 7

Decrement D1 by $n$. This instruction is always executed in HEX mode. Adjusts Carry.

D1=HEX hh - Load D1 with hex constant hh
---------
opcode: 1Dhh
cycles: 4
Load the low-order two nibbles of D1 with the hex constant hh. The upper nibbles of D1 remain unchanged. The digits of hh are stored in the opcode in reverse order so that when the instruction is
executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.

D1=HEX hhhh - Load D1 with hex constant hhhh

| opcode: | 1Ehhhh |
| :---: | :---: |
| cycles: | 6 |

Load the low-order four nibbles of $D 1$ with the hex constant hhhh. The upper nibble of D1 remains unchanged. The digits of hhhh are stored in the opcode in reverse order so that when the instruction is executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.

D1=HEX hhhhh - Load D1 with hex constant hhhhh
$\qquad$
opcode: 1Fhhhhh
cycles: 7
Load all five nibbles of D1 with the hex constant hhhhh. The digits of hhhhh are stored in the opcode in reverse order so that when the instruction is executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.

```
D=-D fs - Two's complement of D into D
fs=A opcode: FB
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: BbB
    cycles: }3+
```

Complement the specified fs field of D. Complement is two's

```
complement if in HEX mode, ten's complement if in DEC mode. Carry
is set if the field is not zero, else Carry is cleared.
```

| $\mathrm{fs}=\mathrm{A}$ | opcode: <br> cycles: | $\begin{array}{r} \mathrm{FF} \\ 7 \end{array}$ |
| :---: | :---: | :---: |
| $f s=(P, W P, X S, X, S, M, B, W)$ | opcode: cycles: | $\mathrm{BbF}$ $3+d$ |

Perform a one's complement on the specified fs field of D. Carry is always cleared.

```
D=0 fs - Set D equal to 0
fs = A opcode: D3
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: Ab3
    cycles: 3 + d
```

Set the specified fs field of $D$ to zero. Carry is not affected.
$D=C$ fs - Copy $C$ to $D$
fs $=\mathrm{A} \quad$ opcode: D7
cycles: 7
fs $=(P, W P, X S, X, S, M, B, W) \quad$ opcode: Ab7
cycles: $3+d$

Copy the fsield of register $C$ into the corresponding field of register D. Carry is not affected.

```
D=C-D fs - C minus D into D
fs}=
opcode: EF
cycles: 7
fs}=(P,WP,XS,X,S,M,B,W) opcode: Ba
cycles: }3+
```

Set the specified fs field of register $D$ to the inverse difference between itself and the corresponding field of register C. Adjusts Carry.
$D=D!C$ fs $-D O R C$ into $D$
------
fs $=\mathrm{A}$ opcode: OEFB cycles: $\quad 4+d$
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: OEaB cycles: $\quad 4+d$

Set the fs field of register $D$ to its logical OR with the corresponding field of register C. Carry is not affected.
$\mathrm{D}=\mathrm{D} \& \mathrm{C}$ fs - D AND C into D
fs $=\mathrm{A}$ opcode: OEF3
cycles: $\quad 4+d$
$f_{S}=(P, W P, X S, X, S, M, B, W) \quad$ opcode: OEa3
cycles: $\quad 4+d$
Set the fs field of register D to its logical AND with the corresponding field of register $C$. Carry is not affected.

```
D=D+1 fs - Increment D
---------
fs = A opcode: E7
    cycles: 7
fs = (P,WP,XS,X,S,M,B,W) opcode: Ba7
    cycles: }\quad3+
Increment the specified fs field of register D by one. Adjusts
Carry.
```

```
D=D+C fs - Sum of D and C into D
fs=A opcode: C3
cycles: 7
fs}=(P,WP,XS,X,S,M,B,W) opcode: Aa3
cycles: }3+
```

Set the specified fs field of register $D$ to the sum of itself and the corresponding field of register C. Adjusts Carry.

```
D=D+D fs - Sum of D and D into D
```

```
fs = A opcode: C7
    cycles: 7
fs}=(P,WP,XS,X,S,M,B,W) opcode: Aa.
    cycles: }3+
```

Double the specified is field of register D. Adjusts Carry.

```
D=D-1 fs - Decrement D
*
fs = A
fs = (P,WP,XS,X,S,M,B,W) opcode: AaF
cycles: 3 + d
Decrement the specified fs field of register D by one. Adjusts
Carry.
D=D-C fs - D minus C into D
---------
fs = A opcode: E3
    cycles: 7
fs}=(P,WP,XS,X,S,M,B,W) opcode: Ba
    cycles: }3+
Set the specified fs field of register \(D\) to the difference between itself and the corresponding field of register C. Adjusts Carry.
```

```
DATO=A fsd - Store into memory from A
----------
fs =A opcode: 140
cycles: 17
fs}=\textrm{B
opcode: 148
cycles: 14
fs = (P,WP,XS,X,S,M,W) opcode: 150a
cycles: 16 + d
fs = d opcode: 158x (x=d-1)
cycles: 15 + d
```

The amount of data (d nibbles) specified by fsd will be written to the memory address pointed to by DO from the specified field of register $A$. The lowest-order nibble of the register field will be written to the lowest-addressed nibble of memory, proceeding toward the higher-order nibbles. If $f s=d, d$ nibbles are written to

```
memory starting from nibble 0 of the register. See the section on "Storing Data Into Memory" earlier in this chapter.
```

$D A T O=C$ fsd - Store into memory from $C$
---------
$f s=A$
$f s=B$
$f s=(P, W P, X S, X, S, M, W)$
$f s=d$
opcode: 144
cycles: 17
opcode: 14C cycles: 14
opcode: 154a
cycles: $16+d$
opcode: $15 \mathrm{Cx}(\mathrm{x}=\mathrm{d}-1)$
cycles: $15+d$

The amount of data (d nibbles) specified by fsd will be written to the memory address pointed to by DO from the specified field of register C. The lowest-order nibble of the register field will be written to the lowest-addressed nibble of memory, proceeding toward the higher-order nibbles. If fs $=d, d$ nibbles are written to memory starting from nibble 0 of the register. See the section on "Storing Data Into Memory" earlier in this chapter.

```
DAT1=A fs - Store into memory from A
fs = A opcode: 141
    cycles: 17
fs = B opcode: 149
cycles: 14
fs = (P,WP,XS,X,S,M,W) opcode: 151a
    cycles: }16+
fs = d opcode: 159x (x=d-1)
cycles: 15 + d
```

The amount of data (d nibbles) specified by fsd will be written to the memory address pointed to by D1 from the specified field of
register A. The lowest-order nibble of the register field will be written to the lowest-addressed nibble of memory, proceeding toward the higher-order nibbles. If fs $=d$, $d$ nibbles are written to memory starting from nibble 0 of the register. See the section on "Storing Data Into Memory" earlier in this chapter.

```
DAT1=C fsd - Store into memory from C
--------
fs = A opcode: 145
cycles: 17
fs = B
opcode: 14D
cycles: 14
fs = (P,WP,XS,X,S,M,W) opcode: 155a
    cycles: 16 + d
fs = d opcode: 15Dx (x=d-1)
cycles: 15 + d
```

The amount of data (d nibbles) specified by fsd will be written to the memory address pointed to by D1 from the specified field of register C. The lowest-order nibble of the register field will be written to the lowest-addressed nibble of memory, proceeding toward the higher-order nibbles. If $f=d$, $d$ nibbles are written to memory starting from nibble 0 of the register. See the section on "Storing Data Into Memory" earlier in this chapter.

DCEX fs - Exchange Registers $D$ and $C$
fs $=\mathrm{A} \quad$ opeode: DF
cycles: 7
$\mathrm{fs}=(\mathrm{P}, \mathrm{WP}, \mathrm{XS}, \mathrm{X}, \mathrm{S}, \mathrm{M}, \mathrm{B}, \mathrm{W}) \quad$ opcode: AbF

$$
\text { cycles: } \quad 3+d
$$

Exchange the fis fields of registers of D and C. Carry is not a.ffected.

```
DSL fs - D Shift Left
```

-----
fs $=\mathrm{A} \quad$ opcode: F3
cycles: 7
$f s=(P, W P, X S, X, S, M, B, W) \quad$ opcode: Bb3
cycles: $\quad 3+d$

Shift the contents of the specified fs field of register D left one nibble, without affecting the rest of the register. The nibble shifted off the left end of the field is lost. The new low-order nibble of the field is zero. The Sticky Bit (SB) is not affected.

## DSLC - D Shift Left Circular

$\qquad$
opcode: 813
cycles: 21
Circular shift register $D$ left one nibble. Operates on all 16 digits. The Sticky Bit (SB) is not affected.

DSR fs - D Shift Right
---------
fs = A opcode: F7
cycles: 7
$\mathrm{fs}=(\mathrm{P}, \mathrm{WP}, \mathrm{XS}, \mathrm{X}, \mathrm{S}, \mathrm{M}, \mathrm{B}, \mathrm{W}) \quad$ opcode: Bb7
cycles: $3+d$
Shift the contents of the specified fs rield of register $D$ right one nibble, without affecting the rest of the register. The nibble shifted off the right end of the field is lost, but the Sticky Bit (SB) is set if the nibble was non-zero. The new high-order nibble of the field is zero.

Shift register D right one bit. Operates on all 16 digits. The bit shifted off the end is lost, but the Sticky Bit (SB) is set if it was non-zero. The new high-order bit of the register is zero.

DSRC - D Shift Right Circular

| opcode: | 817 |
| ---: | ---: |
| cycles: | 21 |

Circular shift register $D$ right one nibble. Operates on all 16 digits. The Sticky Bit (SB) is set if the nibble shifted from low-order around to high-order position was non-zero.

GOC label - Go relative on carry
----------
opcode: 4 aa
cycles: 10 (GO)
3 (NO)
Short relative jump to label if Carry is set. label must be in the range:

$$
\text { addr }-128<=\text { label }<=\text { addr }+127
$$

where addr is the address of the second nibhle of the opcode. The address offset aa is in two's complement form and is relative to addr.

## GOLONG label - Go Long

opcode: 8 Caaaa
cycles: 14

Long relative jump to label unconditionally. label must be in the range:

```
    addr - 32768 <= label <= addr + 32767
```

where addr is the address of the third nibble of the opcode. The address offset aaa is in two's complement form and is relative to addr.

GONC label - Go relative on no carry


Short relative jump to label if Carry is clear. label must be in the range:

```
    addr - 128 <= label <= addr + 127
```

where addr is the address of the second nibble of the opcode. The address offset aa is in two's complement form and is relative to addr.

GOSBVL label - Gosub very long to label
opcode: 8Faaaaa
cycles: 15

Absolute subroutine jump to aaaaa, which is the absolute address of label. See the GOSUB mnemonic.

GOSUB label - Gosub to label
$\qquad$
opcode: 7aaa
cycles: 12
Relative subroutine jump to label. label must be in the range:

```
addr - 2048 <= label <= addr + 2047
```

where addr is the starting address of the next instruction. The address offset aaa is in two's complement form and is relative to addr.

As with all subroutine jumps, the address (addr) of the instruction following the gosub opcode is pushed onto the hardware return stack, so that when a corresponding return is executed, control resumes with the instruction at address addr.

As the return address is pushed onto the return stack, the bottom-most address on the stack is discarded. Therefore, the return stack always contains 8 addresses, and if pushes exceed pops by 8 levels, the bottom-most return addresses are lost. Since the interrupt system requires one level to process interrupts, only 7 levels of the return stack can be used by code which must execute when interrupts are enabled. See the RTN mnemonic for further information.

GOSUBL label - Gosub long to label

> opcode: BEaaaa
> cycles: 15

Long relative subroutine jump to label. label must be in the range:

$$
\text { addr }-32768 \leq=\text { label }<=\text { addr }+32767
$$

where addr is the starting adaress of the next instruction. The address offset aaa is in two's complement form and is relative to addr. See the GOSUB memonic.

GOTO label - Jump relative
$\qquad$
opcode: 6aaa
cycles: 11

Relative jump to label unconditionally. label must be in the range:

```
addr - 2048 <= label <= addr + 2047
```

where addr is the address of the second nibble of the opcode. The address offset aaa is in two's complement form and is relative to addr.

GOVLNG label - Jump very long


#### Abstract

opcode: 8Daaaaa cycles: 14 Unconditional jump to aaaa, which is the absolute address of label.


## GOYES label - Jump if Test is True

opcode: yy
cycles: included in the accompaning Test mnemonic cycle time.

GOYES is a mnemonic to specify part of a CPU test opcode. GOYES must always follow a test mnemonic. If the condition of the test is met, a jump is performed to label with Carry set. label must be in the range

```
addr - 128 <= label <= addr + 127
```

where addr is the starting address of the jump offset yy. If the test condition is not met, Carry is cleared and control passes to the next instruction. Compare with RTNYES.

```
INTOFF
\(\qquad\)
opcode: 808 F
cycles: 5

Disable the keyboard interrupt system. However, INTOFF does not disable the "ON" key or "INT line interrupts. See the "HP-71 Hardware Specification" for more information.

INTON - Interrupt On
--------
opcode: 8080
cycles: 5
Enable the keyboard interrupt system. See the "HP-71 Hardware Specification" for more information.
\(\operatorname{LC}(m)\) n..n - Load \(C\) with constant ( \(1<=m<=6\) )
-----------
\(\begin{array}{lc}\text { opcode: } & 3 \times n \ldots n(x=m-1) \\ \text { cycles: } & 3+m\end{array}\)
Load \(m\) digits of the expression \(n . . n\) to the \(C\) register beginning at the \(P\) pointer position, and proceeding toward higher-order nibbles, with the ability to wrap around the register. See the section on "Loading Data From Memory" earlier in this chapter.

LCASC \(\backslash\) A..A \(\backslash\) Load \(C\) with ASCII constant
opcode: 3mc..c
( \(\mathrm{m}=2^{*}\) (\# of chars) -1 ;
c..c = ASCII codes)
cycles: \(3+2^{*}(\#\) of chars)
Load up to 8 ASCII characters to the \(C\) register beginning at the \(P\)
pointer position, and proceeding toward higher-order nibbles, with the ability to wrap around the register. Each A represents an ASCII character. The ASCII characters are stored in the opcode in reverse order so that when the instruction is executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.

LCHEX h..h - Load \(C\) with hex constant
```

opcode: 3nh..h (n=\# of digits-1)

```
cycles: \(4+n\)

Load up to 16 hex digits into the \(C\) register beginning at the \(P\) pointer position, and proceeding toward higher-order nibbles, with the ability to wrap around the register. The hex digits are stored in the opcode in reverse order so that when the instruction is executed the data will be loaded into the register with the intended orientation. See the section on "Loading Data From Memory" earlier in this chapter.

MP=0 - Clear Module Pulled bit (MP)
\(\begin{array}{rr}\text { opcode: } & 828 \\ \text { cycles: } & 3\end{array}\)

Clears the Module Pulled bit (MP) and pulls the Module Pulled Interrupt line low. See CLRHST mnemonic.

\section*{HP-71 Hardware IDS -- Detailed Design Description}

NOP3 - Three nibble No-op
```

opcode: 420
cycles: 10 (GO/RTNYES)
3(NO)

```

This mnemonic generates a GOC to the next instruction, effectively skipping three nibbles.

NOP4 - Four nibble No-op
```

----

```
    opcode: 6300
    cycles: 11

This mnemonic generates a GOTO to the next instruction, effectively skipping four nibbles.

NOP5 - Five nibble No-op
--
opcode: 64000
cycles: 11
This mnemonic generates a relative GOTO to +4 nibbles. The fifth nibble in the opcode is a place holder and is jumped over. The mnemonic effectively skips five nibbles.

OUT=C - Load 3 nibbles of OR
\(\qquad\)
opcode: 801
cycles: 6
All nibbles of the Output register are loaded with the low-order three nibbles of \(C\) ( \(X\) field).
```

OUT=CS - Load 1 nibble of OR

```
\begin{tabular}{lr} 
opcode: & 800 \\
cycles: & 4
\end{tabular}

The least significant nibble of the Output register is loaded with the least significant nibble of the \(C\) register.
```

P=C n - Copy P pointer from C at Nibble n

```
---------
    opcode: 80Dn
    cycles: 6
Copy nibble \(n\) of register \(C\) into the \(P\) pointer.
\(\mathrm{P}=\mathrm{P}+1\) - Increment P Pointer
--------
opcode: OC
cycles: 3
Increment the \(P\) pointer. If \(P\) is incremented past \(F\) it will automatically wrap around to 0 . This instruction is always executed in HEX mode. Adjusts carry.
\(\mathrm{P}=\mathrm{P}-1 \quad\) - Decrement P Pointer
---------
opcode: OD
cycles: 3
Decrement the \(P\) pointer. If \(P\) is decremented past 0 it automatically wraps around to F. This instruction is always executed in HEX mode. Adjusts Carry.
```

P= n - Set P Pointer to n
---------

```
    opcode: 2n
    cycles: 2
Set the \(P\) pointer to \(n\).
\(\mathrm{RO}=\mathrm{A}\) - Copy A to register RO
---------
                                    opcode: 100
                                    cycles: 19

The contents of the working register \(A\) is copied to the scratch register RO.
```

RO=C - Copy C to register R0

```
----------
    opcode: 108
    cycles: 19

The contents of the working register \(C\) is copied to the scratch register RO.

R1=A - Copy A to register R1
opcode: 101
cycles: 19
The contents of the working register \(A\) is copied \%o the scratch register R1.
```

R1=C - Copy C to register R1
---------

```
    opcode: 109
    cycles: 19
The contents of the working register \(C\) is copied to the scratch
register R1.
```

R2=A - Copy A to register R2
---------
opcode: 102
cycles: 19
The contents of the working register A is copied to the scratch
register R2.

```
\(R 2=C \quad-C o p y C\) to register R2---------
opcode: ..... 10A
cycles: ..... 19
The contents of the working register \(C\) is copied to the scratchregister R2.
R3=A - Copy A to register R3----------
opcode: ..... 103
cycles: ..... 19

The contents of the working register \(A\) is copied to the scratch register R3.

\section*{HP-71 Hardware IDS -- Detailed Design Description}
```

R3=C - Copy C to register R3

```
---.-.-.
    opcode: 10B
    cycles: 19

The contents of the working register \(C\) is copied to the scratch register R3.

R4=A - Copy A to register R4
opcode: 104
cycles: 19

The contents of the working register \(A\) is copied to the scratch register R4.

R4=C - Copy C to register R4
\(\qquad\)
opcode: 10C
cycles: 19

The contents of the working register \(C\) is copied to the scratch register R4.

\section*{RESET - System reset}
opcode: 80A
cycles:
6

The System Reset Bus Command is issued with all chips performing a local reset. The reset function will vary according to the chip type. See the "HP-71 Hardware Specification" for more information.
```

RSIK=C - Push C to Return Stack

```
```

opcode: 06
cycles: 8

```

Push the low-order 5 nibbles (A field) of the \(C\) register onto the Return Stack. The \(C\) register remains unchanged. See the GOSUB mnemonic.
```

RTI - Return from interrupt

```
---------
opcode: OF
cycles: 9
Return and re-enable the interrupt system. See the RTN mnemonic, and also the "HP-71 Hardware Specification."

RTN - Return
--------
opcode: 01
cycles: 9
Return control to the top address on the hardware return stack. The top address on the hardware return stack is popped off and placed in the program counter PC. As the address is popped of \(f\) the stack, a zero address is inserted at the bottom of the stack.

Therefore the hardware return stack always contains 8 addresses, and if more pops (returns) than pushes (gosubs) are performed, zeros will be read off the stack. Such an attempt to "return" to address 0 results in a memory reset, since the memory reset code of the operating system resides at address 0 .
RTNC -Return on carry
opcode: 400cycles: 10 (RTN)
3 (NO)
Return if Carry is set. See RTN mnemonic.
RTNCC - Return, clear carry
--------
opcode: ..... 03
cycles: ..... 9
Return and clear Carry. See RTN mnemonic.
RTNNC - Return on no carry
opcode: 500 cycles: 10 (RTN) 3 (NO)
Return if Carry is not set. See R'IN mnemonic.
RTNSC - Return, set carry
opcode: ..... 02
cycles: ..... 9
Return and set Carry. See R'N mnemonic.
```

RTNSXM - Return, set External Module Missing bit (XM)
---------
opcode: 00
cycles: 9
Return and set the External Module Missing bit (XM). Since the
opcode is zero, this mnemonic is executed on a jump to a
non-existent memory device. See the "HP-71 Hardware Specification"
for more information. See also the RTN mnemonic.

```
RTNYES - Return if Test is True
opcode: 00
cycles: included in the accompaning
mnemonic cycle time.

RTNYES is a mnemonic to specify part of a CFU test opcode. RTNYES must always follow a test mnemonic. If the test condition is met, Carry is set and a return is executed. If the test condition is not met, control passes to the instruction following the RTNYES. Compare with the RTN and GOYES mnemonics.
\(\mathrm{SB}=0 \quad\) - Clear Sticky Bit (SB)
opcode: 822
cycles: 3

Clear the Sticky Bit (SB). See CLRHST mnemonic.

\section*{SETDEC - Set decimal}
---.-.---

\section*{opcode: 05}
cycles: 3
Set CPU arithmetic mode to decimal.

SETHEX - Set hexadecimal mode
opcode: 04 cycles: 3

Set CPU arithmetic mode to hexadecimal.

SHUIDN - System Shutdown
opcode: 807
cycles: 5
When this mnemonic is executed the CPU sends out the Shutdown Bus Command and stops its clock. Issuing the SHUTDN command with the least significant bit of the output register equal to 0 (i.e., if bit0=0) will immediately cause a cold start -- the PC is set to 0 and the CPU is not halted. This action is to insure that the ON key will be able to wake up the CPU. See the "HP-71 Hardware Specification" for more information.
\(S R=0 \quad\) - Clear Service Request bit (SR)
---------
opcode: 824
cycles: 3
Clear the Service Request bit (SR). See the CLRHST mnemonic.
```

SREQ? - Service Request

```
---------
opcode: 80E
cycles: 7

This mnemonic sets the Service Request bit (SR) if any chip on the system bus requests service. When it is executed, a Service Request Bus Command is issued on the system bus to poll all chips for a Service Request. If any chip requests service, a bus line will be pulled high during the next strobe following the Service Request Bus Command. This value of the bus will be latched into the least significant nibble of the \(C\) register. The bus line pulled high determines the device type according to the following table.

Bit Device
--- ------------------------
3 Unused
2 Card Reader
1 HP-IL Mailbox
0 Display Driver (timer)
If any bus line is high, the Service Request bit (SR) will be set. See the "HP-71 Hardware Specification" for more information. See also the ?SREQ and \(S R=0\) mnemonics.
```

ST=0 n - Clear Program Status bit n
---------
opcode: 84n
cycles: 4

```

Clear the Program Status bit selected by \(n\).

ST=1 \(n\) - Set Program Status bit \(n\)
```

opcode: 85n

```
cycles: \(4_{4}\)

Set the Program Status bit selected by \(n\).
```

ST=C - C to Status

```
```

            opcode: OA
                    cycles: 6
    Copy the low-order 12 bits of the C register (X field) into the
low-order 12 bits of the Status register.

```
UNCNFG - Unconfigure
opcode: 804
cycles: 12
Load the low-order 5 nibbles (A field) of the \(C\) register into the data address register of each peripheral chip, with the device addressed by the \(C\) register unconfiguring. See the "HP-71 Hardware Specification" for more information.
XM=0 - Clear External Module Missing bit (XM)
opcode: 821
cycles: ..... 3Clear the External Module Missing bit (XM). This hardware statusbit is set by the RTNSXM mnemonic. See the CLRHST mnemonic.


This chapter contains a description of the 1LF3 display driver chip. The 1LF3 is designed to support the 1LF2 CPU and future processors that operate on the HP-71 bus. The display driver chip is divided into 4 functional areas:
1) An 8-way multiplexed Liquid Crystal Display (LCD) driver capable of functioning with other 1LF3's in a multi-chip display driver system.
2) 1 K nibbles of RAM.
3) A 24 bit, quartz-crystal controlled timer.
4) A low battery indicator.

\subsection*{5.1 Pin Designations}

The display driver's external pins are as follows:
\begin{tabular}{ll} 
PIN & DESCRIPTION \\
VDD & Power supply. \\
GND & System ground. \\
BUS [0:3] & System bus. \\
"STR & "STROBE \\
"CD & "CONTROL-DATA \\
OD & \begin{tabular}{l} 
OUTPUT DISABLE - When high the bus is tristated; \\
pulled low by an internal resistor.
\end{tabular} \\
CB \((0,1)\) & \begin{tabular}{l} 
CONFIGURATION BITs - 2 bits used for address
\end{tabular} \\
"CLK &
\end{tabular}
\begin{tabular}{ll} 
\#DON & "DISPLAY ON - Driven by MASTER; active low. \\
VREF & LCD voltage reference; driven by MASTER. \\
C[1:40] & 40 LCD column drivers. \\
R[1:8] & \begin{tabular}{l}
8 LCD row drivers on MASTER; 8 LCD column \\
drivers on SLAVE.
\end{tabular} \\
OSC \((A, B)\) & 2 pins used for quartz crystal connection.
\end{tabular}

\subsection*{5.2 Bus Commands}

The display driver is a hard configured chip with service request capability. It responds to the bus commands listed in section 2.1.2 with the following exceptions:

1 ID Same as NOP
8 CONFIGURE Same as NOP

9 UNCONFIGURE Same as NOP
A POLL If the timer has timed out (MSB=1) BUSO is pulled high during the next \({ }^{*}\) STR low (see section 2.4).

C BUSCC Same as NOP

E SHUIDOWN If the timer has timed out or times out during shutdown, *CD is pulled low to wake up the CPU (see section 2.4).

F RESET Same as NOP

\subsection*{5.3 Addressing}

The 1 K nibble of RAM requires 10 bits of address space, leaving 10 bits of configuration address. The most significant 8 bits of the RAM configuration are hard programmed. The other 2 bits are specified by CBO and CB1.

The display RAM and timer are configured together and require 8 bits of address space, leaving 12 bits of configuration address. The most significant 10 bits of the display-timer configuration are
hard programmed and the other 2 bits are specified by CBO and CB1.
The address space is mapped as shown below in hex. The reserved space contains no memory and cannot be used elsewhere in the system.
\begin{tabular}{|c|c|c|c|}
\hline & SLAVE1 & Slave2 & MASTER \\
\hline \multicolumn{4}{|l|}{RAM (1024 nibbles)} \\
\hline Starting addr = & 2F400 & 2F800 & 2FC00 \\
\hline Ending addr & 2F7FF & 2 FBFF & 2 FFFF \\
\hline \multicolumn{4}{|l|}{Display RAM (96 nibbles)} \\
\hline Starting addr = & 2 E 100 & 2E200 & 2E300 \\
\hline Ending addr & 2E15F & 2E2FF & 2E35F \\
\hline \multicolumn{4}{|l|}{RESERVED (152 nibbles)} \\
\hline Starting addr = & 2E160 & 2E260 & 2E360 \\
\hline Ending addr & 2EAF7 & 2E2F7 & 2E3F7 \\
\hline \multicolumn{4}{|l|}{Timer (6 nibbles - least significant first)} \\
\hline Starting addr = & 2E1F8 & 2E2F8 & 2E3F8 \\
\hline Ending addr & 2EIED & 2E2FD & 2E3FD \\
\hline \multicolumn{4}{|l|}{Contrast Control Nibble (1 nibble - active only on MASTER)} \\
\hline Addr location \(=\) & 2E1FE & 2E2FE & 2E3FE \\
\hline \multicolumn{4}{|l|}{Display-Timer Control Nibble (1 nibble)} \\
\hline Addr location \(=\) & 2E/FF & 2E2FF & 2E3FF \\
\hline
\end{tabular}

\subsection*{5.4 Display Interface}

The display driver chips function in one of 2 modes: MASTER or SLAVE. The MAS'EER has both CB1 and CBO tied high by the PC board. It drives the timing signal *CLK, the control signal *DON, and the LCD voltage reference VREF. The MASTER's outputs \(R[1: 8]\) drive the 8 rows of the LCD and it's outputs C[1:40] drive the 40 LCD columns on the right side of the display. The SLAVE chips accept \({ }^{*}\) CLK, *DON, and *VREF as inputs. SLAVE1 has CBO tied high and CB1 tied low and it's outputs \(\mathrm{R}[1: 8]\) and \(\mathrm{C}[1: 40]\) drive the 48 LCD columns in the center of the display. SLAVE2 has CBO tied low and CB1 tied high and it's outputs \(R[1: 8]\) and \(C[1: 40]\) drive the 48 LCD columns on the left side of the display.

Each dot of the LCD has a corresponding bit in the display RAM. If that bit is a 1 and the display is on, the corresponding dot will be on.

16 nibbles of the MASTER's display RAM are allocated to row driver data used to control the row driver waveforms. The following data pattern is loaded by the operating system beginning at the hex address 2E350 of the MASTER.

Hex data (low addressed nibble first) \(=8001400220041008\)
If the above data pattern has been loaded, the least significant bit of the low addressed nibble of the display RAM will correspond to the upper-left dot of the display; the next bit of that nibble will be mapped to just below that dot; and dot addresses will continue to increase down and to the right.

The data pattern shown is not the only valid pattern. For example, try entering : POKE '2E350','0180024004200810'. This pattern defines the rows in reverse order and will result in characters (and annunciators) being displayed upside-down. This can be fixed be either poking the correct pattern or by an INIT:1.

\subsection*{5.5 Contrast Control Nibble}

A contrast control nibble is located at address 2E3FE of the MASTER. This nibble adjusts the LCD drive voltage, VREF and thus can be used to control the contrast of the display. The higher the value of the contrast control nibble the darker the dots appear. This nibble can be both read and written and has no effect if the chip is configured as a SLAVE.

\subsection*{5.6 Display-Timer Control Nibble}

A display-timer control nibble is located at address 2EnFF of both MASTER and SLAVE chips. The bits of this control nibble are defined as follows. Note that some bits have different meanings for read and write.
```

BITO - READ \& WRITE : DISPLAY ON; MASTER only; 1=on.
BIT1 - READ \& WRITE : DISPLAY BLINK; MASTER only; 1=blink
if BITO is set.
BIT2 - READ : VERY LOW BATTIERY INDICATOR; MASTER only;
1=very low.
WRITE : DISPLAY TEST; WARNING !!! DO NOT SET THIS BIT
ON MASTER, IT MAY FORCE DISPLAY OUTC OF SYNC.
MASTER only; When set *CLK is synchronized
with *STR (for testing).

```

BIT3 - READ : LOW BATMERY INDICATOR; MASTER only; 1=low. WRITE : TIMER ENABLE; 1=enable.

\subsection*{5.7 RAM}

The 1LF3 RAM consists of 4096 bits of static hard configured memory arranged as 10244 -bit nibbles.

\subsection*{5.8 Timer}

The timer is a 6 nibble read/writeable binary counter located at addresses 2EnF8-2EnFD. The timer is decremented 512 times per second by the *CLK signal. The "CLK signal is derived from the MASTER's 32768 Hz quartz crystal oscillator and is also used to time display output signals. The MASTER chip requires a 32768 Hz crystal connected between OSCA and OSCB.

IMPORTANT: Because the timer decrement is asynchronous with che timing on the HP-71 bus, it is possible that a decrement will occur between the reading or writing of nibbles to the timer. This can cause erroneous values to be read from or written to the timer. It can be avoided by reading the timer twice within \(1 / 512\) second and verifying that both values are the same. If the values are not the same, immediately read the timer a third time. To write the timer, repeatedly read the least significant nibble until it decrements, then immediately write the timer value.

If the timer's most significant bit is a 1 and it is enabled by BIT3 of the display-timer control nibble the chip will respond to a service POLL by pulling BUS[0] 1ow and it will wake up the CPU during shutdown by pulling *CD low. Aftar waking the CPU the chip will not wake the CPU again until either the MSB of the timer or BII'3 of the control nibble has been cleared and set again.

The HP-71 operating system uses the timers to implement the real time clock system. For information on the clock system see the HP-71 software IDS, Volume 1.

\subsection*{5.9 Low Battery Indicator}

The low battery indicator senses 2 power supply levels: low battery and very low lattery. The 2 bits LBI and VLBI are read from BIT2 and BIT3 of the display-timer control nibble and are valid only on the MASTER. After each read of the control nibble a
new sample of the supply voltage begins. This sample requires
loous. After the sample is complete the LBI and VLBI bits are
updated. If the control nibble is read before the sample is
complete the old low battery values are read and a new sample
begins.

The HP-71 operating system samples the LBI bit once each minute. If it is true, the BAT annunciator will be lit. The operating system ignores the VLBI bit.


The 1LG7 ROM is designed to support the 1LF2 CPU and future processors that operated on the HP-71 bus. The ROM core consists of 128 K bits of memory arranged as 32,768 four-bit nibbles.

\subsection*{6.1 Pin Designations}

The pins of the 1LG7 ROM chip are as follows:
PIN FUNCTION
--- --------

VDD Power Supply.
GND System Ground.
BUS [0:3] System bus.
*STR *STROBE
*CD \({ }^{*}\) COMMAND-DATA
OD OUTPUT DISABLE - when driven high, the ROM tristates the BUS[0:3]; OD is passively pulled low on-chip by an internal resister.

DIN Daisy chain input.
DOUT Daisy chain output.

In the HP-71 system, the OD pins of the 4 system ROMs are tied together. This signal is available at PORT1, and with special hardware, at PORT3 and the HP-IL port. By pulling the OD line high, all 4 system ROMs are effectively removed from the bus.

\subsection*{6.2 Bus Commands}

The 1LG7 ROM can be either hard or soft configured and has no service request capability. The ROM responds to the bus commands
described in section 2.1 .2 with the following exceptions:

1 ID If hard configured : same as NOP.
4 PC WRITE The ROM increments its local program counter once each data strobe. No write is performed.

5 DP WRITE The ROM increments its local data pointer once each data strobe. No write is performed.

8 CONFIGURE If hard configured : same as NOP.
9 UNCONFIGURE If hard configured : same as NOP.
A POLL Same as NOP.
\(C\) BUSCC Same as NOP.
E SHUTDOWN Same as NOP.

F RESET If hard configured : same as NOP.
The ROM's ID code is hard programmed (if soft configured). Generally, the ID code will be either 0010A, for one ROM of a multiple ROM set, or 8010A for an individual ROM or the last ROM of a multiple ROM set.

\subsection*{6.3 Addressing}

The 32 K nibbles of ROM require 15 bits of address space, leaving 5 bits of configuration address. The chip is selected when the upper 5 bits of the PC or DP (whichever is active) match the 5 bits stored in its configuration register and its configuration flag is set. The chip uses the remaining 15 bits to address its memory.

The ROM chip is manufactured in both soft and hard configured options (see section 2.2). In the hard configured option the 5 bits of the configuration register as well as the configuration flag are mask programmed to configure the ROM chip to a fixed a.ddress. The HP-71 operating system is stored in 4 mainframe ROMs configured as follows:
\begin{tabular}{lllll} 
& ROMO & ROM1 & ROM2 & ROM3 \\
Starting addr \(=\) & 00000 & 08000 & 10000 & 18000 \\
Ending addr & \(=\) & \(07 F F F\) & OFFFF & 1 17FFF \\
1FFFF
\end{tabular}

\section*{HP-71 Hardware IDS - Detailed Design Description}

Some plug-in ROMs are soft configured. In the soft configured option, the configuration register latchs the configuration address under software control as described in section 2.2.1.


The 1LG8 RAM is designed to support the 1LF2 CPU and future processors that operate on the HP-71 bus. The RAM core consists of 8 K bits of static memory arranged as 2048 four-bit nibbles.

\subsection*{7.1 Pin Designations}

The pins of the 1LG8 RAM chip are identical to the pins of the 1LG7 ROM chip (see section 6.1) with the following addition:

ID If tied high on the hybrid PC board ( last chip) the most significant bit of the most significant nibble of the 5-nibble ID code will be set to a 1 .

\subsection*{7.2 Bus Commands}

The 1LG8 RAM is soft configured and has no service request capability. It responds to the bus commands as described in section 2.1 .2 with the following exceptions:

A POLL Same as NOP.
C BUSCC Same as NOP.
E SHUTDOWN Same as NOP.
The RAM ID code is nOOOE, where \(n=0\) if the ID \(p\) in is tied low and \(n=8\) if the ID pin is tied high. The ID pin is tied high only on the last chip of the 4 -chip hybrid.

\subsection*{7.3 Addressing}

The 2 K nibbles of RAM require 11 bits of address space, leaving 9 bits of configuration address. The chip is selected when the upper 9 bits of the PC or DP (whichever is active) match the 9 bits stored in its configuration register and its configuration flag is
set. The chip then uses the remaining 11 bits to address its memory. As an example of addressing, if a 4-chip RAM hybrid has been configured contiguously starting at 30000 hex, the following would apply:
RAMO RAM1 RAM2 RAM3
\begin{tabular}{llllll} 
Starting addr & \(=\) & 30000 & 30800 & 31000 & 31800 \\
Ending addr & \(=\) & 307 FF & 30 FFF & 317 FF & 31 FFF
\end{tabular}


Figure 8-1. System Block Diagram

Table 8-1. Top-Case Assembly Replaceable Parts
\begin{tabular}{|c|c|c|c|}
\hline \begin{tabular}{l}
INDEX \\
NUMBER, \\
FIGURE
\[
7-1
\]
\end{tabular} & HP PART NUMBER & DESCRIPTION & QUANTITY \\
\hline C4, C5 & 0160-5789 & CAPACITOR, \(33 \mathrm{pF}, 50 \mathrm{~V}, 5 \%\) & 2 \\
\hline C6, C7 & 0160-5790 & CAPACITOR, 0.1 uF, 25V, \(20 \%\) & 2 \\
\hline C9, C10 & 0160-5787 & CAPACITOR, \(1000 \mathrm{pF}, 5 \mathrm{~V}, 20 \%\) & 2 \\
\hline C11, C14 & 0160-5788 & CAPACITOR, \(220 \mathrm{pF}, \pm 5 \%\) & 2 \\
\hline L1 & 9140-0802 & INDUCTOR, \(180 \mathrm{uH}, 5 \%\) & 1 \\
\hline R20, R21 & 0699-1141 & RESISTOR, \(10 \mathrm{~K}, 5 \%, 1 / 8 \mathrm{~W}\) & 2 \\
\hline Y1 & 0410-1381 & CRYSTAL, quartz & 1 \\
\hline
\end{tabular}

Table 8-2. I/O Assembly Replaceable Parts
\begin{tabular}{|c|c|c|c|}
\hline \begin{tabular}{l}
INDEX \\
NUMBER, \\
FIGURE
\[
7-2
\]
\end{tabular} & HP PART NUMBER & DESCRIPTION & QUANTITY \\
\hline C1 & 0180-3351 & CAPACITOR, 470 uF, \(10 \mathrm{~V}, 20 \%\) & 1 \\
\hline C2 & 0180-3352 & CAPACITOR, \(330 \mathrm{uF}, 16 \mathrm{~V}, 20 \%\) & 1 \\
\hline C8 & 0160-0576 & CAPACITOR, 0.1 uF, \(50 \mathrm{~V}, 20 \%\) & 1 \\
\hline C 12 & 0160-3879 & CAPACITOR, \(0.01 \mathrm{UF}, 20 \%\) & 1 \\
\hline C13 & 0160-4441 & CAPACITOR, \(0.47 \mathrm{uF}, 10 \%\) & 1 \\
\hline CR2 & 1901-0999 & DIODE, Schottky & 1 \\
\hline CR3 & 1906-0069 & DIODE, bridge, full-wave, 400V & 1 \\
\hline CR7 & 1901-0704 & RECTIFIER, silicon & 4 \\
\hline L2 & 9140-0794 & INDUCTOR, \(56 \mathrm{mH}, 10 \%\) & 1 \\
\hline Q1 & 1854-0932 & TRANSISTOR, NPN & 1 \\
\hline Q2 & 1854-0973 & TRANSISTOR, & 1 \\
\hline R22 & 0683-3915 & RESISTOR, 390 ohms, \(5 \%, 0.25 \mathrm{~W}\) & 1 \\
\hline R23,R24 & 0683-1035 & RESISTOR, 10 K -ohms, \(5 \%\), 0.25 W & 2 \\
\hline VR7 & 1902-1390 & DIODE, Zener & 1 \\
\hline
\end{tabular}


Figure 8-2. Top-case Assembly Component Location Diagram



Figure 8-4. I/O PCA Component Location Diagram




Figure 8-6. HP-71 Mainframe
[.x.
8-7/8-8




Figure 8-8. HP-71 Memory Module

HP-71 Hardware IDS - Detailed Design Description


Figure 8-9. HP-71 I/O Contact


Figure 8-10. HP-71 Contact Spring Attachment Diagrams


HP-71 Hardware IDS - Detailed Design Description


Figure 8-12. HP-71 HP-IL Module


Figure 8-13. HP-71 Card Reader Port 8-21/8-22

HP-71 Hardware IDS - Detailed Design Description

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Figure 8-14. HP-71 Card Reade
8-23/8-24

SYM PARAMETER MIN TYP MAX UNIT COMMENTS

POWER SUPPLY
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline Vac & AC adapter input voltage & 9.0 & - & 14.0 & V & See Note 1. \\
\hline Vreg & AC adapter voltage at Vdd & 5.6 & - & 6.5 & VDC & \\
\hline Iac & AC adapter current drain & - & - & 30 & mADC & Over mainframe requirements. \\
\hline Vbat & Battery voltage at Vdd & 4.25 & - & 6.0 & VDC & Unregulated. \\
\hline Ibat & Battery current drain & - & - & 100 & mADC & Over mainframe requirements. \\
\hline
\end{tabular}

BUS PARAMETERS
Vih Input logic Vdd-.65 -
level '1'
Vil Input logic - \(\quad\) Vss+. 65 V
level '0'

Voh Output logic Vdd-.5 level '1'

Vol Output logic - - Vss+. 5 V level 'o'

Cout Output capacitance drive capability ( surplus )

SYM PARAMETER MIN TYP MAX UNIT COMMENTS

Ioh High level output source current capability
\begin{tabular}{llllll} 
BUSO-3 & 1.0 & - & - & \(m A D C\) & See Note 3. \\
\({ }^{*} C D,{ }^{H}\) STR & 1.0 & - & - & mADC See Note 3. \\
DIN & 100 & - & - & \(u A D C\)
\end{tabular}

Iol Low level output sink current capability
\begin{tabular}{lrllll} 
BUSO-3 & 1.0 & - & - & \(m A D C\) & See Note 3. \\
\({ }^{\text {HCD, }}\) *STR & 1.0 & - & - & mADC See Note 3. \\
DIN & 2 & - & - & uADC
\end{tabular}

Cin Input capacitance loading
\begin{tabular}{llllll} 
BUSO-3 & - & - & 300 & pF & See Note 2. \\
\({ }^{\text {HCD, }}\) "STR & - & - & 250 & pF & See Note 2. \\
\#INT & - & - & 1200 & pF & \\
IR14 & - & - & 250 & pF & \\
HALT & - & - & 250 & pF \\
OD & - & - & .012 & uF
\end{tabular}

Iih High level input current loading
\begin{tabular}{llllll} 
BUSO-3 & - & - & 50 & uADC & Internal pulldown \\
"CD & - & 0.0 & - & uADC & Internal pullup \\
\#STR & - & 0.0 & - & uADC & Internal pullup \\
HINT & - & 0.0 & - & uADC & 10K pullup \\
IR14 & - & - & 50 & uADC & Internal pulldown \\
HALT & - & - & 1.0 & mADC & 1OK pulldown \\
OD & - & - & 50 & uADC & Internal pulldown
\end{tabular}
SYM PARAMETER MIN TYP MAX UNIT COMMENTS

Iil Low level input current loading
\begin{tabular}{lccccl} 
BUSO-3 & - & 0.0 & - & uADC & Internal pulldown \\
\#CD & - & - & 50 & uADC & Internal pullup \\
\#STR & - & - & 50 & uADC & Internal pullup \\
\#INT & - & - & 1.0 & mADC & 10K pullup \\
IR14 & - & 0.0 & - & uADC & Internal pulldown \\
HALT & - & 0.0 & - & uADC & 10K pulldown \\
OD & - & 0.0 & - & uADC & Internal pulldown
\end{tabular}

BUS TIMING PARAMETERS
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline T0 & \begin{tabular}{l}
*STR cycle \\
time
\end{tabular} & 1.0 & - & - & us & Current HP-71's operate at 600 KHz to 650 KHz , future \\
\hline Tpwl & *STR low & 0.5 & - & - & us & HP-71's will run at up to 1 MHz . \\
\hline Tpwh & *STR high & 0.5 & - & - & us & \\
\hline Tdwc & Data-in valid to *STR high & 200 & - & - & ns & Command cycle \\
\hline Tdwd & \begin{tabular}{l}
Data-in \\
valid to \\
*STR high
\end{tabular} & 100 & - & - & ns & Write cycle \\
\hline Tdh & *STR high to data-in invalid & 100 & - & - & ns & \\
\hline Tacc & *STR low to data-out valid & - & - & 200 & ns & BUS precharged low \\
\hline Toh & \begin{tabular}{l}
*STR high \\
to data-out \\
tristated
\end{tabular} & 20 & - & 100 & ns & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline SYM & PARAMETER & MIN & TYP & MAX & UNIT & COMMENTS \\
\hline Tc1 & \begin{tabular}{l}
*CD low to \\
*STR low
\end{tabular} & 30 & - & - & ns & \\
\hline Tc2 & \begin{tabular}{l}
*STR high to \\
*CD high
\end{tabular} & 50 & - & - & ns & \\
\hline Tc3 & \begin{tabular}{l}
*CD high to \\
*STR low
\end{tabular} & 100 & - & - & ns & \\
\hline \(\operatorname{Tr}\) & Rise time & & & & & \\
\hline & *STR & - & - & 100 & nS & \\
\hline & \({ }^{*} \mathrm{CD}\) & - & - & 100 & \(n \mathrm{~S}\) & \\
\hline & BUSO-3 & - & - & 100 & nS & \\
\hline Tf & Fall time & & & & & \\
\hline & *STR & - & - & 100 & nS & \\
\hline & *CD & - & - & 100 & nS & \\
\hline & BUSO-3 & - & - & 100 & nS & \\
\hline
\end{tabular}

\section*{COMMAND CYCLE}

CD


\section*{READ CYCLE}

\section*{\(\overline{\text { STR }}\)}


DATA DRIVEN BY OTHER SYSTEM IC

CPU
RECHARGE


\section*{WRITE CYCLE}

\section*{\(\overline{S T R}\)}


DATA

\begin{tabular}{|c|c|c|c|c|c|c|}
\hline SYM & PARAMETER & MIN & TYP & MAX & UNIT & COMMENTS \\
\hline Fclk & Display *CLK frequency & - & 512 & - & Hz & \begin{tabular}{l}
\[
+-40 \mathrm{PPM}
\] \\
Real time base.
\end{tabular} \\
\hline \multicolumn{7}{|l|}{LOW BAITERY INDICATION} \\
\hline Vlbi & Low battery trip point & 4.3 & - & 4.5 & V & Low battery bit high when Vdd < Vlbi. \\
\hline Vdta & Vlbi - Very low battery trip point & . 08 & - & . 12 & V & Very low battery bit high when Vdd < Vlbi-Vdta. \\
\hline
\end{tabular}
1) The values specified for Vac limit the peak voltage that may be applied between the 2 AC adapter input pins. This voltage may be either an AC peak voltage or a DC voltage. If a voltage less than 9.0 V is applied the unit's batteries may be discharged. This minimum value for Vac (9.0V) can be reduced to 7.5 V if the batteries are removed from the unit.
2) The values specified under Cout and Cin for BUSO-3, *CD, and *STR allow for 100pF loading by plug-in modules. This can be broken down as:
a) RAM or ROM modules \(=20 \mathrm{pF} \times 4\)
b) HP-IL module \(=12 \mathrm{pF}\)
c) Card reader module \(=8 \mathrm{pF}\)

A HP-71B with no plug-in modules could drive 100 pF more than specified and would load these lines 100 pF less than specified.
3) Reduce Cout ( output capacitance drive capability ) by \(25 \mathrm{pF} / \mathrm{mA}\) of the larger of Ioh and Iol (high level source and low level sink current capability) required of BUSO-3, *CD, and/or \({ }^{*}\) STR. Under no condition should the minimum current capability specified for Ioh and Iol be exceeded.
4) The DIN lines of the I/O ports are driven by the CPU output register ( \(O R\) ). The OR lines are also used to form the keyboard matrix and therefore excessive loading of a DIN line will prevent proper keyboard operation. The maximum capacitive loading of a DIN line may be increased to 50 pF if a pulldown resistor, Rpd, of \(50 \mathrm{Kohm}(+-10 \%\) ) is tied from DIN to ground. For most of the time when the HP-71 is turned on all lines of the \(O R\) are high and thus Rpd will draw current. When the HP-71 is shut off the operating system sets ORO high and all other OR lines low. ORO drives the DIN line of Port 1.

\section*{(h) HEWLETT PACKARD}

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